# Lecture 1 <br> Introduction \& Dimensionality 

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## Announcements \& Agenda

## - Announcements

- People Introductions
- Web site https://cioffi-group.stanford.edu/ee392aa/
- Chapters 1-5 are used, on-line at class web site (Course Reader)
- Review/scan Section 1.3.4-7 ; read 2.1-5 ; 4.1-3
- Chapter 3 (not necessary, equalization and ISI)
- Supplementary files at canvas for your interest/review (contact Yun if


John M. Cioffi Room 363, David Packard

Welcome Course Info Course Reader Lecture Notes Handouts Homework Matlab Code
Spring Quarter 2023
EE 392AA - Multiuser Data Transmission

Instructor : Prof. John Cioffi
Teaching Assistant: Yun Liao
Course Secretary : Helen Niu
Lectures : Monday and Wednesday, 15:30-16:45, in class interested in special section)

## Today

- Course introduction
- The scalar AWGN channel (a foundation)
- The matrix AWGN channel
- Water-filling energy distribution
- Projecting forward

| - Problem Set $1=$ PS1 due Wednesday April 12 at 17:00 |  |  |
| ---: | :--- | :--- |
| 1. | 2.15 | capacity refresher (read "subsymbol" $=$ "symbol" here) |
| 2. | 4.3 | builds intuition on gap-based 1-dimensional channel analysis |
| 3. 4.18$\quad$ DMT water-fill loading |  |  |
| 4. | 4.7 | Simple Water-fill Loading |
| 5. | 4.25 | Matrix AWGN \& vector coding with water-fill |

## Why Communications?

## Next Generation Connectivity



Rural, less-developed connect

Digital Twins used to forecast/emulate each


Defense ("5/6G.mil")
Samsung: 6G "hyper connected"

## Broadband Internet Access (\$1.5T/year)

- Messages
- Internet
- Email
- Text
- video, audio
- Sensor/camera images


OSI Model

this class
April 3, 2023

## VR/AR focus and Bandwidth

## VR and AR require efficient increase in wireless capacity

Constant up/download on an all-day wearable


Richer visual content

- Higher resolution, higher frame rate



## Latency:

Edge $\sim 1 \mathrm{~ms}$
ISP Cloud 20-50 ms

Public Cloud 100 ms

Popular Com Standards Summaries


## Course Introduction

## Communications Depth Sequence

Sig./Syst. II EE 102B


- Modulation and Coding compliment one another
- Modulation = energy assignment to time/frequency/space, is separate from:
- Coding = distinct message mapping
- If both done well, they separate


## Basic Communication (digital)



- The symbol $\boldsymbol{x}$ and messages are in some 1-to-1 relationship
- Finding the best $\widehat{\boldsymbol{x}}$ and designing $\boldsymbol{x}$ well $\rightarrow$ this class (good 1-to-1 assumed)
- Most general channel is represented by the conditional probability $p_{y / x}$.
- Most general source description is $p_{x}$ - together, $p_{x y}$.
- Optimum detector (minimizes ave error probability) is Maximum a Posteriori (MAP), $\max p_{x / y}$
- When input distribution is uniform $\rightarrow \mathrm{ML}$ (maximum likelihood), $\max p_{y / x}$


## 3 Basic Problems to Solve

- CHANNEL IDENTIFICATION - what is $p_{y / x}$ ?
- CODING \& MODULATION - What are good (best) $\boldsymbol{x}$ and $p_{x}$ for a given channel?
- DETECTION - What is a good (best) receiver for deciding which $\boldsymbol{x}$ ?

Especially with more than 1 user (so expanding on 379A)

## Dimensionality

- Input $\boldsymbol{x}$ and channel $\boldsymbol{y}$ are vectors
- Simple dimensions
- time (samples, slots, packets)
- frequency (carriers, tones/subcarriers, bands)
- space ("antennas")

- Exotic Extensions from Physics
- higher-order modes (TM(m,n))
- orbital angular momentum
- quantum communication


## Communication Dimensionality

## - Time-Frequency (any fixed location)

10

- $2 \times$ bandwidth = \# of dimensions/sec (wireless or wired, including "optical" - all are EM waves)

2G(2 dim/ $3.7 \mu \mathrm{~s})$
$3 \mathrm{G}(128 \mathrm{dim} / 533 \mu \mathrm{~s})$

4G/Wi-Fi (160 dim/4 $\mu \mathrm{s}$ )

- ー ー
- Space-Time
- 2D-3D (at least .....)
- Spacing of half wavelength or more
- $10 \mathrm{k}-1 \mathrm{M}$ dimensions per few microseconds
- Number of channels can be up to \# of antennas "spatial streams"
- Also can be crosstalk between wires (e.g., ethernet, telco lines)
- Optical xtalk is between "modes" at same frequency (a spatial effect)

Same 2-path channel over 320 MHz



Matrix
channel
antennas

## Even More Dimensions (smaller wavelengths)



- How do we design these systems for best rates (per energy) use?
- How adaptive do they need to be?


## The scalar AWGN channel

(a foundation: Section 1.3, Section 2.1-3
direct: 2.4.1, 2.4.3)

See PS1.1 (Prob 2.15 - capacity) and PS1. 2 (Prob 4.3 gap)

## Simple Additive White Gaussian Noise Channel

## Detection Problem First, every $T$ seconds (symbol period)



## SNR, QAM, PAM reminders

$$
S N R \triangleq \frac{\bar{\varepsilon}_{x}}{\sigma^{2}}=\frac{\text { single }- \text { sided psd }}{\text { single }- \text { sided psd }}=\frac{\text { two }- \text { sided psd }}{\text { two }- \text { sided psd }}
$$

- SNR must have the same number of dimensions in numerator (signal) and denominator (noise)
- Thus, also $S N R \triangleq \frac{\bar{\varepsilon}_{x}}{\sigma^{2}}=\frac{2 \cdot \bar{\varepsilon}_{x}}{\mathcal{N}_{0}}=\frac{\varepsilon_{x}}{N \cdot \sigma^{2}} \quad$ where $\bar{\varepsilon}_{x}$ is energy/real-dimension.
- Energy/dimension essentially generalizes the term power/ Hz (= energy) so that is why these quantities are related to powerspectral densities (psd's)
- 1 -sided $\rightarrow$ power is integral over positive frequencies of psd
- 2 -sided $\rightarrow$ power is integral over all frequencies of psd
- These two powers are the same
- So - $40 \mathrm{dBm} / \mathrm{Hz}$ (one-sided) psd over 1 MHz is 20 dBm , or 100 mWatts of power, practice PS1.1 (Prob 2.15) and Homework Helper 1's first part
- PAM is always real baseband. QAM is always complex baseband (2 real dimensions)
- When QAM has only 1 bit (2 points) in constellation, it is called BPSK (not binary PAM).
- PAM's positive-frequency bandwidth is $[0,1 / 2 \mathrm{~T}) \quad \ldots$... $\quad$ (
- QAM's positive-frequency bandwidth is $\left[-1 / 2 T+f_{c}, 1 / 2 T+f_{c}\right)$
- The PAM system looks like it uses only $1 / 2$ the bandwidth, but the QAM system is really transmitting two dimensions per symbol (so really like 2 PAM systems in parallel with symbol rate 1/T each), so no wonder it takes twice the bandwidth of a single PAM to do so


## Codes and Gaps

Shannon's maximum reliable data rate "capacity"

$$
\mathcal{C}=\log _{2}(1+S N R) \text { bits/complex-subsymbol } \quad \text { AWGN Max bits/sub-sym for } P_{e} \rightarrow 0 \text { (reliably decodable) }
$$



$$
\text { bits } / \operatorname{dim}=\bar{b}=b /{ }_{N} ; \text { bits/subsym }=\tilde{b}={ }^{b} /{ }_{N}=\widetilde{N} \cdot \bar{b}
$$

- QAM/PAM operates with given low $P_{e}\left(10^{-6}\right)$ and at a "SNR gap" ( $\Gamma=8.8 \mathrm{~dB} @ 10^{-6}$ ) below capacity
- See basics in Section 1.3.4 - for practice, see Section 2.4; also PS1.2 (Prob 4.3)

$$
\tilde{b}=\log _{2}\left(1+\frac{S N R}{\Gamma}\right) \text { bits/complex-subsymbol } \leq \mathcal{C}
$$

- For all $\tilde{b}>1$, simple square QAM constellations have constant gap ( $=8.8 \mathrm{~dB}$ at $P_{e}=10^{-6}$ )

$$
\frac{3}{2^{\bar{b}}-1} \cdot S N R=13.5 d B \text { (from } P_{e}=10^{-6} \text { formula) }
$$

## Margin

$$
\tilde{b}=\log _{2}\left(1+\frac{S N R}{\Gamma \cdot \gamma_{m}}\right) \text { bits/complex-subsymbol } \leq \mathcal{C}
$$

- The designer wants a little "margin" protection against possible noise-power increase
- MARGIN $\gamma_{m}$ is this protection (usually in dB), $\quad \gamma_{m}=\frac{(S N R / \Gamma)}{2^{\tilde{b}}-1}$


## Positive margin - means performing well ; Negative margin - means not meeting design goals

- AWGN with SNR $=20.5 \mathrm{~dB}$, then $\tilde{\mathcal{C}}=\log _{2}\left(1+10^{2.05}\right)=7 \mathrm{bits} /$ subsymbol
- Suppose that 16-QAM $(\tilde{b}=4)$ is transmitted $@ P_{e}=10^{-6}(\Gamma=8.8 \mathrm{~dB})$, then $\gamma_{m}=\frac{10^{2.05-.88}}{2^{4}-1}=0 \mathrm{~dB}$
- Suppose instead QAM with $\tilde{b}=5 \mathrm{bits} / c o m p l e x-s u b s y m b o l$ with a code and gain 7 dB of gain $(\Gamma \rightarrow 8.8-7=1.8 \mathrm{~dB})$
- $\gamma_{m}=\frac{10^{2.05-.18}}{2^{5}-1}=3.8 \mathrm{~dB}$
- 6 bits/subsymbol with same code? $\rightarrow 0.7 \mathrm{~dB}$ margin - just barely below the desired $P_{e} ; \bar{P}_{e}={ }^{P_{e}} /{ }_{N}$


## EE 392AA (379C)

- The simple single-dimension AWGN is fundamental to most all designs
- All subsequent designs will depend on good codes (small or 0 dB gap) re-use on those single dimension AWGNs
- Designs can be optimized to get highest possible data rates for Gaussian noise
- Single user (of course)
- All multiuser
- Channels with crosstalk between dimensions
- Intersymbol interference
- Crosstalk
- Spatial reflections, multi-paths
- Many users with many antennas, high/low data rates, crosstalking wires and different locations
- This is where the big gains occur


## Gap Plot \& Example

- The gap is constant, independent of the bits/dimension - greatly simplifies "loading" (adapting transmission codes to the channel)



## The Matrix AWGN Channel

## Section 2.3.5

also supplementary lectures S1A and S1B also 379 Help files at Canvas site

## Generating Parallel AWGNs

- Methods from EE379?
- An "equalizer" is one choice
- Parallel channels in time
- $z_{k}=x_{k}-e_{k}$



Section 3.6

- Another?
- Multicarrier is another choice
- Parallel channels in frequency

$$
Y_{n} \cong H_{n} \cdot X_{n} \quad\left(+N_{n}\right)
$$

Sections 1.3.8 and 4.2.1
L1: 23
Stanford University

## In general, a matrix AWGN channel



$$
\varepsilon_{x}=\varepsilon_{x^{\prime}}=\sum_{l=1}^{L_{x}} \varepsilon_{l}
$$

$$
2^{\tilde{b}} \text { possible }
$$ messages


$L_{x}$ dimensions



detected message
$L_{y}$ dimensions

## Vector Coding (MIMO)

$$
\begin{array}{ll}
x_{1}^{\prime} \longrightarrow y_{1} & n_{1} \sim \sigma^{2} \\
x_{L}^{\prime} \longrightarrow y_{1}^{\prime} \\
\lambda_{L} & n_{L} \sim \sigma^{2} \\
&
\end{array}
$$

## $L \leq \min \left(L_{x}, L_{y}\right)$ independent dimensions

$$
R_{n \boldsymbol{n}} \neq I \rightarrow\left(H \rightarrow R_{n \boldsymbol{n}}^{-1 / 2} \cdot H\right)
$$

$$
\begin{aligned}
& H=F \cdot \Lambda \cdot M^{*} \ldots \begin{array}{c}
\text { singular value decomposition (svd in matlab) } \\
F \cdot F^{*}=F^{*} \cdot F=I_{L_{y}} ; M \cdot M^{*}=M^{*} \cdot M=I_{L_{x}}
\end{array} \\
& \Lambda \text {. }\left(L_{y} \times L_{x}\right) \text { is "diagonal" (real) }
\end{aligned}
$$

## Geometric Equivalent Channel




Use it $L$ times like single constant AWGN

- Vector Coding - uses SVD to translate matrix AWGN to set of equivalent parallel AWGN's
- Each can be individually encoded like AWGN (they are independent)
- Geometric-equivalent channel use $L$ times
- Any $H$ and $R_{n n}$
- Any set of input energies (that sum to allowed energy)


## The Detection/Communication Issue

- MAP/ML receiver/detector implementation can be very complex
- An entire body of theory/practice has been devoted to reducing this complexity as well as projecting nearly attainable bounds
$>$ Communication Theory and Information Theory
- Decomposing into multiple channels can simplify design!
> Multiple dimensions are the key to this simplification
$>$ And today, used throughout digital communication (wires, wireless, soon fiber)


# The Water-Filling Energy Distribution 

## Sections 2.3.5, 4.1-4.3

also supplementary lecture S1A

See PS1.3 (Prob 4.18), PS1.4 (Prob 4.7), and PS1.5 (Prob 4.25)

## Rate Maximization and Dual

- Choose energy and bit allocation to maximize sum data rate over the dimensions

$$
\begin{aligned}
& \max _{\varepsilon_{l}} \sum_{l=1}^{L} \log _{2}\left(1+\frac{\varepsilon_{l} \cdot g_{l}}{\Gamma}\right)=\sum_{l=1}^{L} b_{l} \\
& S T: \varepsilon_{x}= \sum_{l=1}^{L_{x}} \varepsilon_{l} \\
& \quad \text { Rate Adaptive (RA) }
\end{aligned}
$$

$$
\begin{array}{lr}
\min _{b_{l}} \sum_{l=1}^{L} \varepsilon_{l} & \text { DUAL } \\
S T: b=\sum_{l=1}^{L_{x}} b_{l} & \\
& \text { Margin Adaptive (MA) }
\end{array}
$$

- Solution (basic calculus - see Section 4.2) ; see also matlab "waterfill.m" at web site to save hand calcs

$$
\varepsilon_{l}+\frac{\Gamma}{g_{l}}=\text { constant } \quad \underset{(\text { Whannon 1948) }}{\text { WATER-FILLING }}
$$

Neither energies allocated nor bits allocated can be negative

## Water-filling Illustrated

- Energy available in a pitcher
- Note re-indexed 0 (DC) to 5

RA: until all energy used
MA: until total bit rate attained

$$
\tilde{b}_{l}=\log _{2}\left(1+\frac{S N R_{l}}{\Gamma}\right)
$$



- Write and sum energy constraints

$$
g_{1} \geq g_{2} \geq \ldots \geq g_{L}
$$

$$
\begin{gathered}
\varepsilon_{1}+\Gamma / g_{1}=K \\
\varepsilon_{2}+\Gamma / g_{2}=K \\
\vdots \\
\varepsilon_{L}+\Gamma / g_{L}=K \\
\sum_{l=1}^{L} \varepsilon_{l}+\Gamma \cdot \sum_{l=1}^{L} 1 / g_{l}=L \cdot K
\end{gathered}
$$

- Solve for Water-Fill Constant

$$
K=\frac{\varepsilon_{x}}{L^{*}}+\frac{\Gamma}{L^{*}} \cdot \sum_{l=1}^{L^{*}} 1 / g_{l}
$$

$L^{*}$ is largest L such that $\varepsilon_{l}>0$ for all $l=1, \ldots, L^{*}$

2 x 2 Antenna System with 0 dB gap


- There is crosstalk between dimensions and $\varepsilon_{x}=2$
- Kind of sounds like a problem then, right?
>> H=[10 4
$21]$;
>> [F , Lambda , Mstar]=svd(H);
>> Lambda =
10.99850
$0 \quad 0.1818$
$\gg$ g2=Lambda $(1,1)^{\wedge} 2=120.9669$
$\gg$ g1=Lambda(2,2)^2 $=0.0331$
$\gg \mathrm{K}=1+0.5^{*}(1 / \mathrm{g} 1+1 / \mathrm{g} 2)=16.1250$
$\gg E 2=K-1 / \mathrm{g} 2=16.1167$
$>$ E1=K-1/g1 = -14.1167<0 (whoops)

Just use dimension $2 \rightarrow \tilde{b}=\log _{2}\left(1+2 * g_{2}\right)=6.93$ bits/subsymbol

In this case water-fill simply puts all energy on the best dimension (returns to scalar/SISO if that is best)


- There is stronger crosstalk between dimensions
- Maybe worse, right? ???
>> H=[10 9
-8 10];
$\gg$ [F, Lambda , Mstar]=svd(H);
>> Lambda =
13.62440
012.6244
$\gg$ g2 $2=\operatorname{Lambda}(1,1)^{\wedge} 2=185.6244$
>> g1=Lambda(2,2)^2 $=159.3756$
$\gg \mathrm{K}=1+0.5^{*}(1 / \mathrm{g} 1+1 / \mathrm{g} 2)=1.0058$
> $\mathrm{E} 2=\mathrm{K}-1 / \mathrm{g} 2=1.0004$
>>E1=K-1/g1 = 0.9996
>> btilde $=\log 2\left(1+E 2^{*} \mathrm{~g} 2\right)+\log 2(1+E 1 * \mathrm{~g} 1)=14.8693$

Actually this is close to $2 x$ the data rate for the previous case
Clearly, the use of both dimensions, and somewhat stronger crosstalk and signal.

In general, the increase is roughly a factor of $L$ in data rate.

## Energy-minimizing Margin-Adaptive Solution

- Energy and sum-bit constraints

$$
g_{1} \geq g_{2} \geq \ldots \geq g_{L}
$$

$$
\begin{aligned}
& \varepsilon_{l}=K-\Gamma / g_{l} \\
& \begin{array}{c}
\tilde{b}=\sum_{l=1}^{L} \tilde{b}_{l}=\sum_{l=1}^{L} \log _{2}\left(1+\frac{\varepsilon_{l} \cdot g_{l}}{\Gamma}\right) \\
=\sum_{l=1}^{L} \log _{2}\left(\frac{K \cdot g_{l}}{\Gamma}\right) \\
=\log _{2}\left(\prod_{l=1}^{L} \frac{K \cdot g_{l}}{\Gamma}\right)
\end{array}
\end{aligned}
$$

- Solve for Water-Fill Constant

$$
K=\Gamma \cdot\left(\frac{2^{\tilde{b}}}{\prod_{l=1}^{L} g_{l}}\right)^{1 / L^{*}}
$$

$L^{*}$ is largest $L$ such that $\varepsilon_{l}>0$ for all $l=1, \ldots, L^{*}$

## 2 x 2 Antenna System with MA



- Attempt $\tilde{b}=14 \frac{\mathrm{bits}}{\mathrm{Hz}}$; The use of 2 antennas exploited channel's crosstalk,
- Without the crosstalk, this channel supports only 7 bits/Hz (either channel has then SNR = 10)


## >> H=[10 9

-8 10];
>> K=sqrt((2^14)/(g1*g2)) $=0.7442$
> E2=K-1/g2 = 0.7388
> E1=K-1/g1 = 0.7379
$\gg$ margin $=10^{*} \log 10(2 /(E 1+E 2))=1.3 \mathrm{~dB}$

This effect magnifies as long as most of the singular values are "decent"

## RA Water-Fill Flow Chart

- Can start with all channels energized
- Compute $K$, test lowest energy
- Reduce number of dimensions incrementally
- Can also start with 1 channel energized
- Compute K, test lowest energy
- Increase number of dimensions incrementally
- The sort is most complex part
- Can use pivots and bi-section
- Avoids sort

Sort subchannels $g_{1} \geq g_{2} \geq \cdots \geq g_{L}$

$$
j=L ; \widetilde{K}=\varepsilon_{x}+\sum_{l=1}^{L} \Gamma / g_{l}
$$



$$
\text { Compute WF Energies } \varepsilon_{l}=K-\Gamma / g_{l} l=1, \ldots, j=L^{*}
$$

Unsort subchannels $\tilde{b}_{l}=\log _{2}\left(1+\frac{\varepsilon_{l} \cdot g_{l}}{\Gamma}\right) \quad l=1, \ldots, j=L^{*}$

## Margin Adaptive Flowchart

$$
\gamma_{\max }=\frac{\varepsilon_{x}}{\sum_{l=1}^{L^{*}} \varepsilon_{l}}
$$

## Projecting Forward

## Water-filling from 100k feet



- How do we learn and adjust either or both of energy/bits per dimension?
> Dynamically
- $\quad$ Some of very first Al methods in communication (from Stanford)
$>$ "bit-swapping"
> \#3 Stanford patent on value/royalty in Engineering


## Multiple directions in Space



Best energies will also be water-fill over the channel's spatial singular vectors

Essentially matrix form of machine learning From earlier


## End Lecture 1

