

#### *Lecture 9* **Finish MAC & Broadcast Channels** *April 30, 2024*

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#### **Announcements & Agenda**

- Announcements
	- Problem Set #4 due today
	- Midterm in class Thursday

- Agenda (L9)
	- Finish L8 MAC examples
		- mu\_mac.m
	- Simultaneous Water-Filling for MAC max rate sum
		- SWF.m and macmax.m
	- **MAC:** Capacity region for frequency-indexed MACs
	- **BC:** Precoder Basics for the Matrix AWGN
	- Scalar Gaussian BC



#### **MAC Examples**

Sections 2.7.3-4

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### **Matrix AWGN MAC Example 1**  $(L_{x,u} \equiv 1)$

 $\widetilde{H} = \left[ \begin{array}{cc} 5 & 2 \\ 3 & 1 \end{array} \right]$ 



- These are the two vertices for dimension-share (pentagon outer face).
- **•** Two receiver output dimensions for each one-dimensional input  $x<sub>u</sub>$  (instead of 1 output dimension earlier)



 $\sim$  Rf H<sup>\*H</sup>

#### **Example 1 continued**

§ Receiver filters and bias are

**Vertex 1**

 $\gg$  W=inv(S0)\*inv(G') = 0.0286 0 -0.3171 0.8537  $\gg$  Wunb=S0\*inv(S0-eye(2))\*W = 0.0294 0 -2.1667 5.8333 >> MSWMFu=Wunb\*H' =

0.1471 0.0882 0.8333 -0.6667 >> Gunb=eye(2)+S0\*inv(S0-eye(2))\*(G $eye(2)$  = 1.0000 0.3824 0 1.0000

#### **Vertex 2**

 $\Rightarrow$  W=inv(S0)\*inv(G') = 0.1667 0 -0.3171 0.1463  $\Rightarrow$  Wunb=S0\*inv(S0-eye(2))\*W = 0.2000 0 -0.3714 0.1714 >> MSWMFu=Wunb\*H' = 0.4000 0.2000 0.1143 0.1429 >> Gunb=eye(2)+S0\*inv(S0-eye(2))\*(G-eye(2)) = 1.0000 2.6000 0 1.0000 >> MSWMFu\*H= 1.0000 2.6000 0.3714 1.0000

Not really triangular, why?<br>2.1667 1.0000 **2.1667** 

 $\gg$  MSWMFu\*H = 1.0000 0.3824



#### **Easier with mu\_mac.m**



 $1$ ]; [b, GU, WU, S0, MSWMFU] = mu\_mac(H, **eye(2)**, [1 1] , 2);

**b = 2.5646 0.1141**

 **1.0000 0.3824 0 1.0000** 0.0294 0 5.8333  $\overline{0}$ 1714

 **0.1471 0.0882 0.8333 -0.6667**  $FU<sup>*</sup>H =$ 0.3824 1.0000

 $10<sup>*</sup> \log 10(\text{diag}(S0)) =$ 

 $= 2.6788$ 

### **Example 2: 2 x 3 MAC (secondary users)**

#### H=[5 2 1

3 1 1]; **basically added a 3rd user**  $[b, GU, WU, SO, MSWMFU] = mu_mac(H, eye(3), [1 1 1], 2)$ 



- The channel rank is 2 so at least 1 secondary comp =  $3-2$ .
- § But secondary applies to energy-sum MAC (which this is not, yet).
- § If original 2 units of energy is spread over 3 users?

```
\gg [b, GU, WU, S0, MSWMFU] = mu_mac(H, (2/3)^*eye(3), [1 1 1], 2)
b = 2.0050 0.1009 0.0696
GU = 1.0000 0.3824 0.2353
    0 1.0000 0.3878
    0 0 1.0000
WU = 0.0662 0 0
  -2.3878 6.6582 0
  -2.0000 -0.5000 9.8750
S<sub>0</sub> = 16.1111 0 0
     0 1.1502 0
     0 0 1.1013
MSWMFU =
  0.2206 0.1324
  0.9184 -0.3367
  -0.7500 2.2500
>> sum(b) = 2.1755 (lower than 2x2 value of 2.6788)
```
Relatively more energy on secondary-user comp(s), bsum  $\downarrow$ .

#### Section 2.7.2.2 April 25, 2024 **PS4.4 - 2.26 MAC regions** L9: 7



#### **Non-Zero Gap Achievable Region**

- **•** Construct  $C(b)$  with  $\Gamma = 0$  dB.
- Reduce all rates by  $\gamma_h$  relative to boundary points.
- **•** Inscribe smaller region  $C(\boldsymbol{b})$  ( $\gamma_b \odot \mathbf{1}$ ).
- Square constellations instead of spheres (AWGN) loss 1.53 dB in gap above (0.25 bit/dimension).





#### **Simultaneous Water-Filling for MAC max rate sum**

Sections 2.7.3-4

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### **Revisit the rate-sum mutual information**

$$
b = \sum_{u=1}^{U} \tilde{b}_u \le \mathbb{I}(\boldsymbol{x}; \boldsymbol{y}) = \log_2 \frac{|H \cdot R_{\boldsymbol{x}\boldsymbol{x}} \cdot H^* + R_{\boldsymbol{n}\boldsymbol{n}}|}{|R_{\boldsymbol{n}\boldsymbol{n}}|}
$$

- § Maximum rate-sum focuses on the numerator,
	- when optimizing over  $R_{xx}$ .

$$
\frac{\max\limits_{\{R_{xx}(u)\}}\left|H_u \cdot R_{xx}(u) \cdot H_u^* + \sum\limits_{i \neq u} H_i \cdot R_{xx}(i) \cdot H_i^* + R_{nn}\right|}{\underbrace{R_{noise}(u)}}
$$

- § Have we seen this problem before?
	- Yes, it is Vector Coding / Waterfilling , except with  $\widetilde H_{\cal u} \to R_{noise}^{-1/2} \cdot H_{\cal u}$  for each  $u$
- But now it repeats  $U$  times in the same form for each user.
	- Optimum has each user **simultaneously water-fill** by treating all other users (water-fill) spectra as noise.



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**Simutaneous water**  
\n
$$
\mathcal{E}_{u,l} + \frac{1}{g_{u,l}} = K_u \,\forall \, u = 1, ..., U'
$$
\n
$$
\sum_{l=1}^{L_x} \mathcal{E}_{u,l} = \mathcal{E}_u
$$
\n
$$
\mathcal{E}_{u,l} \ge 0
$$
\n
$$
x_u = M_u \cdot v_u
$$

**Simultaneous Waterfilling**

## **Compute Using Iterative Water-filling**



- § SWC problem is convex, and each single-water-fill step is "gradient-like" in improving direction, swf.m
- E-Sum SWC is a saddle point with enlarged region.
	- 2nd optimization is on the allocation of  $\mathcal{E}_{x,u} \to \sum_{u=1}^{U} \mathcal{E}_{x,u} = \mathcal{E}_x$ .

Section 2.7.4.1

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## **SWF.m Program for MAC's max sum rate**

function [Rxx, bsum , bsum\_lin] = SWF(Eu, **H**, Lxu, Rnn, cb)

 Simultaneous water-filling MAC max rate sum (linear and nonlinear GDFE) The input is space-time domain h, and the user can specify a temporal block symbol size N (essentially an FFT size).

#### Inputs:

- Eu U x 1 energy/SAMPLE vector. Single scalar equal energy all users any (N/N+nu) scaling should occur BEFORE input to this program.
- **H The FREQUENCY-DOMAIN** Ly x sum(Lx(u)) x N MIMO channel for all users.
- N is determined from size(H) where  $N = #$  used tones
- Lxu 1xU vector of each user's number of antennas
- Rnn The Ly x Ly x N noise-autocorrelation tensor (last index is **per tone**) cb cb = 1 for complex, cb=2 for real baseband
	- cb=2 corresponds to a frequency range at an sampling rate 1/T' of  $[0, 1/2T']$  while with cb=1, it is  $[0, 1/T']$ . The Rnn entered for these two situations may differ, depending on how H is computed.

#### Outputs:

 Rxx A block-diagonal psd matrix with the input autocorrelation for each user on each tone. Rxx has size (sum( $Lx(u)$ ) x sum( $Lx(u)$ ) x N. sum trace(Rxx) over tones and spatial dimensions equal the Eu

bsum the maximum rate sum.

- bsum bsum\_lin the maximum sum rate with a linear receiver
- b is an internal convergence sum rate value, not output

 This program significantly modifies one originally supplied by student Chris Baca

- Eu is each user's energy**/sample.**
- For now,  $N = 1$ , so time/freg are same:
	- $\cdot$  H=h.
- Lxu is number of antennas for each user.
- § Separate specification of Rnn removes need for noise whitening.
- $\blacksquare$  cb=1 for complex, =2 for real.



## **Revisit Previous example (slides L8: 26-29)**

```
H = 5 2 1
    3 1 1
\Rightarrow [Rxx, bsum, bsum_lin] = SWF([1 1 1], H, [1 1 1], eye(2), 2)
Rxx = 1 0 0
           \Omega\overline{\phantom{0}} 1
bsum = 2.7925
bsum lin = 1.4349
```
- § Same result as L8:29, so each user waterfills with all others as noise; this is trivial when each user has only 1 input dimension. (Why?)
- This is for input energy-vector constraint.
- § Note linear solution (no feedback, so matrix MMSE-LE) loses roughly ½ the data rate.
- SWF becomes more interesting when  $N > 1$  tones or if  $L_{x,y} > 1$  antennas.

#### For  $L_{x,u} = 2$ ;  $u = 1,2$ ?



- Note block-diagonal Rxx.
- Linear-only loses about 25% in data rate (for this channel).



## **Or use Macmax.m for Esum MAC**

function [Rxx, bsum , bsum\_lin] = macmax(Eu, **h**, Lxu, N , cb)

 Simultaneous water-filling Esum MAC max rate sum (linear & nonlinear GDFE) The input is space-time domain h, and the user can specify a temporal block symbol size N (essentially an FFT size).

This program uses the CVX package

the inputs are:

- Eu The sum-user energy/SAMPLE scalar.
- This will be increased by the number of tones N by this program. Each user energy should be scaled by N/(N+nu)if there is cyclic prefix This energy is the sum trace of the corresponding users' Rxx (u). The sum energy is computed as the sum of the Eu components internally.
- **h The TIME-DOMAIN** Ly x sum(Lx(u)) x N channel for all users Lxu The number of antennas for each user 1 x U
- N The number of used tones (equally spaced over (0,1/T) at N/T. cb cb = 1 for complex, cb=2 for real baseband

#### the outputs are:

 Rxx A block-diagonal psd matrix with the input autocorrelation for each user on each tone. Rxx has size (sum( $Lx(u)$ ) x sum( $Lx(u)$ ) x N. sum trace(Rxx) over tones and spatial dimensions equal the Eu bsum the maximum rate sum.

bsum bsum lin - the maximum sum rate with a linear receiver

b is an internal convergence (vector, rms) value, but not sum rate



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- § ENERGY-SUM input (per sample)
	- Lxu = numbers of xmit antennas/user
- § Time-domain (noise-whitened) h
- This is actually a double loop that:
	- water-fills each and every user for some current set of per-user energies and
	- adjusts energies so they sum to total but increase the rate sum.
- It corresponds to a saddle point.
	- It is not convex (although each sub loop is).
	- It has a solution and converges anyway.
- This will be easier understood later as a dual of a broadcast problem as to why this is true.

Section 2.7.4.3

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#### **Back to Example**

```
>> H3(:,:,1)=HH = 5 2 1
  3 1 1
>> [Rxx, bmacmax, bmaclin]=macmax(3/2, H, [1 1 1], 1, 2)
Rxx = 3.0000 0 0
    0 0.0000 0
        0 0.0000
bmacmax = 3.3432
bm = 3.3432
```
- This produces a larger data rate because there is less energy restriction.
- Rxx energizes just user 3! (It's all primary user component, and users 1 and 2 are secondary)
- § Linear is the same. Why?



## **Capacity region for frequency-indexed MACs**

Sections 2.7.4.1-2

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## $\mathcal{C}(\bm{b})$  is union of  $\mathcal{S}_x(f)$ -indexed Pentagons



- Each pentagon corresponds to an  $S_{\mathbf{x}}(f)$  choice.
	- The pentagons become triangles for the sum-energy MAC.
- The union (convex hull is union when inputs are Gaussian) can dimension-share in frequency as  $N \to \infty$ .



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#### **T MAC**



- **•** The users have continuous-time/frequency channels  $\rightarrow$  use MT on each, theoretically.
- This really means dimensionality is infinite (or very large) so "dimension-sharing" may be inherent.
- SWF applies, but with some interpretation (like power instead of energy and power per dimension instead of power-spectral density, etc. ).



#### **Decoders and SWF**



- **FDM is clearly simplest decoder for max rate-sum case.**
- Both users (and all components in case c) are primary.



Section 2.7.4.2

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#### **Symmetric 2-user channel and SWF**



- Symmetric means  $H_1(f) = H_2(f)$  (noise is one-dimensional and added to sum)
- Each of points A, B, and C have different SWF spectra all have same (max) rate sum

Section 2.7.4.2

# **Basic Precoders and the Matrix AWGN**

PS5.1 - 2.28 modulo precoding function

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#### **Broadcast Channel (BC)**



- The BC is the "Dual" of the MAC.
- Receivers are in different places and so cannot "co-process"  $\{y_n\}$ .
- **•** Transmitter can co-encode/generate  $x$ , although input messages remain independent.
	- Who encodes first? (may be at disadvantage)
	- Who encodes last? (knowing other users' signals is an advantage)
	- What then is the **order**?



### **BC is "reversed" MAC**





- The MAC's uncoordinated user input is a kind of "worst case" transmitter, reducing data rate.
	- With only an energy-sum constraint, these worst-case inputs' users best pass as primary user components; secondary components "freeload" on the primary's passage.
- The BC similarly will effectively correspond to a worst-case noise for which receiver coordination is useless, reducing data rate.
	- With worst-case noise, the channel best passes the primary components'; secondary components freeload on the primary's passage.

## **Triangular Matrices - Innovations and Prediction**

§ Prediction for some **user order** separates a modulated input to independent message components.

$$
\boldsymbol{v}_u = \boldsymbol{x}_u - \widehat{\boldsymbol{x}}_{u/\{x_{u+1} ... x_U\}}
$$

Innovations or predictions, but for BC  $v_u$  become the independent-users' subsymbols, with normalization  $R_{uv}(u) = I$ .

§ This is a triangular relationship (**inverse of upper triangular is also upper triangular**).

$$
\boldsymbol{\nu} = \begin{bmatrix} \boldsymbol{\nu}_1 \\ \boldsymbol{\nu}_2 \\ \vdots \\ \boldsymbol{\nu}_U \end{bmatrix} = \begin{bmatrix} 1 & g_{1,2} & \dots & g_{1,U} \\ 0 & 1 & \dots & g_{2,U} \\ \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{bmatrix} \begin{bmatrix} \boldsymbol{x}_1 \\ \boldsymbol{x}_2 \\ \vdots \\ \boldsymbol{x}_U \end{bmatrix} = G^{-1} \cdot \boldsymbol{x}
$$

- OR,  $x = G \cdot v$  (G is also upper triangular).
- § *Generating from can increase energy (~ enhance noise in MAC rcvr) if implemented directly (linearly).*

(order reversal is intentional)



Section 2.8.1.1 and D.3.6.1.1 **L9:24** April 30, 2024 **L9:24** 

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# **Voronoi Regions and Modulo Addition (Sec 2.1)**

- A lattice is a (countable) group of vectors  $\Lambda = \{x\}$  that is closed under an operation addition, so that
	- If  $x_1$  ∈  $\Lambda$  and  $x_2$  ∈  $\Lambda$ , then  $x_1 + x_2 \in \Lambda$ . (Section 2.2.1.1 and Appendix B.2)
	- A constellation is a finite subset of a lattice, plus a constant (coset)  $C \subset \Lambda + \lambda_0$ . ( $\lambda_0$  ensures average value is zero.)



- Voronoi Region of a lattice,  $V(\Lambda_c)$  is the decision region around any point with volume  $V(\Lambda_c)$ .
	- $\Lambda_c$  is the "coding" lattice; codes try to pack more points into limited space (volume/area). HEX is better than SQ.
- A constellation C typically selects points in one (coding-gain) lattice,  $\Lambda_c$ , within the  $V(\Lambda_s)$  of another (shaping-gain) lattice  $\Lambda'$  that is larger (can be scaled versions of one another or possibly different). (Subtract any nonzero vector mean to save energy.)
	- All points in  $\Lambda_c$  outside of  $\mathcal{V}(\Lambda_s)$  map into a point inside  $\mathcal{V}(\Lambda_s)$  disguised detector problem.

Section 2.3, App B.2 **L9:25** April 30, 2024 **L9:25** 

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### **More general precoder (than Tomlinson/Laroia …)**

§ Generalize **Modulo Operation:** 

$$
\boxed{(\boldsymbol{v})_{\Lambda_S} = \boldsymbol{e} \ni \min_{\boldsymbol{\lambda} \in \Lambda_S} \|\boldsymbol{e}\|^2 \text{ where } \boldsymbol{e} = \boldsymbol{v} - \boldsymbol{\lambda}}
$$

- **e** does not necessarily need to be a point in  $\Lambda_c$ ; instead, it is a point in  $\mathcal{V}(\Lambda_s)$ .<br>• It's essentially the error between input and output of decoder with decision region equal to  $\mathcal{V}(\Lambda_s)$  sbs in trivial ca
	-

§ useful

**Lemma 2.8.1 (distribution of modulo addition)** Modulo addition distributes as

$$
(\boldsymbol{\mu} + \boldsymbol{\nu})_{\Lambda} = (\boldsymbol{\mu})_{\Lambda} \oplus_{\Lambda} (\boldsymbol{\mu})_{\Lambda} . \qquad (2.371)
$$

#### Side info is s.

Section 2.8.1.2

- *is known (ISI for causal G(D) = in Tomlinson/DFE)*
- Pre-subtract (precode) and use modulo to set  $\mathcal{E}_x$  level (no energy increase for **s**).
- Any  $\Lambda_s$  shaping gain also applies here (8.5,  $s=0$ ).
	- $v$  then is any (e.g., SQ) input constellation.
	- In shaping case,  $\hat{v}$  also needs a subsequent detector yet for whatever code is used on  $v$ .
- x will effectively have continuous uniform distributions over  $\Lambda_s$ .





#### **For BC:** *s* will be the earlier users (but their xtalk cancels) ~ **order.**

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PS 5.1 (2.28) – modulo precoder L9:26

## **With nontrivial channel, need MMSE version**

#### **Forney's Crypto Lemma – 2003 (Section 2.8.1.2)**



- The MMSE part can be important in non-trivial cases (often missed in most info-theory texts).
	- It's reshaping the channel crosstalk and/or ISI in MMSE (not zero-forcing) sense.

**No xmit energy increase Simplifies ML detection**

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- When s is uniform over  $\mathcal{V}(\Lambda_s)$ , then so is  $v$ , AND  $v$  is independent of both s and  $v$  (like encryption), s is the "key"<br>• Or "writing on dirty paper" (s is the dirt, v is the writing, and the second modulo cleans
	-
- Sometimes the channel adds  $s$  (ISI/xtalk), sometimes the transmitter adds  $s$  (xmit case,  $s$  shares dimensions and energy with  $x$ ).
	- The add-at-xmit case has  $s$  as other users, effectively (with a twist  $\ldots$  later).

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#### **Non-Causal ?**

- Subtly, the lattice  $\Lambda_{\rm s}$  has a dimensionality N over which s and x are uniformly distributed.
- Wise dimension use with fixed energy  $\varepsilon_x$  suggests  $\Lambda_s$  has a hyper-spherical boundary, as  $N \to \infty$ .
	- This infinite-length precoder then also obtains full 1.53 shaping gain.
- Asymtotically, the modulo has infinite number of dimensions, so requires infinite delay for  $s$  to be fully known in the formation of  $x$ ; whence "non-causal."
	- Approximated with finite delay in practice, s becomes another user's encoded signal known first ( $\sim$  non-causal)  $\rightarrow$  order.



- Mod holds energy at  $\varepsilon_r$  (Gaussian in any finite number of dimensions, uniform in infinite dimensional hypersphere).
- If  $\Lambda_s$  is hypercube, Forney's crypto still holds but with SNR loss of (up to) 1.53 dB (the maximum shaping gain).
	- So reuse code with  $\Gamma \to 0$  dB, with QAM constellations and the (up to) 1.53 dB loss remains (greatly simplifies precoder implementation),
		- but everything else works the same.



#### **Midterm Review**

- Questions
- Advice





# **End Lecture 9**