



STANFORD

*Lecture 1*

# **Introduction & Dimensionality**

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# Announcements & Agenda

## ■ Announcements

- People Introductions
- Web site <https://cioffi-group.stanford.edu/ee379b/>
- Chapters 1-9 are used, on-line at class web site (Course Reader)
- Read Chapter 4
- EE379A website is also available for review
  - <https://cioffi-group.stanford.edu/ee379a/>

## ■ Today

- Multiuser Communication
- The scalar AWGN channel (a foundation)
- The matrix AWGN channel
- Water-filling energy distribution
- Projecting forward

## ■ Problem Set 1 = PS1 due Wednesday April 7 at 17:00

1. 2.15 capacity refresher
2. 4.3 gap-based 1-dimensional channel analysis
3. 4.18 DMT water-fill loading
4. 4.7 Simple Water-fill Loading
5. 4.25 Matrix AWGN & vector coding with water-fill

Send to [kfardi@stanford.edu](mailto:kfardi@stanford.edu)



# Multiuser Communication

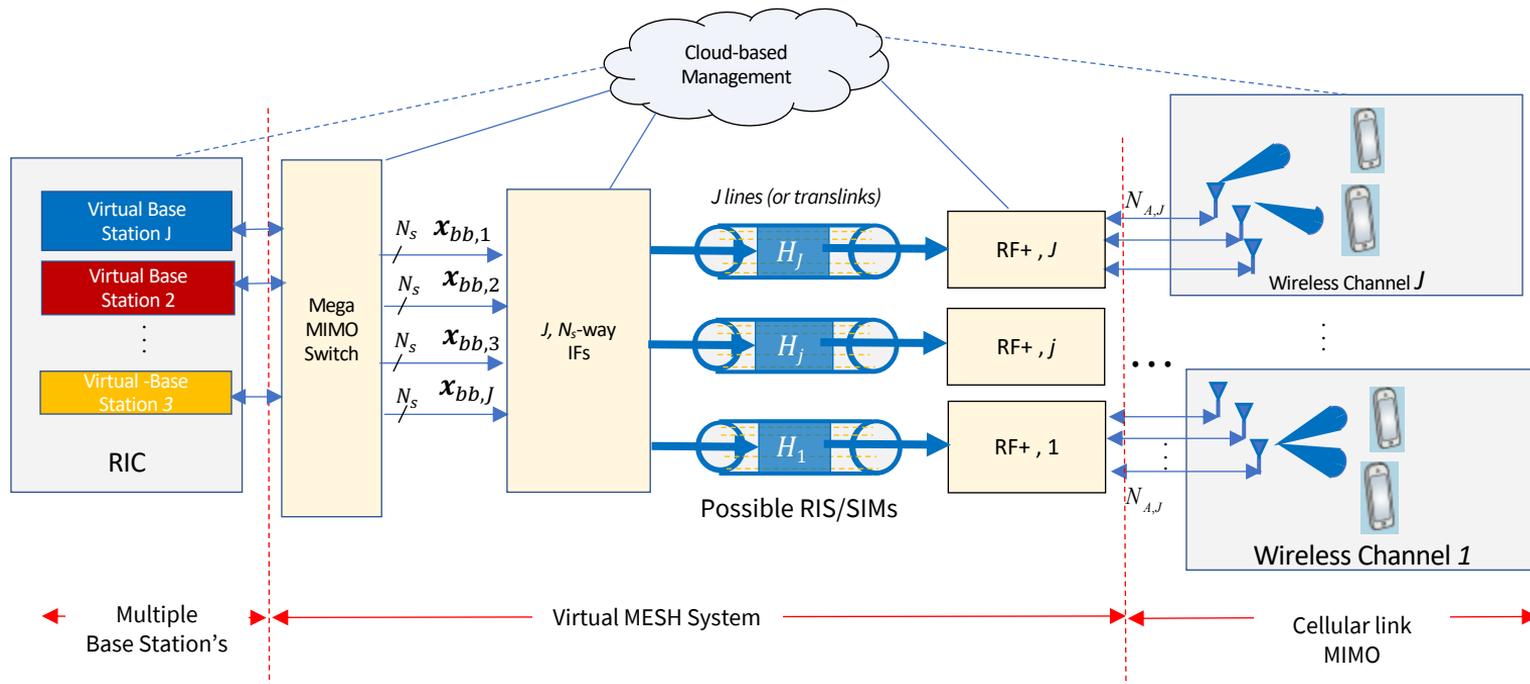
# Multiuser Examples: Broadband and Cellular



- Downlink/stream – one to many is “**broadcast**” multiuser channel.
- Uplink/stream – many to one is **multiple access** multiuser channel.
- Overlapping combinations (Wi-Fi, or cell, or really all) is an “**interference**” channel.
- Relaying the signals makes a **mesh** channel.



# Mega MIMO – Translink Convergence (“Xhaul”)



- This supports “cell-free,” a concept in next-generation systems.
- It exploits “Virtualization” (software modulation, coding too!) moves to data-center/edge.



# Metaverse Distributed Rendering

- Remote rendering: Devices (smart phones, glasses/goggles) used to augment current environment

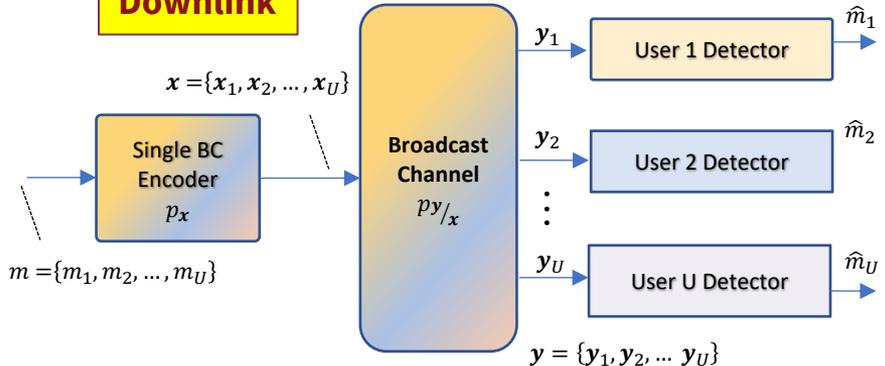


- Games
- Education
  - Instructions
- Health
- Multiple contributors**

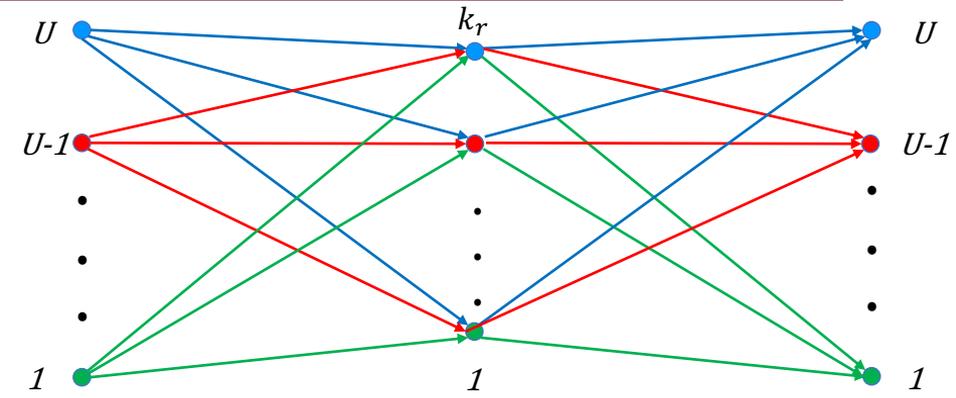
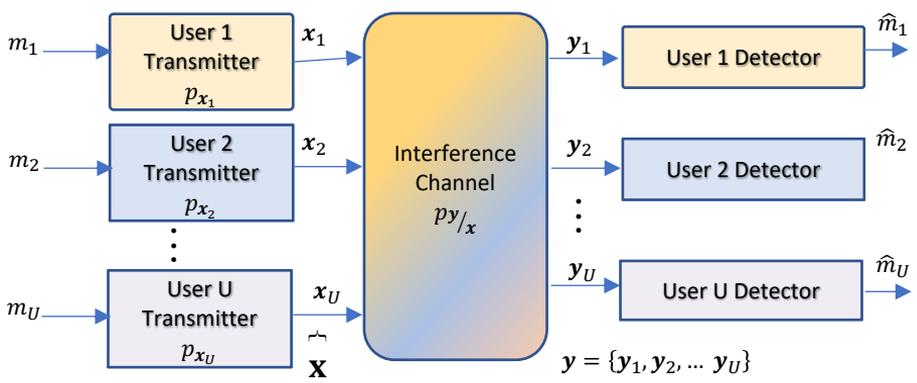
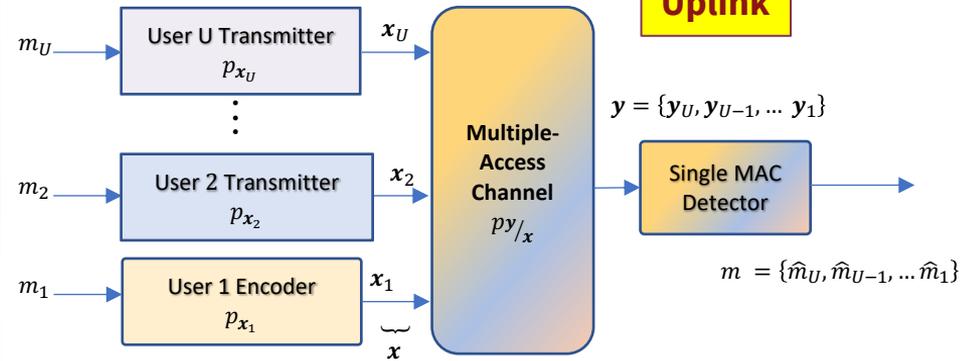


# Multiuser Channel Basics (all others are combos)

## Downlink



## Uplink



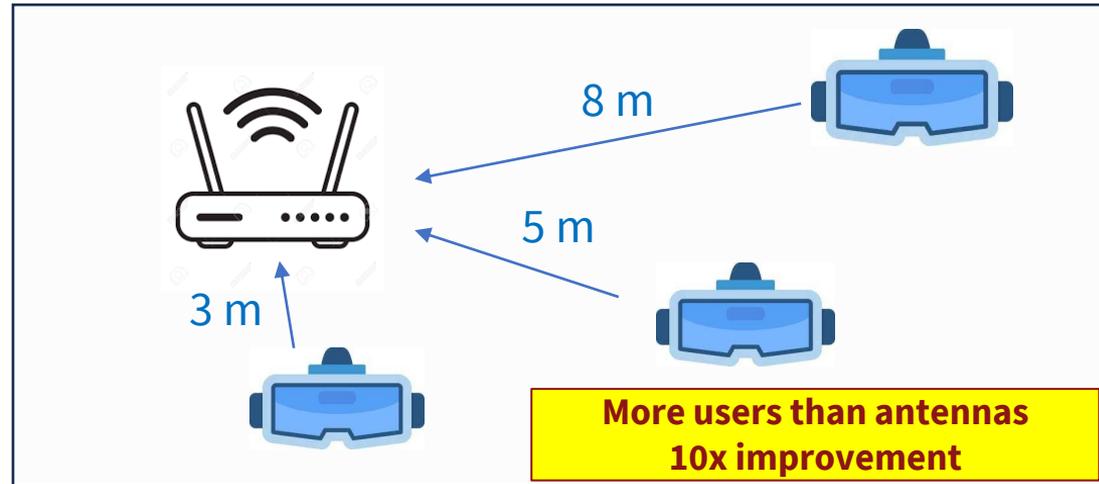
## Mesh/Relay

## Inter-cell (and roaming)

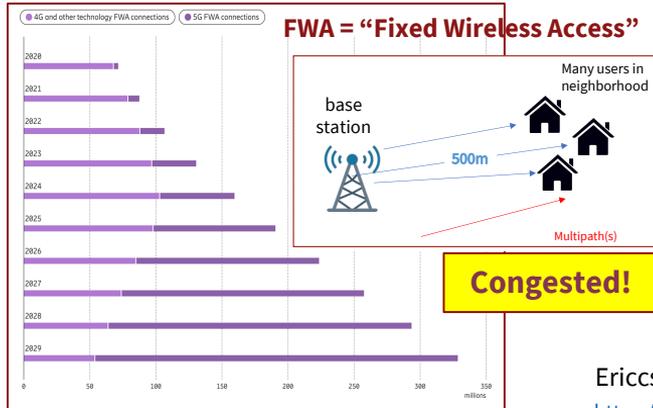


# Wi-Fi Multiuser Channel

- Minimum Distributed AR rate
  - 500 Mbps/user in same 80 MHz channel.



- Similar in Fixed Wireless Access



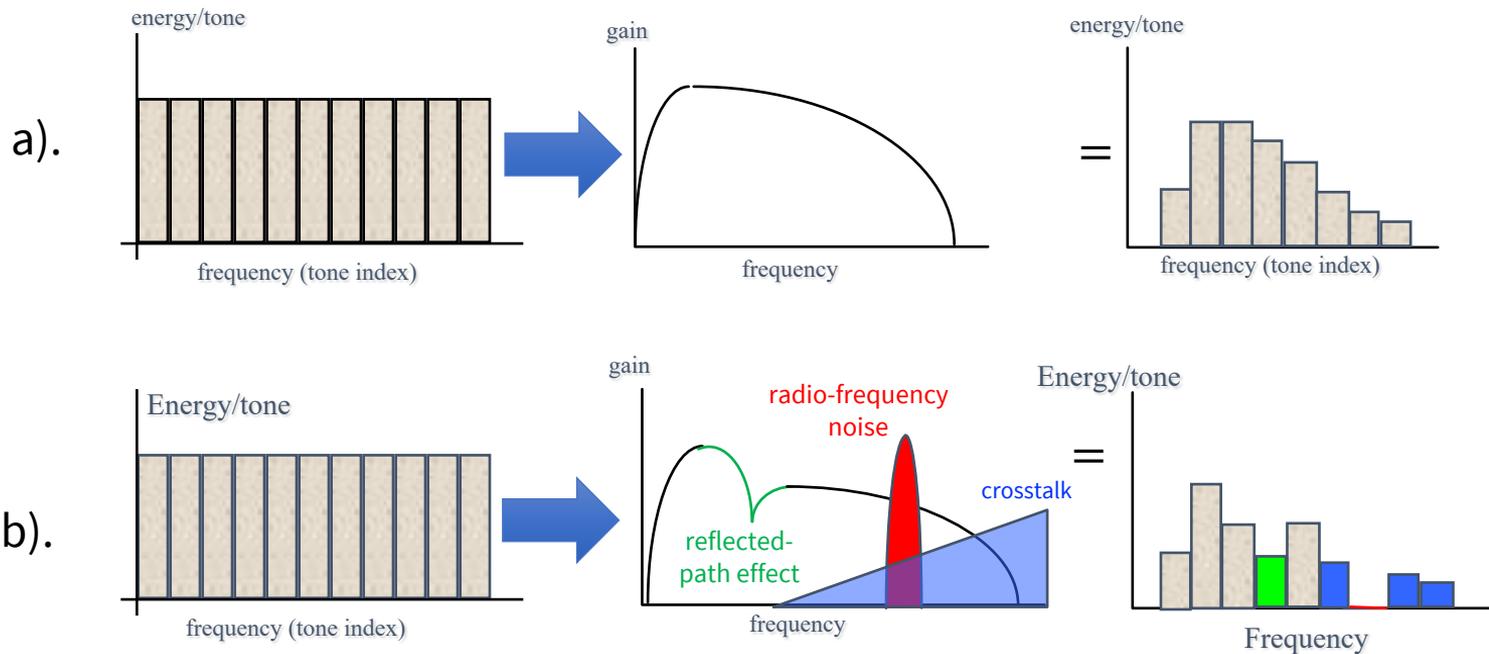
**EE379B has custom Matlab software that optimizes for best wireless performance.**

Ericsson Mobility Report 2024

<https://www.ericsson.com/en/reports-and-papers/mobility-report/dataforecasts/fwa-outlook#:~:text=Over%20330%20million%20FWA%20connections%20by%202029,expected%20to%20be%20over%205G.>

# Multicarrier Adaptive Transmitters

- Dynamic adaptation of transmit resource use to the channel situation, “loading” or resource allocation.

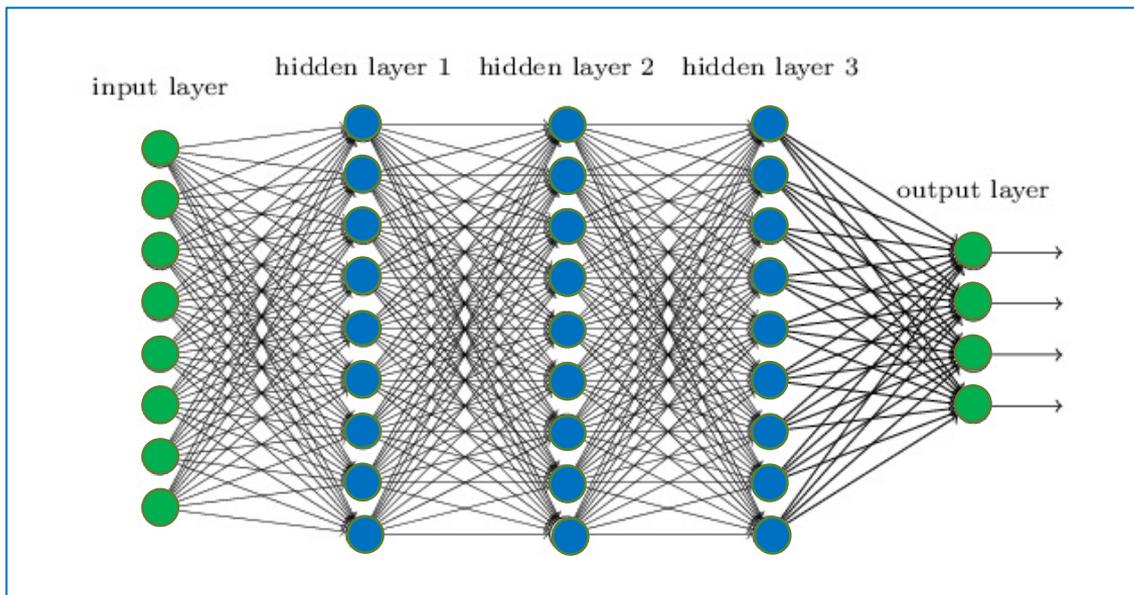


- Frequency expands to include spatial dimensions.



# Use of Machine/Deep Learning?

- **Transceivers** will use two **basic operations** (unitary Q and triangular-inverse  $G^{-1}$ ) in real time.
  1. **Filtering** – beamforming, spectrum adjustment (Q)
  2. **Recursive feedback** – (G)
- **Controller** that assigns resources (energy/information/bits), guides Q and G also.



*Controller also evolves this way – definitely adaptive optimization is very important – at edge.*

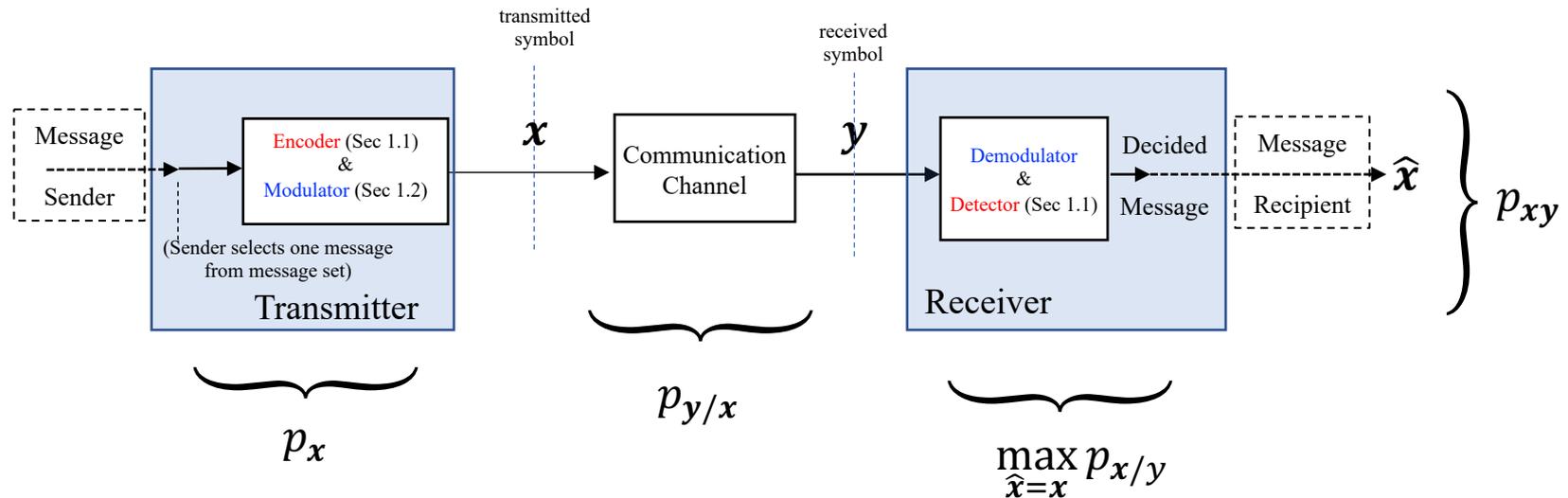


# The scalar AWGN channel

*(a foundation: Section 1.3, Section 2.1-3  
direct: 2.4.1, 2.4.3)*

*See PS1.1 (Prob 2.15 - capacity) and PS1.2 (Prob 4.3 gap)*

# Basic Communication (digital)



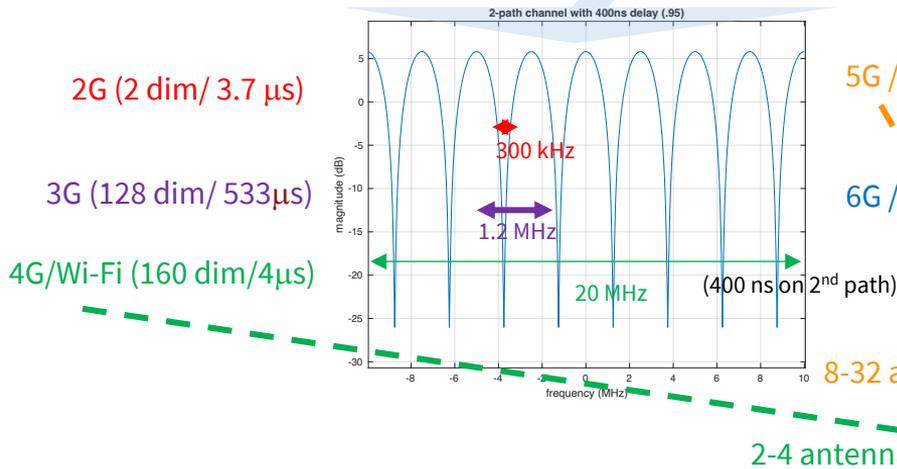
- The symbol  $x$  and messages are in some 1-to-1 relationship.
- Finding the best  $\hat{x}$  and designing  $x$  well  $\rightarrow$  this class (good 1-to-1 assumed).
- Most general channel is represented by the conditional probability  $p_{y/x}$ .
- Most general source description is  $p_x$  - together,  $p_{xy}$ .
- Optimum detector (minimizes ave error probability) is Maximum a Posteriori (MAP),  $\max p_{x/y}$ .
  - When input distribution is uniform  $\rightarrow$  ML (maximum likelihood),  $\max p_{y/x}$ .



# Communication Dimensionality

- Temporal** is Time-Frequency (at any fixed location).

- 2 x bandwidth = # of dimensions/sec (wireless or wired, including “optical” – all are EM waves)

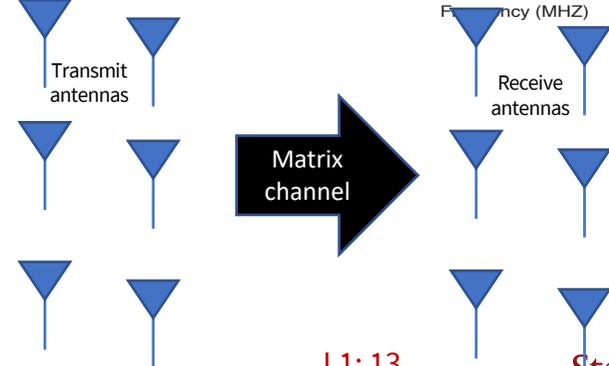
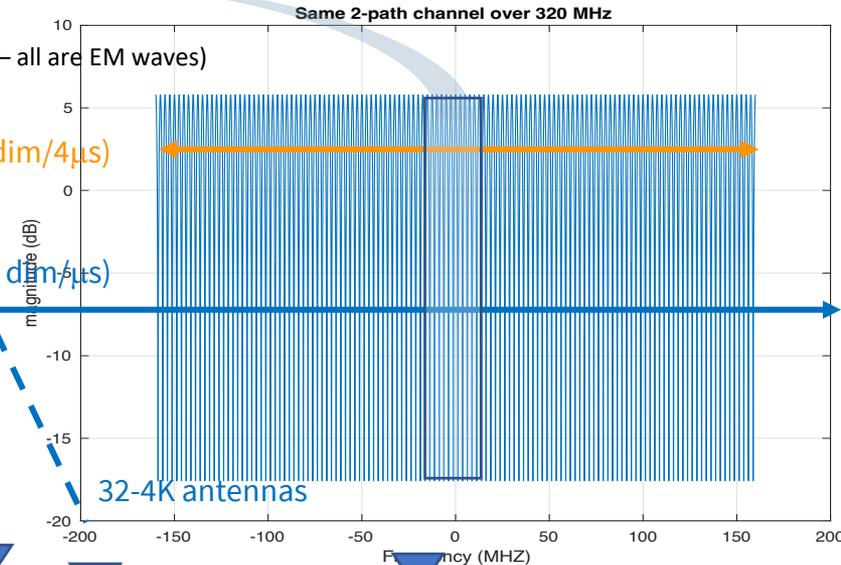


5G / Wi-Fi6 (1280 dim/4  $\mu$ s)

6G / Wi-Fi7 (~1000 dim/ $\mu$ s)

8-32 antennas

2-4 antennas



- Spatial** is space-time (at any frequency).

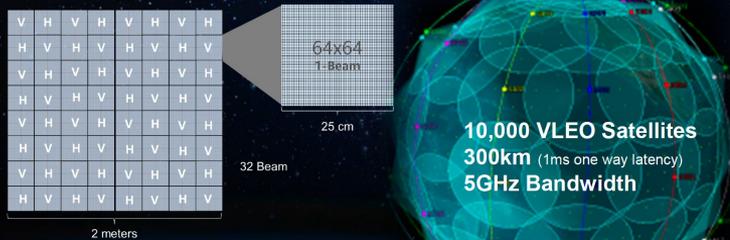
- 2D - 3D (at least ....)
  - Antenna spacing is half wavelength or more
  - 10k – 1M dimensions exist over a few microseconds.
  - Number of channels is up to # of antennas “**spatial streams**.”
  - There can be **crosstalk** between wires (e.g., ethernet, telco lines).
    - Optical xtalk is between “modes” at same frequency (a spatial effect).



# Even More Dimensions (smaller wavelengths)

## Massive VLEO Satellites

Very Low Earth Orbit (VLEO) Constellation



Courtesy Huawei

## THz Sensing and Imaging

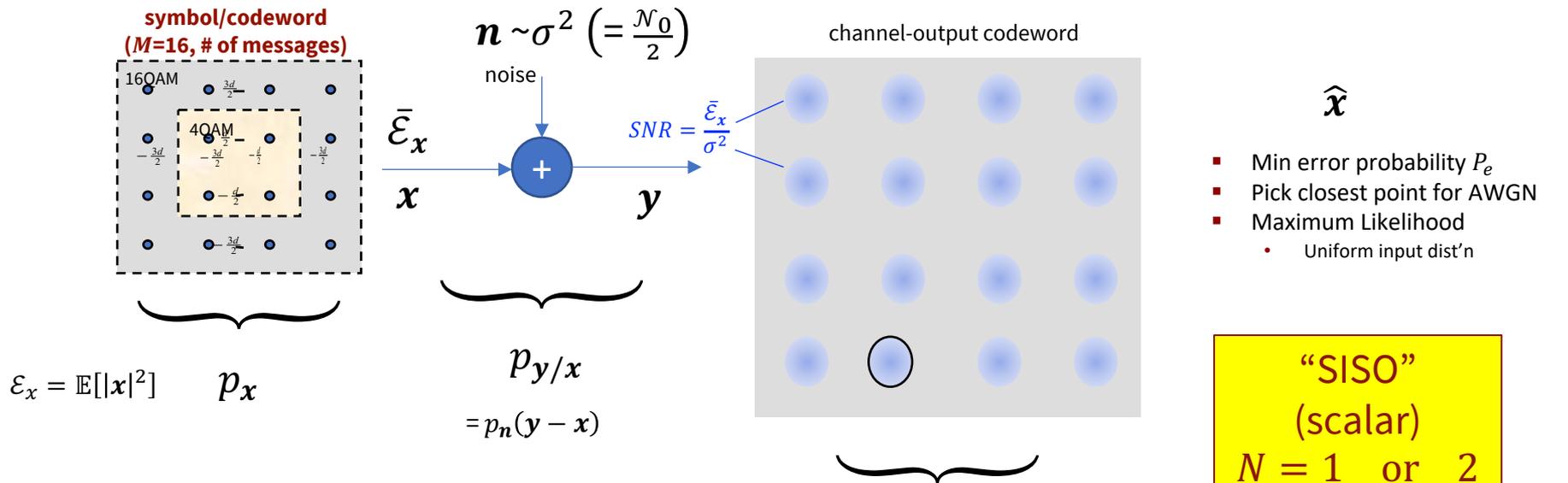
Frequency	Array Size
70GHz	1,024
140GHz	4,096
280GHz	16,348
560GHz	65,392
1.0THz	262,140
2.0THz	1,046,272
4.0THz	4,185,088

- How do we design these systems for best rates (per energy) use?
- How adaptive do they need to be?



# Simple Additive White Gaussian Noise Channel

Detection Problem First, every  $T$  seconds (symbol period)



$$\mathcal{E}_x = \mathbb{E}[|x|^2] \quad p_x$$

$$p_{y/x} = p_n(y - x)$$

$$\max_{\hat{x}=x} p_{y/x}$$

- QAM  $\rightarrow$  2 dimensional
- Uniform input (usually)  $p_x = \frac{1}{M}$
- $b = \log_2 M$  bits/symbol
- $R = \frac{b}{T}$  bits/second (data rate)

- Add noise
- Zero mean
- Variance  $\sigma^2$  (= 2-sided PSD)

$$P_e = 4 \cdot \left(1 - \frac{1}{\sqrt{M}}\right) \cdot Q\left(\sqrt{\frac{3 \cdot \text{SNR}}{M-1}}\right)$$

Subsymbol if coded  
 $x \rightarrow \tilde{x} \in \mathbb{R}$ ;  $x \rightarrow \tilde{x} \in \mathbb{C}$   
 $x$  has  $N$  real dimensions in general,  
 and has  $\bar{N}$  subsymbols, of dim  $\bar{N}$ .



# SNR, QAM, PAM reminders

$$SNR \triangleq \frac{\bar{\mathcal{E}}_x}{\sigma^2} = \frac{\text{single - sided psd}}{\text{single - sided psd}} = \frac{\text{two - sided psd}}{\text{two - sided psd}}$$

- SNR must have the same number of dimensions in numerator (signal) and denominator (noise).
- Thus, also  $SNR \triangleq \frac{\bar{\mathcal{E}}_x}{\sigma^2} = \frac{2 \cdot \bar{\mathcal{E}}_x}{N_0} = \frac{\mathcal{E}_x}{N \cdot \sigma^2}$  where  $\bar{\mathcal{E}}_x$  is energy/real-dimension.
- Energy/dimension essentially generalizes the term power/Hz (= energy) so that is why these quantities are related to power-spectral densities (psd's)
  - 1-sided  $\rightarrow$  power is integral over positive frequencies of psd.
  - 2-sided  $\rightarrow$  power is integral over all frequencies of psd.
  - These two powers are the same.
  - So -40 dBm/Hz (one-sided) psd over 1 MHz is 20 dBm, or 100 mWatts of power , [practice PS1.1 \(Prob 2.15\) and Homework Helper 1's first part.](#)
- PAM is always real baseband. QAM is always complex baseband (2 real dimensions).
  - **When QAM** has only 1 bit (2 points) in constellation, it is called BPSK (not binary PAM).
  - PAM's positive-frequency bandwidth is  $[0, 1/2T) \dots \times (1 + \alpha)$  when there is  $(100 \cdot \alpha)$  percent excess bandwidth.
  - QAM's positive-frequency bandwidth is  $[-1/2T + f_c, 1/2T + f_c) \dots$  “
  - The PAM system looks like it uses only 1/2 the bandwidth, but the QAM system is really transmitting two dimensions per symbol (so really like 2 PAM systems in parallel with symbol rate 1/T each), so no wonder it takes twice the bandwidth of a single PAM to do so.

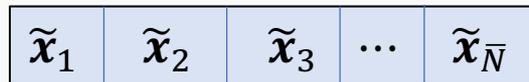


# Codes and Gaps

Shannon's maximum reliable data rate "capacity" is

$$\mathcal{C} = \log_2(1 + \text{SNR}) \text{ bits/complex-subsymbol.}$$

AWGN Max bits/sub-sym for  $P_e \rightarrow 0$  (reliably decodable)



codeword (symbol)  $x$

Good Code  $\tilde{b} \rightarrow \tilde{\mathcal{C}}$  as  $\bar{N} \rightarrow \infty$ .

subsymbols  $N = \bar{N} \cdot \tilde{N} = \# \text{ subsymbols} \times (\text{dim/subsymbol})$

$$\text{bits/dim} = \bar{b} = b/N; \text{bits/subsym} = \tilde{b} = b/\tilde{N} = \tilde{N} \cdot \bar{b}$$

Code construction

[Section 2.1.1; also PS1.1 \(Prob 2.15\)](#)

- QAM/PAM operates with given low  $P_e$  ( $10^{-6}$ ) and at a "SNR gap" ( $\Gamma = 8.8 \text{ dB} @ 10^{-6}$ ) below capacity.
  - See basics in Section 1.3.4 – [for practice, see Section 2.4; also PS1.2 \(Prob 4.3\).](#)

$$\tilde{b} = \log_2 \left( 1 + \frac{\text{SNR}}{\Gamma} \right) \text{ bits/complex-subsymbol} \leq \tilde{\mathcal{C}}.$$

$$\frac{3}{2^{\tilde{b}-1}} \cdot \text{SNR} = 13.5 \text{ dB (from } P_e = 10^{-6} \text{ formula)}$$

- For all  $\tilde{b} > 1$ , simple square QAM constellations have constant gap (= 8.8 dB at  $P_e = 10^{-6}$ ).

It's like noise increased or power decreased for  $P_e$   
(where  $\Gamma$  approaches 0 dB for best codes)  
Gap is function of code and of  $P_e$ , not  $\tilde{b}$ .



# Margin

$$\tilde{b} = \log_2 \left( 1 + \frac{SNR}{\Gamma \cdot \gamma_m} \right) \text{ bits/complex-subsymbol} \leq \tilde{C}.$$

[See also PS1.2 \(Prob 4.3\)](#)

- The designer wants “margin” protection against possible noise-power increase.
- **MARGIN**  $\gamma_m$  is this protection (usually in dB),  $\gamma_m = \frac{(SNR/\Gamma)}{2^{\tilde{b}} - 1}$ .

**Positive margin** – means performing well; **Negative margin** – means not meeting design goals.

- AWGN with SNR = 20.5 dB, then  $\tilde{C} = \log_2 (1 + 10^{2.05}) = 7$  bits/subsymbol.
- Suppose that 16-QAM ( $\tilde{b} = 4$ ) is transmitted @  $P_e = 10^{-6}$  ( $\Gamma = 8.8$  dB), then  $\gamma_m = \frac{10^{2.05 - 8.8}}{2^4 - 1} = 0$  dB.
- Suppose instead QAM with  $\tilde{b} = 5$  bits/complex-subsymbol with code of 7 dB gain ( $\Gamma \rightarrow 8.8 - 7 = 1.8$  dB).
  - $\gamma_m = \frac{10^{2.05 - 1.8}}{2^5 - 1} = 3.8$  dB.
- 6 bits/subsymbol with same code?  $\rightarrow 0.7$  dB margin – just barely below the desired  $P_e$  ;  $\bar{P}_e = P_e / N$ .

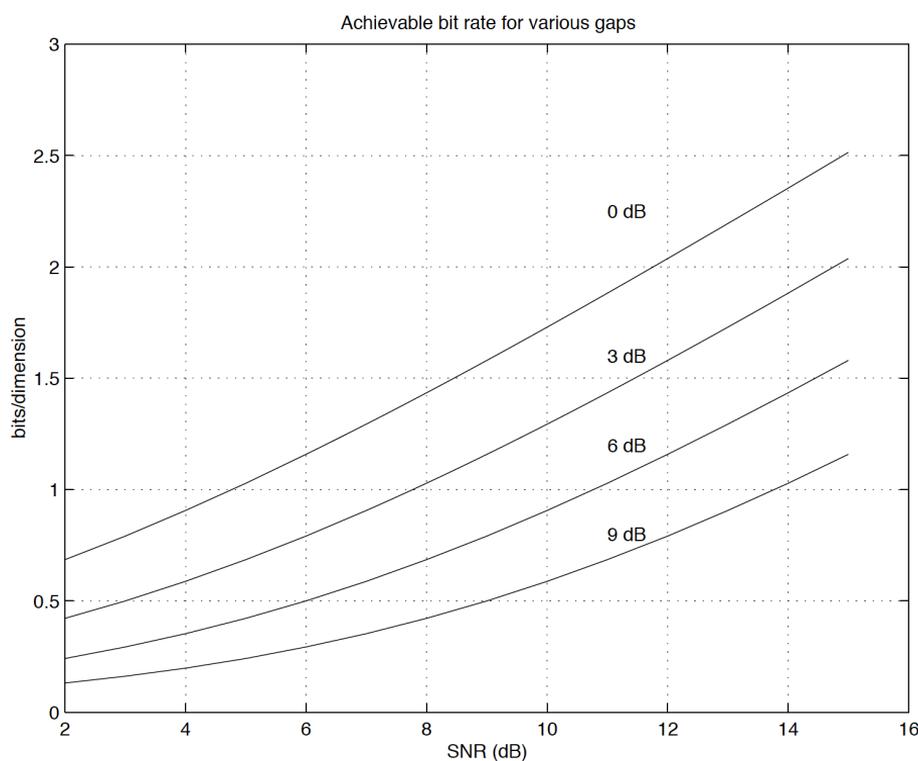


- The simple single-dimension AWGN is fundamental to most all designs.
- All subsequent designs will depend on good codes (small or 0 dB gap) re-use on those single dimension AWGNs (see EE379A).
- Designs can be optimized to get highest possible data rates for Gaussian noise:
  - single user (of course),
  - **all multiuser**,
  - channels with interference between dimensions, which includes
    - intersymbol interference (temporal),
    - crosstalk (spatial), &
    - modal (electromagnetic information theory – near field).
  - Designs are for many users with many antennas, high/low data rates, crosstalking wires, and different locations.
- The gains can be enormous (particularly with respect to EE379A coding gains).



# Gap Plot & Example

- The gap is constant, independent of the bits/dimension – greatly simplifies “loading” (adapting transmission codes to the channel).

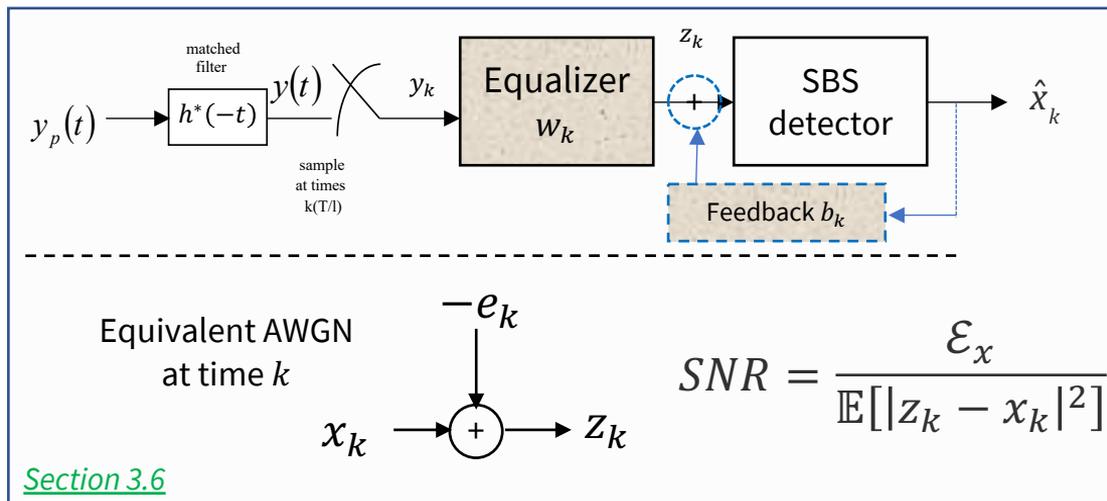
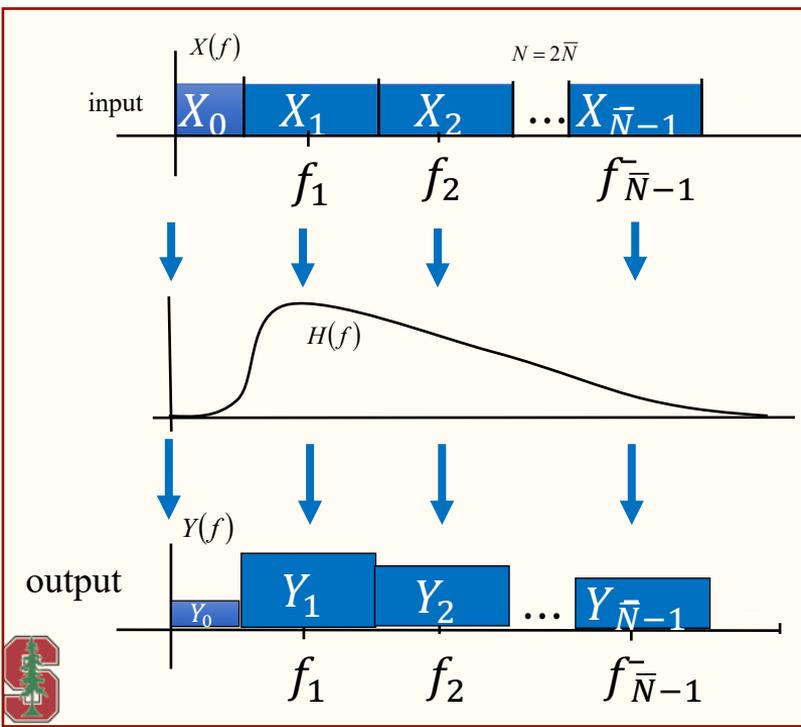


# The Matrix AWGN Channel

## *Section 2.3.5*

# Generating Parallel AWGNs

- Methods from EE379A?
- An “equalizer” is one choice and
  - creates parallel channels in time.
  - $z_k = x_k - e_k$ .

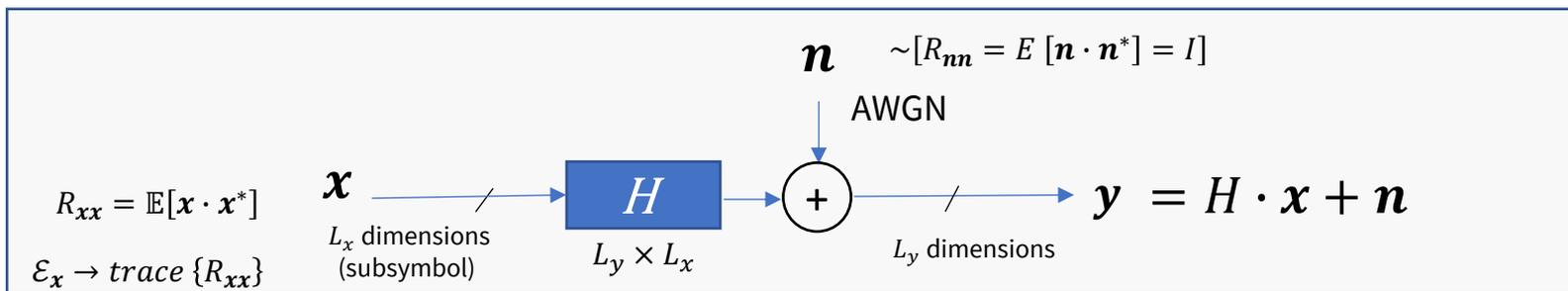


- Another?
- Multicarrier is another choice and
  - creates parallel channels in frequency.

$$Y_n \cong H_n \cdot X_n \quad (+N_n)$$

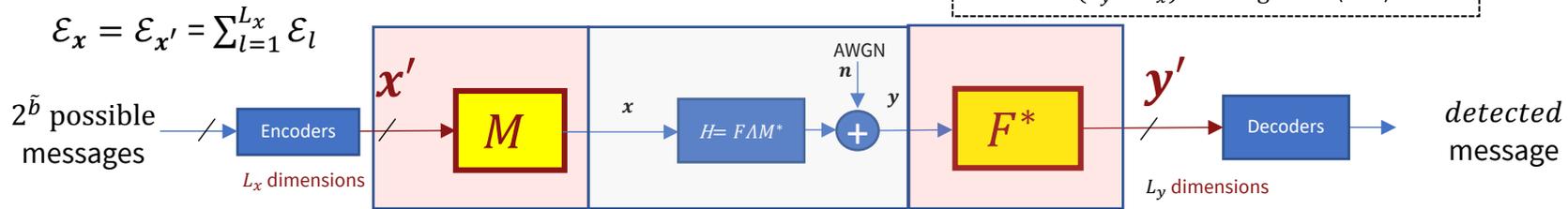
[Sections 1.3.8 and 4.2.1](#)

# In general, a matrix AWGN channel

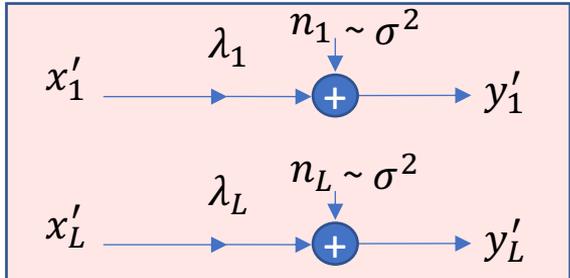


$$H = F \cdot \Lambda \cdot M^*$$

**singular value decomposition** (svd in matlab)  
 $F \cdot F^* = F^* \cdot F = I_{L_y}$ ;  $M \cdot M^* = M^* \cdot M = I_{L_x}$   
 $\Lambda$ . ( $L_y \times L_x$ ) is "diagonal" (real)



**Vector Coding (MIMO)**  
 Kasturia 1989



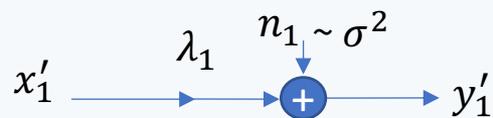
$L \leq \min(L_x, L_y)$  independent dimensions

$$R_{nn} \neq I \rightarrow (H \rightarrow R_{nn}^{-1/2} \cdot H)$$



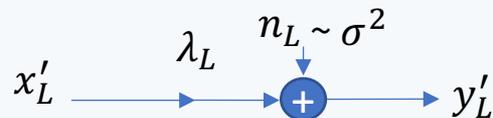
# Geometric Equivalent Channel

Parallel Independent

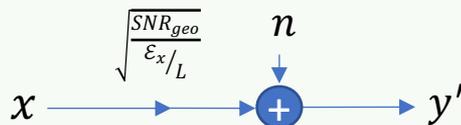


$$\tilde{b}_l = C_l = \log_2(1 + SNR_l)$$

$$SNR_l = \frac{\epsilon_l \cdot \lambda_l^2}{\sigma^2} = \epsilon_l \cdot g_l$$



$$\tilde{b} = \sum_{l=1}^L \tilde{b}_l = \sum_{l=1}^L \log_2(1 + SNR_l) = L \cdot \log_2(1 + SNR_{geo})$$



$$SNR_{geo} = \left[ \overbrace{\prod_{l=1}^L (1 + SNR_l)}^{\text{geometric average}} \right]^{1/L} - 1$$

Use it  $L$  times like single constant AWGN

- Vector Coding – uses SVD to translate matrix AWGN to set of equivalent parallel AWGN's.
  - Each can be individually encoded like AWGN (they are independent).
- Geometric-equivalent channel is used  $L$  times,
  - any  $H$  and  $R_{nn}$ , &
  - any set of input energies (that sum to allowed energy).

**We will be able in 379B to design codes for the geo-average channel.**

# The Detection/Communication Issue

- **Generally**, MAP/ML receiver/detector implementation can be very complex.
- ***Decomposing into multiple channels can simplify design!***
  - Multiple dimensions are the key to this simplification.
  - And today, used throughout digital communication (wires, wireless, soon fiber).
  - And, with proper design, there is no loss in so doing.



# The Water-Filling Energy Distribution

*Sections 2.3.5, 4.1-4.3*  
also supplementary lecture S1A

[See PS1.3 \(Prob 4.18\), PS1.4 \(Prob 4.7\), and PS1.5 \(Prob 4.25\)](#)

# Rate Maximization and Dual

- Transmitter chooses energy and bit allocation to maximize sum data rate over the dimensions  $g_l = [H_l]^2 / \sigma^2$ .

$$\max_{\bar{\epsilon}_l} \sum_{l=1}^L \frac{1}{2} \log_2 \left( 1 + \frac{\bar{\epsilon}_l \cdot g_l}{\Gamma} \right) = \sum_{l=1}^L \bar{b}_l$$

$$ST: \mathcal{E}_x = \sum_{l=1}^{L_x} \bar{\epsilon}_l$$

Rate Adaptive (RA)

$$\min_{\bar{b}_l} \sum_{l=1}^L \bar{\epsilon}_l$$

$$ST: b = \sum_{l=1}^{L_x} \tilde{b}_l$$

**DUAL**

Margin Adaptive (MA)

- Solution (basic calculus – see Section 4.2) ; see also matlab “waterfill.m” at web site to save hand calcs.

$$\bar{\epsilon}_l + \frac{\Gamma}{g_l} = \text{constant.}$$

**WATER-FILLING**

(Shannon 1948)

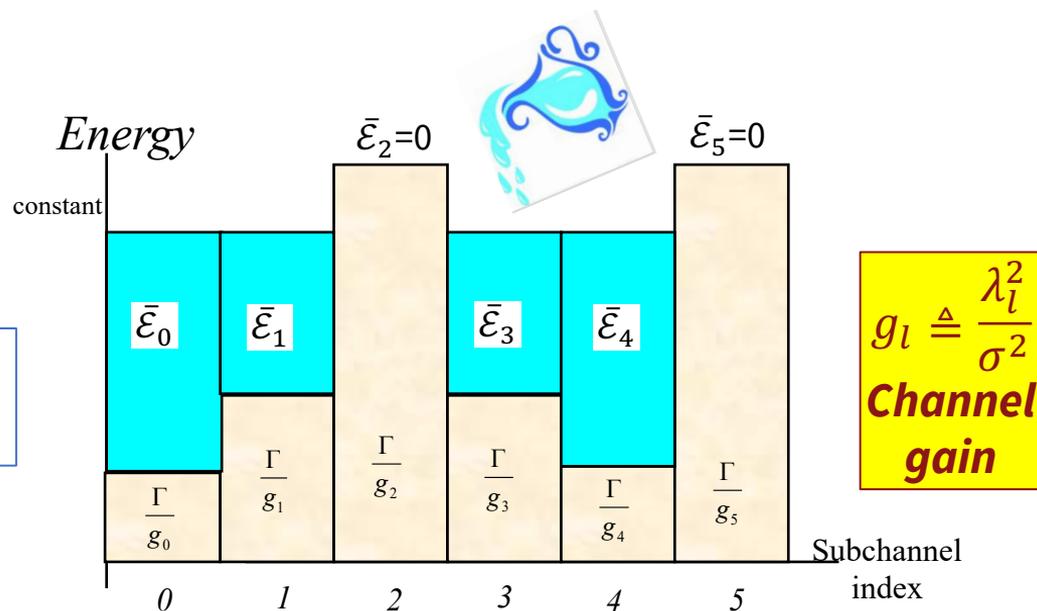
Neither energies allocated nor bits allocated can be negative.



# Water-filling Illustrated

- Energy is available in a pitcher:
  - note re-indexed 0 (DC) to 5.

RA: until all energy used.  
MA: until target bit rate attained.



$$\tilde{b}_l = \log_2 \left( 1 + \frac{SNR_l}{\Gamma} \right) \text{ where } SNR_l = \bar{\epsilon}_l \cdot g_l.$$



# Rate Adaptive Solution

$$g_1 \geq g_2 \geq \dots \geq g_L$$

- Write and sum energy constraints:

$$\varepsilon_1 + \Gamma/g_1 = K$$

$$\varepsilon_2 + \Gamma/g_2 = K$$

$$\vdots$$

$$\varepsilon_L + \Gamma/g_L = K$$

---

$$\sum_{l=1}^L \varepsilon_l + \Gamma \cdot \sum_{l=1}^L 1/g_l = L \cdot K$$

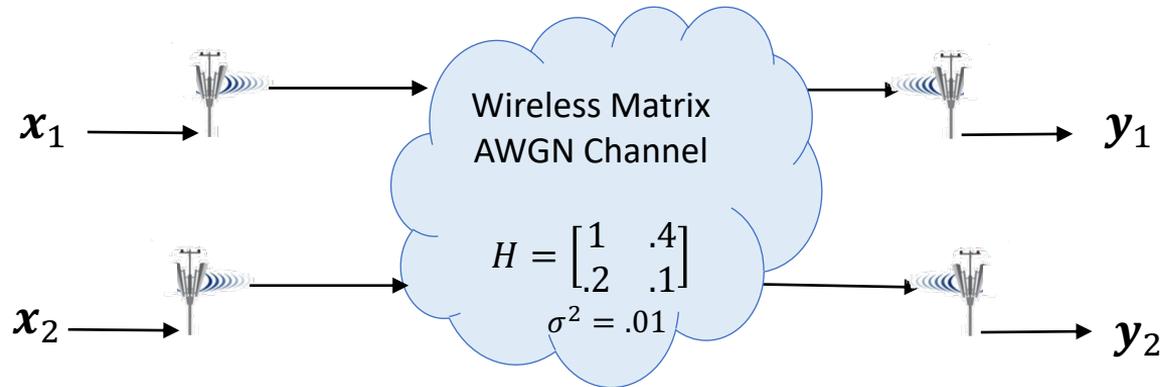
- Solve for Water-Fill Constant:

$$K = \frac{\varepsilon_x}{L^*} + \frac{\Gamma}{L^*} \cdot \sum_{l=1}^{L^*} 1/g_l$$

$L^*$  is largest  $L$  such that  $\varepsilon_l > 0$  for all  $l = 1, \dots, L^*$ .



# 2 x 2 Antenna System with 0 dB gap



- There is crosstalk between dimensions and  $\mathcal{E}_x=2$ .
  - Kind of sounds like a problem then, right?

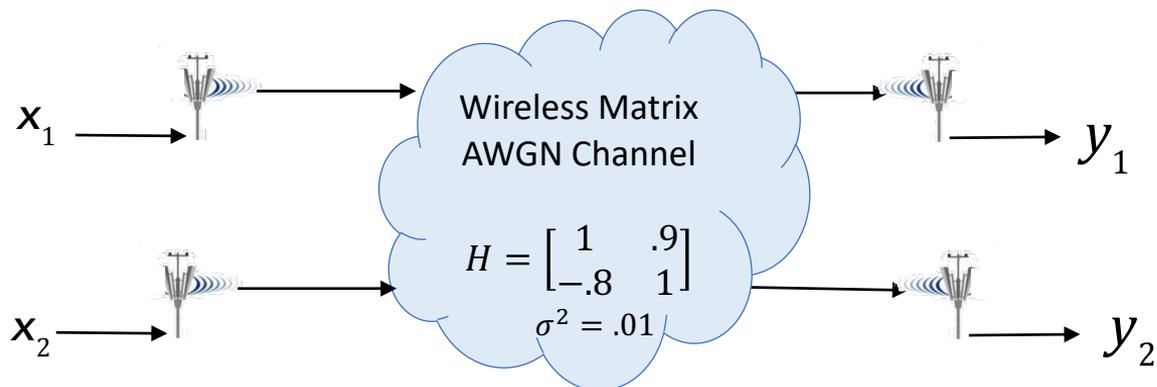
```
>> H=[10 4  
2 1];  
>> [F, Lambda, Mstar]=svd(H);  
>> Lambda =  
10.9985 0  
0 0.1818  
>> g2=Lambda(1,1)^2 = 120.9669  
>> g1=Lambda(2,2)^2 = 0.0331  
>> K=1+0.5*(1/g1+1/g2) = 16.1250  
>> E2=K-1/g2 = 16.1167  
>> E1=K-1/g1 = -14.1167 <0 (whoops)
```

Just use dimension 2  $\rightarrow \tilde{b} = \log_2(1 + 2 * g_2) = 6.93$  bits/subsymbol.

In this case water-fill simply puts all energy on the best dimension (returns to scalar/SISO if that is best).



# 2 x 2 Antenna System



- There is stronger crosstalk between dimensions.
  - Maybe worse, right? ???

```
>> H=[10 9  
-8 10];  
>> [F, Lambda, Mstar]=svd(H);  
>> Lambda =  
13.6244 0  
0 12.6244  
>> g2=Lambda(1,1)^2 = 185.6244  
>> g1=Lambda(2,2)^2 = 159.3756  
>> K=1+0.5*(1/g1+1/g2) = 1.0058  
>> E2=K-1/g2 = 1.0004  
>> E1=K-1/g1 = 0.9996  
>> btilde = log2(1+E2*g2)+log2(1+E1*g1) = 14.8693
```

Actually this is close to 2x the data rate for the previous case. Clearly, the use of both dimensions, and somewhat stronger crosstalk and signal **improves the best rate**.

In general, the increase is roughly a factor of  $L$  in data rate if  $H$  has rank  $L$ .



# Energy-minimizing Margin-Adaptive Solution

$$g_1 \geq g_2 \geq \dots \geq g_L$$

- Energy and sum-bit constraints

$$\bar{\epsilon}_l = K - \Gamma/g_l$$

$$\begin{aligned} \tilde{b} = \sum_{l=1}^L \tilde{b}_l &= \sum_{l=1}^L \log_2 \left( 1 + \frac{\bar{\epsilon}_l \cdot g_l}{\Gamma} \right) \\ &= \sum_{l=1}^L \log_2 \left( \frac{K \cdot g_l}{\Gamma} \right) \\ &= \log_2 \left( \prod_{l=1}^L \frac{K \cdot g_l}{\Gamma} \right) \end{aligned}$$

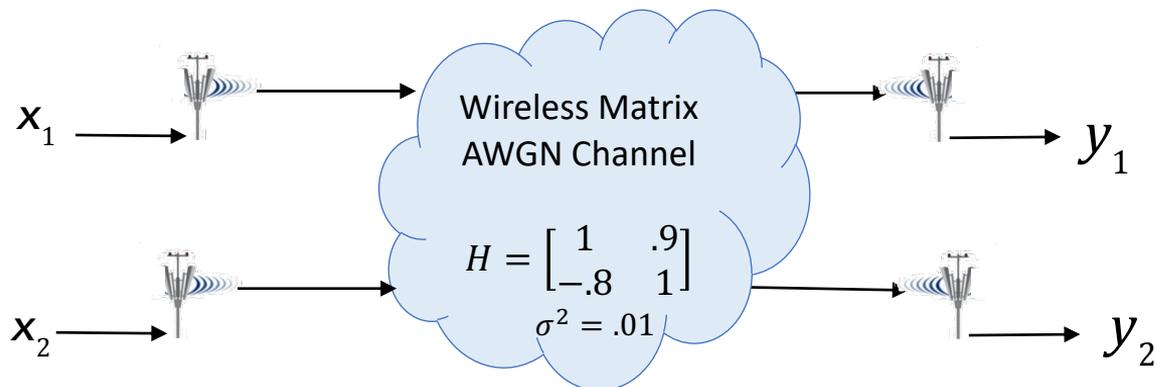
- Solve for Water-Fill Constant

$$K = \Gamma \cdot \left( \frac{2^{\tilde{b}}}{\prod_{l=1}^{L^*} g_l} \right)^{1/L^*}$$

$L^*$  is largest  $L$  such that  $\bar{\epsilon}_l > 0$  for all  $l = 1, \dots, L^*$ .



# 2 x 2 Antenna System with MA



- Attempt  $\tilde{b} = 14 \frac{\text{bits}}{\text{Hz}}$ ; The use of 2 antennas exploited channel's crosstalk,
  - Without the crosstalk, this channel supports only 7 bits/Hz (either channel has then SNR = 10).

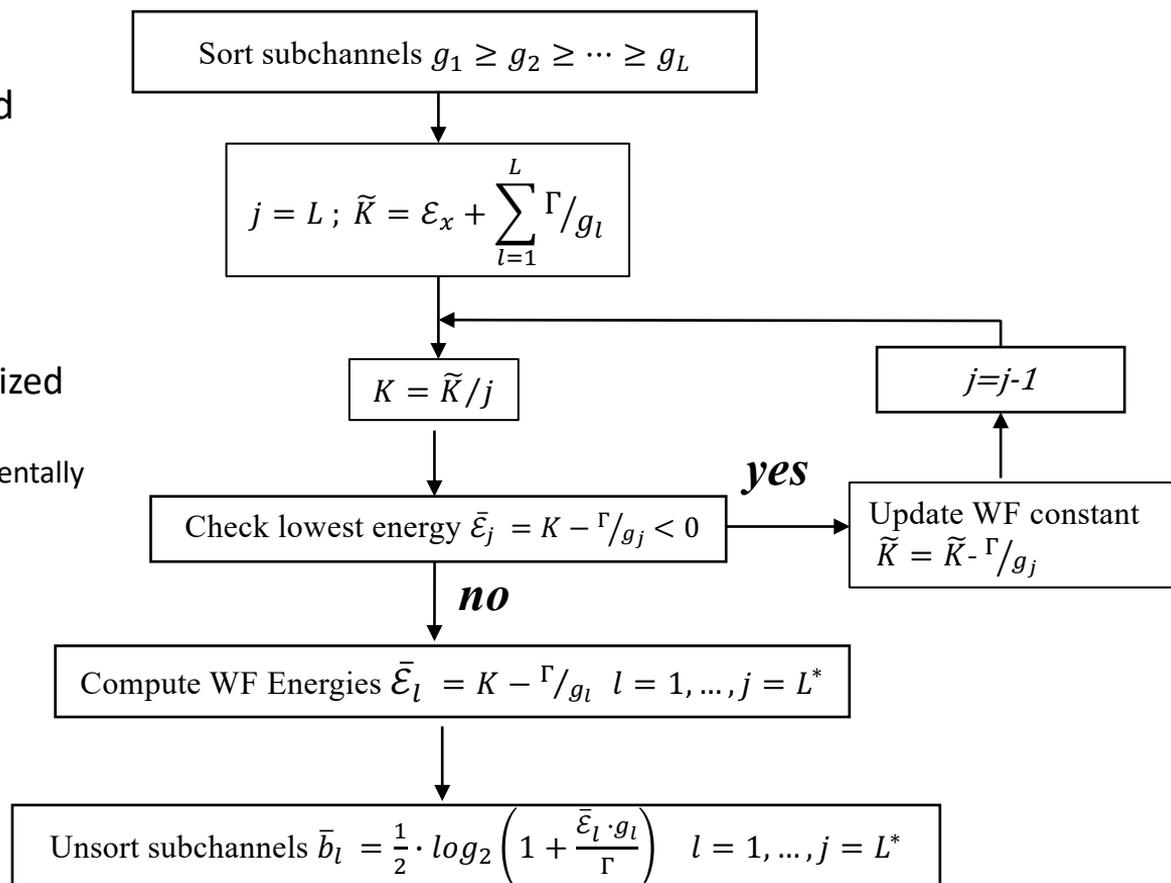
```
>> H=[10 9  
-8 10];  
>> K=sqrt((2^14)/(g1*g2)) = 0.7442  
>> E2=K-1/g2 = 0.7388  
>> E1=K-1/g1 = 0.7379  
>> margin = 10*log10(2/(E1+E2)) = 1.3 dB
```

This effect magnifies as long as most of the singular values are “decent.”

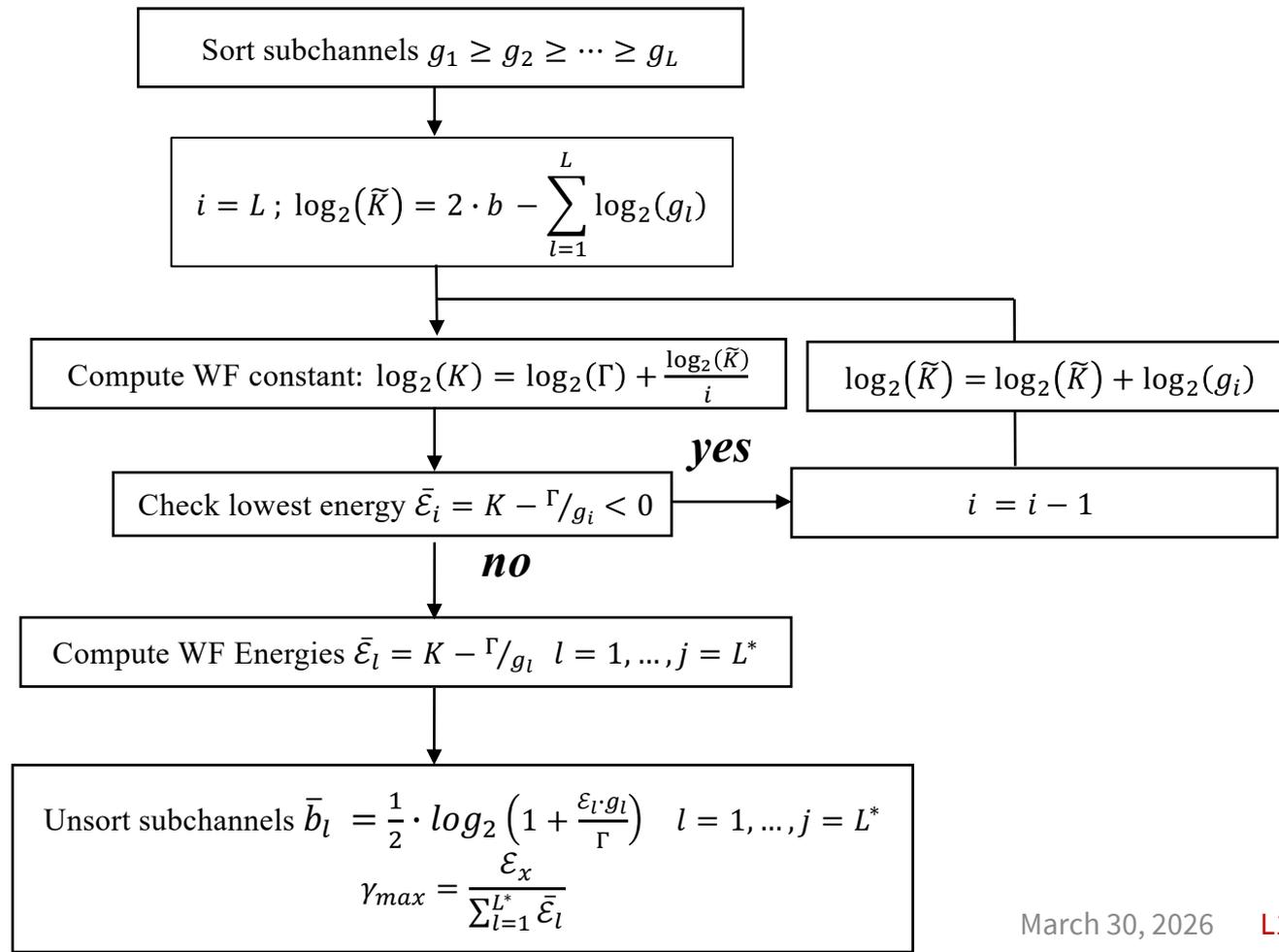


# RA Water-Fill Flow Chart

- Can start with all channels energized
  - Compute  $K$ , test lowest energy
  - Reduce number of dimensions incrementally
- Can also start with 1 channel energized
  - Compute  $K$ , test lowest energy
  - Increase number of dimensions incrementally
- The sort is most complex part
  - Can use pivots and bi-section
  - Avoids sort



# Margin Adaptive Flowchart





# End Lecture 1