

Lecture 11 BC, IC, and Other MU Channels May 9, 2024

JOHN M. CIOFFI

Hitachi Professor Emeritus of Engineering

Instructor EE379B – Spring 2024

Announcements & Agenda

- Announcements
 - PS #5 due May 14

Agenda

- Maximum BC rate sum
- Scalar Duality (BC and MAC)
- Continuous-time scalar BC
- MAC-set approach to IC
 - Examples
- IC maximum rate sum and energy sums
- Nesting, DAS, cellfree, & relay



Worst-Case Noise Examples/Uses – for algebra, see text.

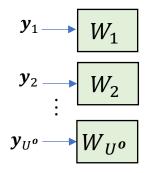
Summary of algebra appears in L10:30-32.

The (single-user) best receiver with WCN

- The MMSE receiver is block diagonal(!)
 - for WCN only, but is
 - just what the BC needs

$$W = \underbrace{S_0^{-1} \cdot G^{-*}}_{1-to-1} \cdot \underbrace{A^* \cdot H^* \cdot R_{wcn}^{-1}}_{\text{noise-white-match}} \cdot \underbrace{Q_{wcn}^* \cdot Q_{wcn}}_{I}$$
$$= S_0^{-1} \cdot G^{-*} \cdot \Phi^* \cdot Q^* \cdot Q \cdot R^* \cdot Q_{wcn}$$
$$= S_0^{-1} \cdot G^{-*} \cdot \Phi^* \cdot R^* \cdot Q_{wcn}$$
$$= S_0^{-1} \cdot G^{-*} \cdot G^{-1} \cdot D_A \cdot Q_{wcn}$$
$$= S_0^{-1} \cdot D_A \cdot Q_{wcn} ,$$

Design has same bias removal as with all MMSE.





L11: 4

BC WCN-Design Steps Summary (2.8.3.3)

Special Square Root

- Find R_{wcn} this step also finds S_{wcn} and also the primary/secondary users and $b_{max}(R_{xx})$.
 - Delete rows/columns (secondary sub user dimensions) with zeros from S_{wcn} , and correspondingly then in R_{wcn} .
- If S_{wcn} is non-trivial (block diagonal MIMO), form $S_{wcn}=Q_{wcn}^* \cdot S_{wcn}' \cdot Q_{wcn}$ (eigen decomp).
- Perform QR factorization on $Q_{wcn} \cdot R_{wcn}^{-1} \cdot H = R \cdot Q^*$ where R is upper triangular, and Q is unitary.
- Perform Cholesky Factorization on $Q^* \cdot R_{xx} \cdot Q = \Phi \cdot \Phi^*$ where Φ is also upper triangular.
- And now, the special square root is $R_{xx}^{1/2} = Q \cdot \Phi$ (see diagram L10:22 = A).

Precoder and Diagonal Receiver

- Find the diagonal matrix $D_A = \text{Diag}\{R \cdot \Phi\}$.
- Find the (primary sub-user) precoder $G = D_A^{-1} \cdot R \cdot \Phi$ (monic upper triangular).
- Find the backward MMSE (block) diagonal matrix $S_0 = D_A \cdot (S')_{wcn}^{-1} \cdot D_A$ (note, $R_b^{-1} = G^* \cdot S_0 \cdot G$).
- Block diagonal (unbiased) receiver is $W_{unb} = (S_0^{-1} I)^{-1} \cdot D_A \cdot Q_{wcn}$.
- Can check, but $b_{max}(R_{xx})$ from WCN will be $\mathbb{I}_{wcn}(x; y) = \log_2 |S_0| = \sum_{u=1}^{U^o} \log_2 (1 + SNR_{BC, wcn, u}).$

Other data rate vectors b then share this system between primary/secondary.

Section 2.8.3.5

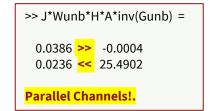
Example – all primary

• Energy \mathcal{E}_x =2 , L_x = 2

>> H = [80 70 ; 50 60]; >>Rxx=[1 .8 ; .8 1];

>> [Rwcn,b]=wcnoise(Rxx,H,1) Rwcn = 1.0000 0.0232 **Nonsingular Rwcn** 0.0232 1.0000 b = 9.6430>> Swcn = inv(Rwcn)-inv(H*Rxx*H'+Rwcn) = 0.9835 0.0000 0.0000 0.9688 >> Htilde=inv(Rwcn)*H = 78.8817 68.6440 48.1687 58.4064 >> [R,Q,P]=rq(Htilde)R = -12.4389 -74.6780 -104.56730 Q = 0.6565 -0.7544 -0.7544 -0.6565 P = 2 1**ORDER IS REVERSED SO SWITCH USERS!** J=[0 1; 1 0];

>> Rxxrot=Q'*Rxx*Q; >> Phi=lohc(Rxxrot) = 0.4482 0.0825 0 1.3388 >> DA=diag(diag(R*Phi)); >> G=inv(DA)*R*Phi = 1.0000 18.1182 1.0000 0 $>> A=Q^{inv}(R)^{*}DA^{*}G =$ 0.2942 -0.9557 -0.3381 -0.9411 >> S0=DA*inv(Swcn)*DA = 1.0e+04 * 0.0032 -0.0000 -0.0000 2.0229 Wunb=inv((S0)-eye(2))*DA*J -0.0000 -0.1822 -0.0069 -0.0000 Indeed diagonal with order switch! >> Gunb=eye(2)+S0*inv(S0-eye(2))*(G-eye(2)) = 1.0000 18.7103 1.0000 0 >> b=0.5*log2(diag(S0))' = 2.4909 7.1521 >> sum(b) = 9.6430 (checks



Try different Input Rxx, See text, Ex 2.8.7

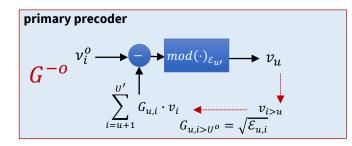


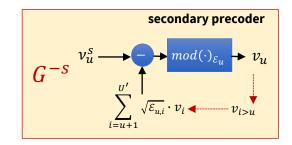
Section 2.8.3.5

May 7, 2024

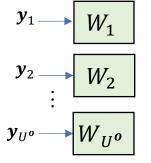
Return to Design

The design can allocate R_{xx} energy to secondary and primary users as





- The receivers are GDFE for primary users.
- Just decode with user-specific MMSE design for secondary receivers, for which corresponding users have codes with lower rates than could be decoded at the primary receivers.





Section 2.8.3.5

Another example – singular 3x3 BC (Ex 2.8.8)

>> H=[80 60 40 60 45 30 20 20 20]; >> rank(H) = 2>> Rxx=diag([3 4 2]);>> [Rwcn1, b]=wcnoise(Rxx, H, 1, 1e-5, 1e-4); >> Rwcn1 1.0000 0.7500 0.0016 0.7500 1.0000 0.0012 0.0016 0.0012 1.0000 >> b = 11.3777 >> Swcn=inv(Rwcn1)-inv(H*Rxx*H'+Rwcn1) = 0.9995 0.0000 0.0000 0.0000 -0.0000 0.0000 0.0000 0.0000 0.9948 User 2 is secondary - remove for now >> H1=[H(1,1:3)]H(3,1:3)] =80 60 40 20 20 20 >> [Rwcn, b]=wcnoise(Rxx, H1, 1, 1e-5, 1e-4); >> Rwcn = 1.0000 0.0016 0.0016 1.0000 >>b = 11.3777 >> Swcn=inv(Rwcn)-inv(H1*Rxx*H1'+Rwcn) = 0.9995 0.0000 0.0000 0.9948 Primary/Secondary



Section 2.8.3.5

May 7, 2024

>> [R,Q,P]=rq(inv(Rwcn)*H1) R = 0 9.1016 -33.2537 0 0-107.6507 O = 0.4082 -0.5306 -0.7429 -0.8165 0.1517 -0.5571 0.4082 0.8340 -0.3713 P= 2 1 ORDER IS REVERSED (Here it is order of users 1 and 3 since 2 was eliminated) >> R1=R(1:2,2:3); >> 01=0(1:3.2:3): >> Rxxrot=01' *Rxx*01 = 2.3275 0.2251 0.2251 3.1725 >> Phi=lohc(Rxxrot); >> DA=diag(diag(R1*Phi)) = 13.8379 0 0-191.7414 >> G=inv(DA)*R1*Phi = 1.0000 -4.1971 0 1.0000 >> A=Q1*inv(R1)*DA*G = -0.8067 -1.3902 0.2306 -0.9730 1.2679 -0.5559 >> A*A' = 2.5833 1.1667 -0.2500 1.1667 1.0000 0.8333 -0.2500 0.8333 1.9167

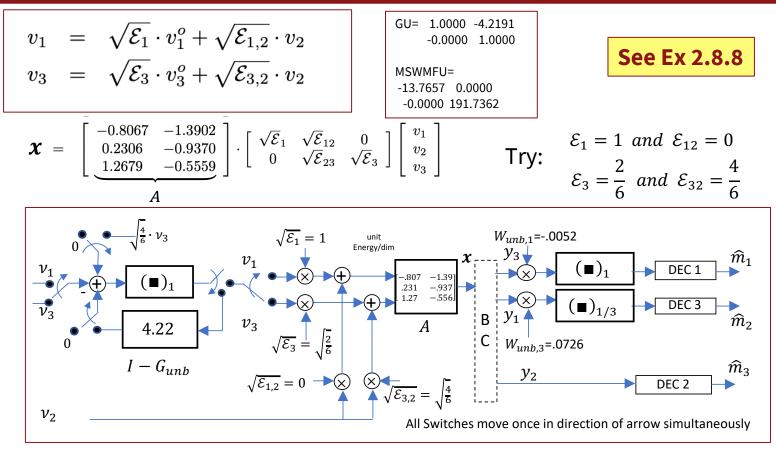
Sa Root & Precoder

>> S0=DA*inv(Swcn)*DA = 1.0e+04 * 0.0192 0.0000 Rcvr & Data Rate 0.0000 3.6957 >> MSWMFunb=inv((S0)-eye(2))*DA*J = 0.0000 0.0726 -0.0052 -0.0000 >> Gunb=eye(2)+S0*inv(S0-eye(2))*(G-eye(2)) = 1.0000 -4.2191 -0.0000 1.0000 >> b=0.5*log2(diag(S0))' = 3.7909 7.5868 >> sum(b) = 11.3777 checks >> H*A = 0.0219 - 191.8333 $\sqrt{\frac{2}{3}} + \frac{1}{2}$ 0.0164 -143.8749 13.8379 -58.3825 See Example 2.8.8 or details of below Assign 1 energy unit to User 1, 1/3 to user 3, and now squeeze in 2/3 energy on user 2 >> b=0.5*log2(diag([11/3]) *diag(S0)) = 3.7909 6.7943 Crosstalk is >> ct=1/3*143.9^2 = 6.8928e+03 >> b2=0.5*log2(1+(2/3)*60^2/6892.8) = 0.2155 >> b2+sum(b) = 10.8007 < 11.3777 Energy on secondary reduces rate sum

Not equal to Rxx Energy not inserted into null space (same on part that is in pass space)

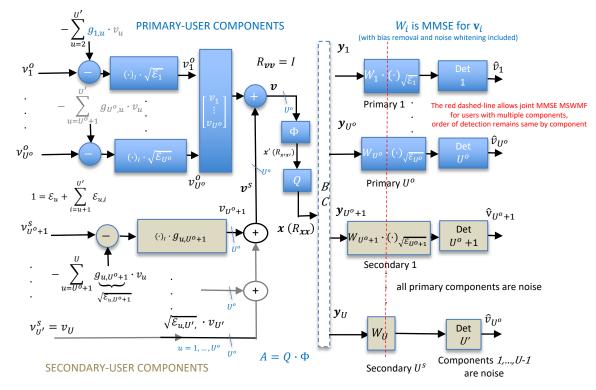
L11: 8

System Diagram for this WCN design





Gaussian Vector BC System Diagram



This design is for any R_{xx} , but the square root $Q \cdot \Phi$ is very special and unique; this design is for the R_{wcn} , no matter the real correlation between receiver noises; U' is number of user components.



Section 2.8.3.4 May

May 7, 2024

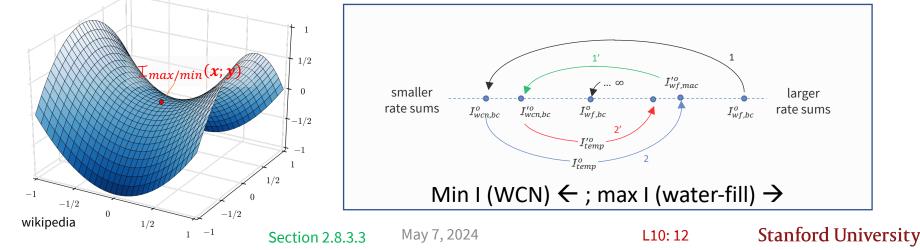
L10: 10

Maximum BC rate sum

Maximum BC rate sum

- Maximize I(x; y) through water-filling (but ... presumes receivers can coordinate).
 - This is concave problem that always can be solved for the best input autocorrelation R_{xx} .
- Minimize $I_{min}(x; y)$ through worst-case-noise to get $I_{wcn}(x; y)$.
 - This is a convex problem that always can be solved for worst noise autocorrelation R_{wcn} .

This is a saddle-point problem that produces a max-min = min-max:



bcmax.m

function [Rxx, Rwcn, bmax] = bcmax(iRxx, H, Lyu) Uses cvx_wcnoise.m and rate-adaptive waterfill.m (Lagrange Multiplier based) Inputs: - iRxx: initial input autocorrelation array, size is Lx x Lx x N. Only the sum of traces matters, so can initialize to any valid autocorrelation matrix Rxx to run wcnoise. needs to include factor N/(N+nu) if nu ~= 0 - H: channel response, size is Ly x Lx x N, w/o sqrt(N) normalization - Lyu: array number of antennas at each user; scalar Lyu means same for all **Outputs:** - Rxx: optimized input autocorrelation, Lx x Lx x N - Rwcn: optimized worst-case noise autocorrelation, with white local noise Ly x Ly x N so IF H is noise-whitened for Rnn, then actual noise is Rwcn^(1/2)*Rnn*Rwcn^(*/2) - b: maximum sum rate/real-dimension - user must mult by 2 for complex case

Revisit example from slide L11:8

iRxx =
3 0 0
0 4 0
0 0 2
>>H=
80 60 40
60 45 30
20 20 20
[RxxA, RwcnA, bmax] = bcmax(iRxx, H, 1)
RxxA =
3.7515 1.5032 -0.7451
1.5032 1.5019 1.5007
-0.7451 1.5007 3.7465
RwcnA =
1.0000 0.7500 0.0008
0.7500 1.0000 0.0006
0.0008 0.0006 1.0000
bmax = 12.1084 (> 11.3777 that occurred earlier)

Secondary components' energy is zeroed for this maximum rate sum.



Revisit example from L11:6

```
H = [ 80 70
50 60];
>> iRxx=[1.8
.81];
>> [Rxx, Rwcn, bmax] = bcmax(iRxx, H, Lyu)
Rxx =
1.0001 0.0082
0.0082 0.9999
Rwcn =
1.0000 0.0049
0.0049 1.0000
bmax = 10.3517 > 9.6430
```

Usually converges pretty quickly, not always though – CVX can get finicky when singularity involved.



New Example – Singular Rwcnopt

>> H 1.0719 -0.8627 -0.1901 0.2952	inv(Rwcnopt) - inv(H*Rxxopt*H'+Rwcnopt) -69.9931 62.4186 -26.7687 -6.3512	>> [V,D]=eig(Rxxopt)
1.0498 -0.7245 0.2568 0.2757	62.4186 -54.9656 23.8385 5.6560	V =
-0.4586 0.5595 1.0027 0.0530 0.4107 0.0496 0.3965 -0.7740	-26.7687 23.8385 -9.5873 -2.4256 -6.3512 5.6560 -2.4256 -0.1050	-0.5342 -0.7457 -0.3635 -0.1626
[F,L,M]=svd(H);	[Rwcnopt, sumRatebar, S1, S2, S3, S4] =	-0.7873 0.6052 -0.0344 -0.1124 0.2071 0.2556 -0.9213 0.2073
L =	cvx_wcnoise(Rxxopt, H, [1 1 1 1])	-0.2278 -0.1108 0.1336 0.9581
2.0671 0 0 0	Rwcnopt =	
0 1.1449 0 0	1.0000 0.9163 -0.4792 -0.0191	D =
0 0 0.8130 0	0.9163 1.0000 -0.0991 0.1251	D -
0 0 0.0000	-0.4792 -0.0991 1.0000 0.1044	<mark>0.0000</mark> 0 0 0
H=F(:,1)*M(:,1)'*L(1,1)+F(:,2)*M(:,2)'*L(2,2)+F(:,3)*M(:,3)'*L(3,3);	-0.0191 0.1251 0.1044 1.0000	$\begin{array}{cccccccccccccccccccccccccccccccccccc$
	sumRatebar = 2.1911	0 0 0 0.9256
[Rxxopt, Rwcnopt, bmax] = bcmax(eye(4), H, 1)	rank(H) = 3	
	>> \$3+\$4(1:4,1:4) =	
Rxxopt =	<mark>0.0978</mark> 000	Confirms that best input is also
1.1559 -0.7381 0.0995 -0.0691	0 0.6204 0 0	singular - it should never have higher
-0.7381 0.6399 0.2862 -0.2207	0 0 0.6360 0	rank than number of primary user
0.0995 0.2862 1.3091 -0.0326		(components).
-0.0691 -0.2207 -0.0326 0.8951	rank(H) = 3	
Rwcnopt =	>> Htilde=pinv(Rwcnopt)*H;	
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	>> [R,Q,P]=rq(Htilde); R =	
-0.4792 -0.0991 1.0000 0.1044	10000000.0211 - 0.2595 0.1213	
-0.0191 0.1251 0.1044 1.0000	0 -0.9411 -0.0435 0.0071	
bmax = 2.1911	0 0 -1.0542 0.0496	
>> det(Rwcnopt) % = 1.5337e-09 Singular Rwcn	0 0 0 -1.1499	
	P = 1 4 2 3	
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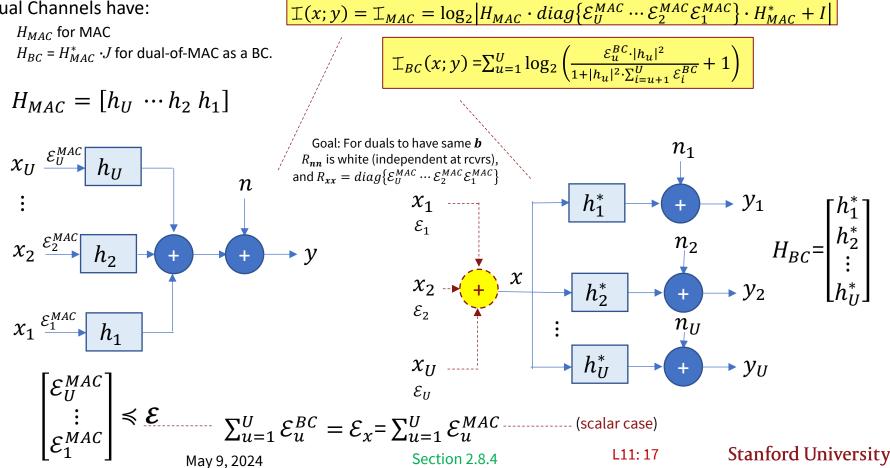
Scalar Duality (BC and MAC)

PS5.2 - 2.29 scalar BC region

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Scalar Dual Channels – Same I(x; y)

- Dual Channels have:
 - H_{MAC} for MAC
 - $H_{BC} = H^*_{MAC} \cdot J$ for dual-of-MAC as a BC.



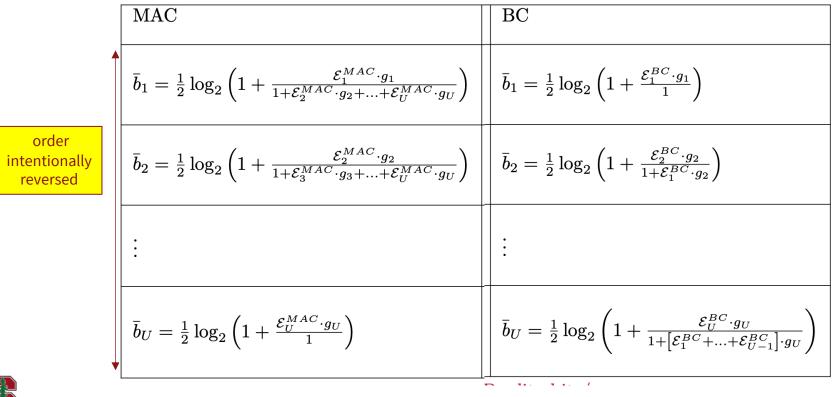


X

Scalar Duality

• Set data rates equal and solve for $\mathcal{E}_u^{MAC/BC}$:

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Section 2.8.4

L11: 18

Corresponding Energies

$$\begin{split} \mathcal{E}_{1}^{BC} &= \mathcal{E}_{1}^{MAC} \cdot \frac{1}{1 + \mathcal{E}_{2}^{MAC} \cdot g_{2} + \ldots + \mathcal{E}_{U}^{MAC} \cdot g_{U}} \\ \mathcal{E}_{2}^{BC} &= \mathcal{E}_{2}^{MAC} \cdot \frac{1 + \mathcal{E}_{1}^{BC} \cdot g_{2}}{1 + \mathcal{E}_{3}^{MAC} \cdot g_{3} + \ldots + \mathcal{E}_{U}^{MAC} \cdot g_{U}} \\ \vdots &= \vdots \\ \mathcal{E}_{U}^{BC} &= \mathcal{E}_{U}^{MAC} \cdot \left(1 + \left[\mathcal{E}_{1}^{BC} + \ldots + \mathcal{E}_{U-1}^{BC}\right] \cdot g_{U}\right) \end{split}$$

- By selecting these energies, all user rates are the same (with the order reversal) and running through all such energies that sum to total produces the SAME energy-sum capacity region.
- See proof in text (Theorem 2.8.2 in Section 2.8.4).



Section 2.8.4

L11: 19

Revisit Scalar Example

- Total energy is 1, instead use dual MAC to investigate BC with:
 - $\mathcal{E}_2^{BC} = 0.25$ (bottom of BC),
 - $\mathcal{E}_1^{BC} = 0.75$ (top BC), &
 - reversing order $g_1 = 6400$ and $g_2 = 2500$.

$$\mathcal{E}_{2}^{MAC} = \frac{\mathcal{E}_{2}^{BC}}{1 + \mathcal{E}_{1}^{BC} \cdot g_{2}} = \frac{.25}{1 + 2500 \cdot (.75)} = \frac{1}{7504} = 1.3326 \times 10^{-4} \text{ (top MAC)}$$
$$\mathcal{E}_{1}^{MAC} = \mathcal{E}_{1}^{BC} \cdot \left(1 + g_{2} \cdot \mathcal{E}_{2}^{MAC}\right) = .75 \cdot (1 + 2500/7504) = \frac{7503}{7504} = .9999 = 1 - \mathcal{E}_{2}^{MAC} \text{ (bottom MAC)}$$

User data rates for this combination are (and were in earlier table found directly for BC).

$$b_{1} = \frac{1}{2} \cdot \log_{2} \left(1 + \frac{\mathcal{E}_{1}^{MAC} \cdot g_{1}}{1 + \mathcal{E}_{2}^{MAC} \cdot g_{2}} \right) = 6.1144$$
$$b_{2} = \frac{1}{2} \cdot \log_{2} \left(1 + \frac{\mathcal{E}_{2}^{MAC} \cdot g_{2}}{1} \right) = .2074$$

• Can use the easier MAC developments to analyze the BC through duality.



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Section 2.8.4

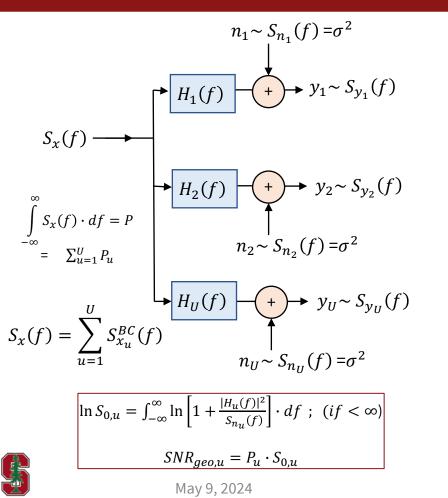
L11: 20

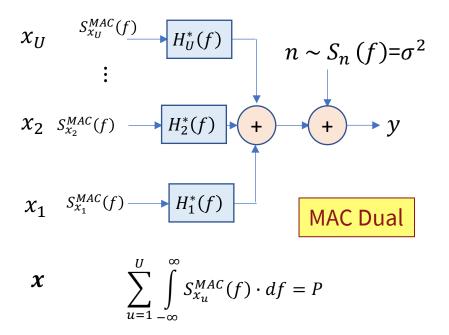
Continuous-time Scalar BC

Section 2.8.5

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Continuous time/freq Scalar BC





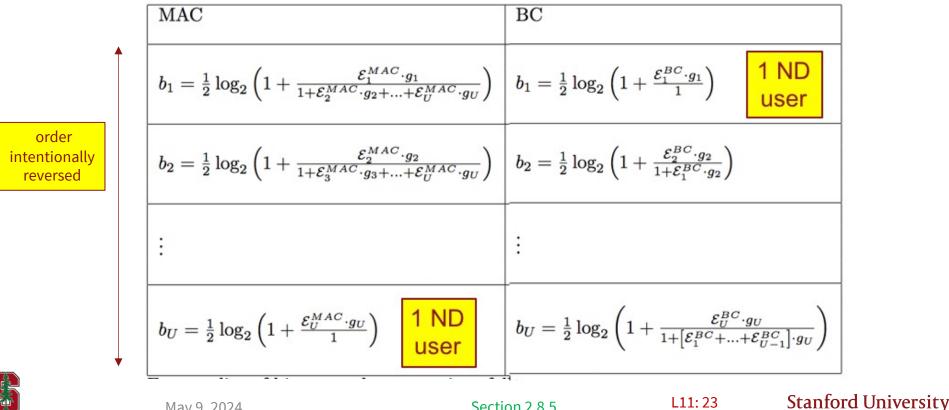
Design for this MAC, and then find dual

Section 2.8.5

Scalar Duality

Replace with integrals and $\mathcal{E}_{u}^{MAC/BC} \rightarrow S_{u}^{MAC/Bc}(f)$

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Section 2.8.5

L11:23



Corresponding PSD's

•
$$\mathcal{E}_u^{MAC/BC} \to S_u^{MAC/BC}(f)$$

• See proof in notes (Theorem 2.8.2 in Section 2.8.4), but execute with PSD's $S_u^{MAC/BC}(f)$

$$\begin{array}{lcl} \mathcal{E}_{1}^{BC} & = & \mathcal{E}_{1}^{MAC} \cdot \frac{1}{1 + \mathcal{E}_{2}^{MAC} \cdot g_{2} + \ldots + \mathcal{E}_{U}^{MAC} \cdot g_{U}} \\ \mathcal{E}_{2}^{BC} & = & \mathcal{E}_{2}^{MAC} \cdot \frac{1 + \mathcal{E}_{2}^{BC} \cdot g_{2}}{1 + \mathcal{E}_{3}^{MAC} \cdot g_{3} + \ldots + \mathcal{E}_{U}^{MAC} \cdot g_{U}} \\ \vdots & = & \vdots \\ \mathcal{E}_{U}^{BC} & = & \mathcal{E}_{U}^{MAC} \cdot \left(1 + \left[\mathcal{E}_{1}^{BC} + \ldots + \mathcal{E}_{U-1}^{BC}\right] \cdot g_{U}\right) \end{array}$$



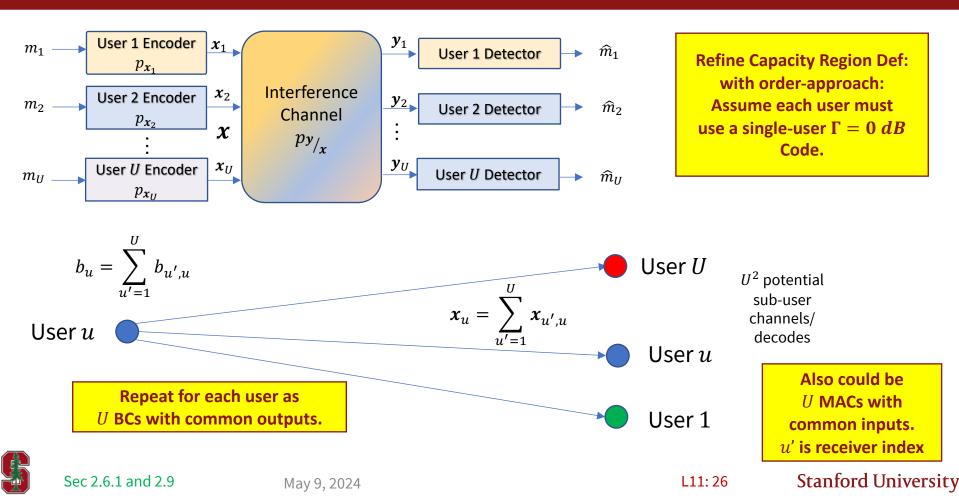
L11: 24

MAC-set Approach to IC

Sec 2.9

May 9, 2024

The Interference Channel (IC)

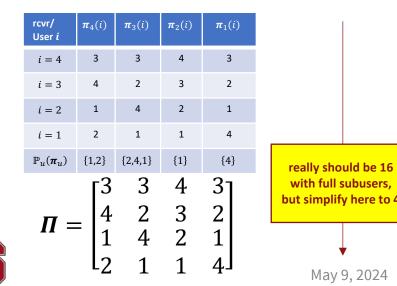


Prior-User Set (repeat from L7)

Order vector and its inverse are:

$$\boldsymbol{\pi}_{u} = \begin{bmatrix} \pi(U') \\ \vdots \\ \pi(1) \end{bmatrix} \quad \boldsymbol{\pi}_{u}^{-1} = \begin{bmatrix} U' \\ \vdots \\ 1 \end{bmatrix} \qquad j = \pi(i) \rightarrow i = \pi^{-1} e^{ij}$$

- Prior-User Set is $\mathbb{P}_u(\pi) = \{j \mid \pi^{-1}(j) < \pi^{-1}(u)\}$
 - That is "all the users before the desired user u in the given order π .
 - Receiver *u* best decodes these "prior" users and removes them, while "post" users are noise
 - π can be any order in $\mathbb{P}_u(\pi)$, but the most interesting is usually π_u (receiver u's order)



 Data rates (mutual information bounds) average only those users who are not cancelled are xtalk noise.

	I	I 4	I 3	I 2	\Im_1
_	top	00	I ₃ (3/1,2,4) 20	8	00
		⊥ ₄ (4/1,2) 10	⊥ ₃ (2/1,4) 9	8	Ø
4		⊥ ₄ (1/2) 5	⊥ ₃ (4/1) 4	⊥ ₂ (2/1) 4	⊥ ₁ (1/4) 2
	bottom	⊥ ₄ (2) 1	I(1) 2	⊥ ₂ (1) 2	⊥ ₁ (4) 5
	Sectior	1 2.6.2		L11:	27

$$\mathbb{I}_{min}(\boldsymbol{\Pi}, \boldsymbol{p}_{\boldsymbol{x}\boldsymbol{y}}) = \begin{bmatrix} 4\\20\\1\\2 \end{bmatrix}$$

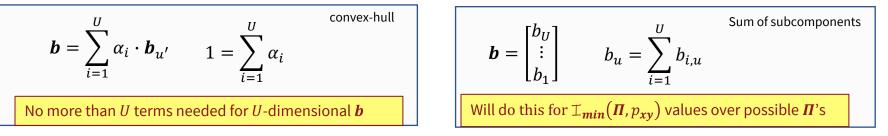
Maximum number of subusers U²

User u has maximum bit rate, when all other users are given (cancelled):

$$b_u \leq \mathbb{I}(\boldsymbol{x}_u; \boldsymbol{y}_u / \boldsymbol{x}_{\boldsymbol{U} \setminus u})$$

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- Although -- this may not be good for the other users $i \neq u$.
 - $I_{min}(\Pi, p_{xy})$ calculation, given a Π , precedes a subsequent convex-hull search over all Π to obtain the achievable region $\mathcal{A}(\boldsymbol{b}, p_{xy})$



- For U ≥ 2, the other users' subuser components may be desirable to decode, but not all → U'! ≤ (U²)! for each receiver's order may need search/evaluation.
- Π maximally has $(U'!)^U$ possible choices (in most general case).

At any receiver u , the subuser components separate into two groups for any given order π_u :

(1) those cancelled (or generally conditional probability has specific given values for those components), and

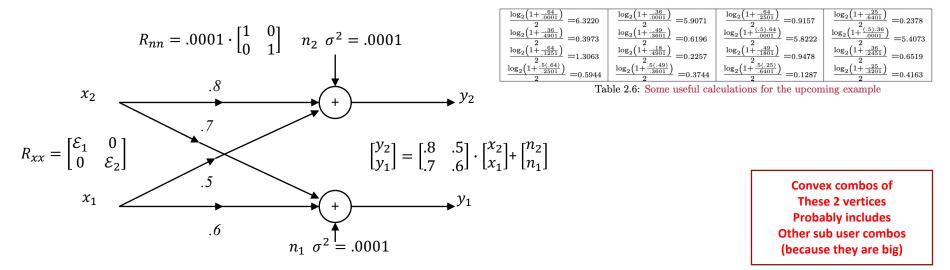
(2) those not cancelled, which are averaged out generally in marginal distributions.

If that choice is made for each user for each receiver, there are thus maximally U choices across all receivers into (1) or (2), so $U' \leq U^2$.



L11: 28

Example channel – Scalar Gaussian IC



- Earlier H, but this time as an IC
 - .8>.7 and .6 > .5

Sec 2.9.1

- Not complete set of orders (4 instead of (4!)² = 576)
- Shaded points are interior to line formed by unshaded points

Ι	Ι	u	\mathcal{E}_{u}	$\mathbb{P}_u(oldsymbol{\pi}_u)$	$\mathscr{I}_u(x_u;y_u/\mathbb{P}_u(oldsymbol{\pi}_u))$	$\mathscr{I}_{\neg u}(x_u; y_{\neg u}/\mathbb{P}_{\neg u}({\pmb{\pi}}_{\neg u}))$	$\mathcal{I}_{min,u}(oldsymbol{\Pi}, oldsymbol{\mathcal{E}})$
$\lceil 2 \rceil$	2	2	1	{1}	6.322	∞	6.322
$\lfloor 1$	$1 \rfloor$	1	1	Ø	.3973	.2378	.2378
[1	1]	2	1	Ø	.9157	.6196	.6196
$\lfloor 2$	$2 \rfloor$	1	1	$\{2\}$	5.9071	∞	5.9071
[1	$2\rceil$	2	1	Ø	.9157	∞	.9157
$\lfloor 2$	$1 \rfloor$	1	1	Ø	.3973	∞	.3973
$\lceil 2 \rceil$	1]	2	1	{1}	6.322	.2378	.2378
[1	$2 \rfloor$	1	1	$\{2\}$	5.9071	.6196	.6196
Table 2.7: Evaluation of \mathcal{I}_{min} for different orders.							



L11: 29

IC Rate Region Examples

Sec 2.9.1

May 9, 2024

Example continued

Γ	I	u	\mathcal{E}_{u}	$\mathbb{P}_u({m \pi}_u)$	$\mathscr{I}_u(x_u;y_u/\mathbb{P}_u(oldsymbol{\pi}_u)$	$\mathscr{I}_{u}(x_{\neg u};y_{\neg u}/\mathbb{P}_{\neg u}(\boldsymbol{\pi}_{\neg u}))$	$\mathcal{I}_{min,u}(oldsymbol{\Pi}, oldsymbol{\mathcal{E}})$
$\lceil 2 \rceil$	2	2	1	{1}	6.322	∞	6.322
[1	1	1	0.5	Ø	.2257	.1287	.1287
[1	1]	2	1	Ø	1.3063	.9478	.9478
$\lfloor 2$	2	1	0.5	$\{2\}$	5.4073	∞	5.4073
$\lceil 2 \rceil$	2	2	0.5	{1}	5.822	∞	5.822
L1	1	1	1	Ø	.6519	.4163	.4163
[1	1]	2	0.5	Ø	5.822	.3744	.3744
$\lfloor 2$	$2 \rfloor$	1	1	{1}	5.9071	∞	5.9071
2	2	2	1	$\{1,2\}$	6.322	∞	6.322
[1	1	1	0.95	$\{1\}$.3818	.2276	.2276
[1	1]	2	1	$\{2\}$.9425	.6412	.6412
$\lfloor 2$	$2 \rfloor$	1	0.95	$\{1,2\}$	5.8701	∞	5.8701

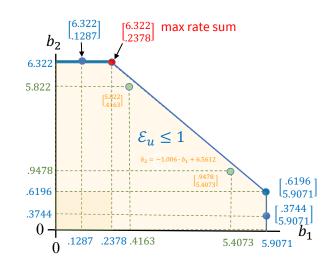


Table 2.8: More example points with the best orders

- The dimension-sharing of the large- b_u points dominates the other points on the interior.
- Also check of vertices' derivatives relative to the dimension-sharing line (-1.006) try upper.

$$\ln(2) \cdot \frac{db_2}{d\mathcal{E}_2} = \frac{3200}{6400 \cdot \mathcal{E}_2 + 1} = 0.4999 \qquad \qquad \frac{db_2}{db_1} = -3.56$$

$$\ln(2) \cdot \frac{db_1}{d\mathcal{E}_2} = -\frac{3200 \cdot 2500 \cdot \mathcal{E}_1}{6400 \cdot (\mathcal{E}_2 + 2500 \cdot \mathcal{E}_1 + 1) \cdot (6400 \cdot \mathcal{E}_2 + 1)} = -.1404 \qquad \frac{db_2}{db_1} = -3.56$$

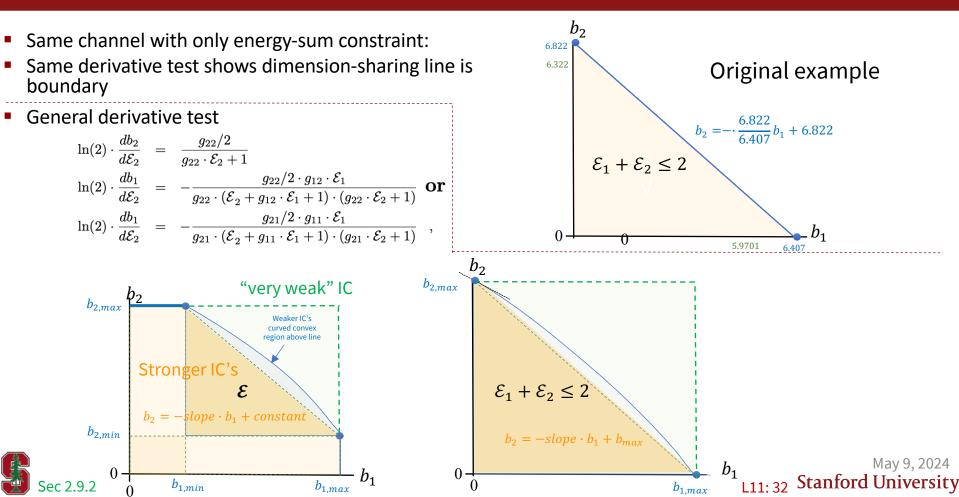
Vertices all () inside pentagon for this example

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- For 2 vertices if magnitude of slope is less than 1, then upper point and otherwise lower point (or whole line).
 - Could check other vertex also, but if curvature is within the line already (convex), then no need. Sec 2.9.1 May 9, 2024

L11: 31

Energy-sum IC extension



So-called "weak" symmetric IC

Achievable Region when $\mathcal{E}_1 = \mathcal{E}_2 = 1$

 $\boldsymbol{b} = \begin{vmatrix} \frac{1}{2} \\ 1 \end{vmatrix}$

 $\boldsymbol{b} = \begin{bmatrix} \frac{1}{2} \cdot \log_2\left(1 + \frac{c_2}{1 + \alpha^2 \cdot \varepsilon_1}\right) \\ \frac{1}{2} \cdot \log_2\left(1 + \frac{\varepsilon_1}{1 + \alpha^2 \cdot \varepsilon_2}\right) \end{bmatrix}$

 $\boldsymbol{\Pi} = \begin{bmatrix} 2 & 2 \\ 1 & 1 \end{bmatrix}$

- $\begin{bmatrix} y_2 \\ y_1 \end{bmatrix} = \begin{bmatrix} 1 & \alpha \\ \alpha & 1 \end{bmatrix} \cdot \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} + \begin{bmatrix} n_2 \\ n_1 \end{bmatrix}$ $R_{nn} = I \text{ and } \mathcal{E} \leq \mathbf{1}$
- When $\alpha \to 0$, there is no crosstalk and so $C_{IC}(\mathbf{b})$ is a square.
- When $\alpha > 1$, $C_{IC}(\mathbf{b})$ is a pentagon.

Vector Gaussian IC Example

• 2 users and H is 4 x 2
$$y = \begin{bmatrix} y_2 \\ y_1 \end{bmatrix} = \begin{bmatrix} H_2 \\ H_1 \end{bmatrix} \cdot \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} + \begin{bmatrix} n_2 \\ n_1 \end{bmatrix}$$

$$R_{nn}$$
=.01 · l

$$H_1 = \begin{bmatrix} \mathbf{h}_{12} & \mathbf{h}_{11} \end{bmatrix} = \begin{bmatrix} .8 & .7 \\ .6 & .5 \end{bmatrix}$$

 $H_2 = [h_{22} \ h_{21}] = \begin{bmatrix} .9 & .3 \\ .2 & .2 \end{bmatrix}$

>> H2 = [9 3 3 8]; >> Rb2inv=H2'*H2+diag([1 1]); >> Gbar2=chol(Rb2inv); >> G2=inv(diag(diag(Gbar2)))*Gbar2; >> S02=diag(diag(Gbar2))*diag(diag(Gbar2)); >> 0.5*log2(diag(S02)) = b2 = 3.2539b1 = 2.7526>> H1 = [8 7 5]; 6 >> Rb1inv=H1'*H1+diag([1 1]); >> Gbar1=chol(Rb1inv); >> S01=diag(diag(Gbar1))*diag(diag(Gbar1)); >> 0.5*log2(diag(S01)) = b2 = 3.3291b1 = 0.4128

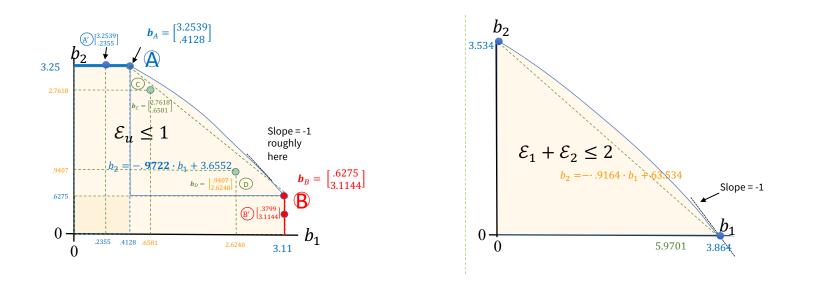
```
>> J2=hankel([0 1]);
>> Rb2inv=J2*H2'*H2*J2+diag([1 1]);
  74 51
  51 91
>> Gbar2=chol(Rb2inv);
>> S02=diag(diag(Gbar2))*diag(diag(Gbar2));
>> 0.5*log2(diag(S02)) =
b1 = 3.1047
b2 = 2.9018
>> Rb1inv=J2*H1'*H1*J2+diag([1 1]);
>> Gbar1=chol(Rb1inv);
>> S01=diag(diag(Gbar1))*diag(diag(Gbar1));
>> 0.5*log2(diag(S01)) =
b1 = 3.1144
b2 = 0.6275
```

Try various energy points With the same orders



Sec 2.9.2

4x2 IC Example continued



Vary energies near max points to see if local points above or below dimension-sharing, so curved boundary or flat

• Since the line has slope magnitude less than 1, then the curvature is above this line with max rate sum at magnitude 1

PS 5.4 (2.31) – IC channel has mix, one 2x2 user and one scalar user

May 9, 2024

Sec 2.9.2

L11: 35

Chris AF's MU_IC.m

function [b, GU, WU, S0, MSWMFU] = mu_ic(H, A, Lxu, Lyu, cb) Per-tonal (temporal dimension) multiuser interference channel receiver and per-user bits - Chris Adelico Ferrarin - 2023	 IC needs Lxu and Lyu A is like mu_mac Per tone Arrange input Hu matrices according to order Π 	_
 Inputs: H, A, Lxy, Lyu, cb Outputs: b, GU, WU, SO, MSWMFU Definitions: H: noise-whitened channel matrix [HUU ··· HU1] sum-Lyu x sum-Lxu : `.`! H2U ··· H21 [H1U ··· H11] A: Block Diag sq-root sum-Lxu x sum-Lxu discrete modulators, blkdiag([AU A1]); The Au entries derive from each IC user's Lxu x Lxu input autocorrelation matrix, where the trace of each such autocorrelation matrix is user u's energy/symbol. This is per-tone. Lxu: # of input dimensions for each user U 1 in 1 x U row vector Lyu: # of output dimensions for each user U 1 in 1 x U row vector cb: = 1 if complex baseband or 2 if real baseband channel GU: unbiased feedback matrix sum-Lxu x sum-Lxu x U with matrices indexed from user U (e.g. GU(:,:,2) gives GU for user U-1). S0: sub-channel channel gains sum-Lxu x sum-Lxu x U with matrices indexed from user U (e.g. GU(:,:,2) gives SO for user U-1). S0: sub-channel channel gains sum-Lxu x sum-Lxu x U with matrices indexed from user U (e.g. GU(:,:,2) gives SO for user U-1). MSWMFU: unbiased mean-squared whitened matched filter, sum-Lxu x Ly x U with matrices indexed from user U (e.g. MSWMFU(:,:,2) gives MSWMFU for user U-1). b. user u's bits/symbol 1 x U the user should recompute b if there is a cyclic prefix 	$ \begin{array}{c} H2 = [0.9 \ 0.3; \ 0.3 \ 0.8]; \mbox{\% From L11:34} \\ H1 = [0.8 \ 0.7; \ 0.6 \ 0.5]; \\ sigma2 = 0.01; \\ Ht = [H2; H1] / \ sqrt(sigma2); \ \mbox{\% this is } 4x2 \ matrix (2 \ outputs/input) \\ A = [1 \ 0; \ 0 \ 1]; \\ Lxu = [1 \ 1]; \\ Lyu = [2 \ 2]; \\ cb = 2; \\ >> [b_A, GU_A, WU_A, SO_A, MSWMFU_A] = mu_ic(Ht, A, Lxu, Lyu, cb); \\ USER 2 \qquad USER 1 \\ >> b_A \ \mbox{\% = } \\ \hline \begin{array}{c} 3.2539 \qquad 0.4128 \\ GU_A(:;:,1) = \qquad GU_A(:;:,2) = \\ 1.0000 0.5667 \qquad 1.0000 0.8600 \\ 0 1.0000 \qquad 0 1.0000 \\ WU_A(:,:,1) = \qquad WU_A(:,:,2) = \\ 0.0111 0 \qquad 0.0100 0 \\ -0.0126 0.0225 \qquad -1.1026 1.2949 \\ SO_A(::,1) = \qquad SO_A(::,2) = \\ 91.0000 0 101.0000 0 \\ 0 45.4176 \qquad 0 1.7723 \\ MSWMFU_A(:,:,1) = \qquad MSWMFU_A(:,:,2) = \\ 0.1000 0.0333 \qquad 0.0800 0.0600 \\ -0.0460 0.1423 \qquad 0.2436 -0.1410 \end{array} \right)$	×
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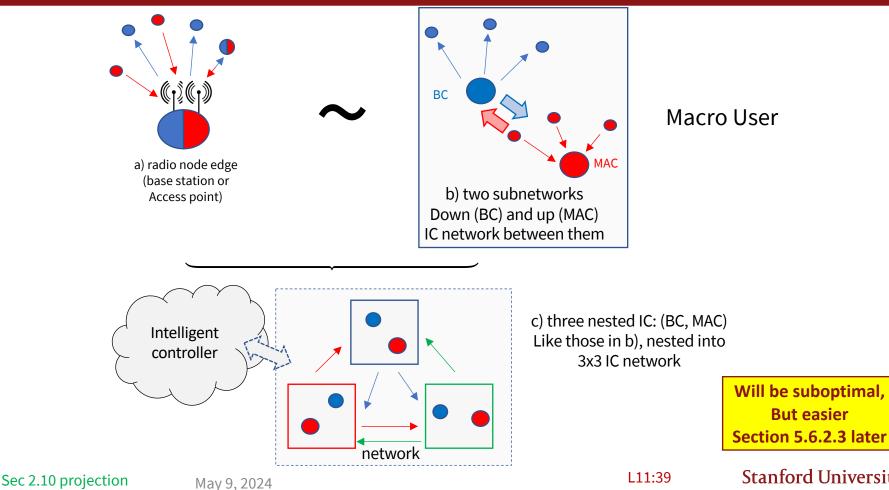




End Lecture 11

Nesting, DAS, cellfree & Relay

Multiuser Nesting



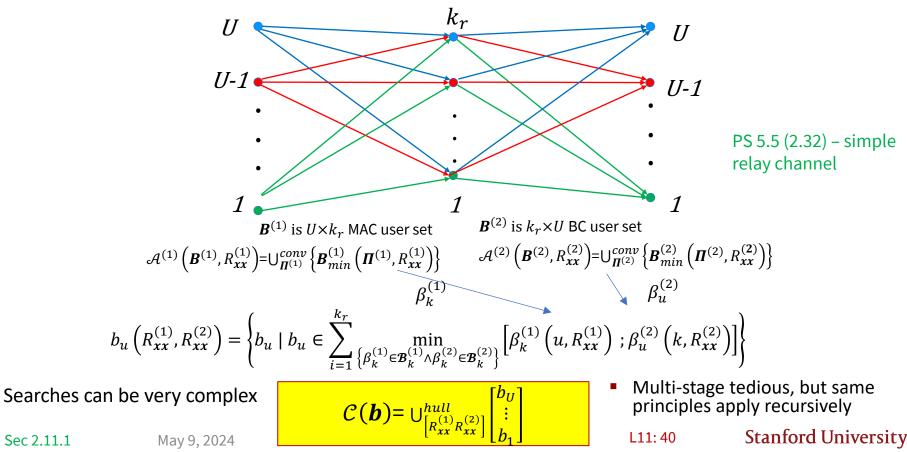
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But easier

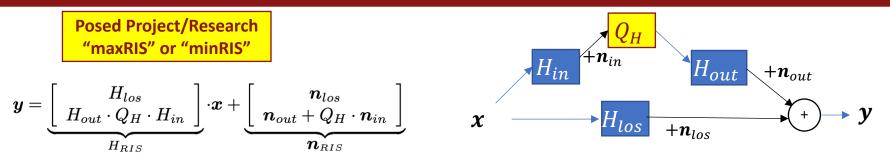
Single-Stage Relay Channel

Conceptually uses what we know already and introduces sub-users at k_r relay points

Sec 2.11.1



Reflective Intelligent Surfaces (RIS)



- The RIS matrix Q_H satisfies $||Q_H||_F^2 \leq G_H$, the RIS gain it may also satisfy
 - *Q_H* is unitary matrix (preserves energy)
 - Q_H is diagonal, and usually also unitary, to be phase/gain-only adjustment on each antenna port (in-to-out)
 - *Q_H* has individual elements restricted
- For a given R_{xx} , maximize over Q_H $\mathcal{I}(\boldsymbol{y}; \boldsymbol{x}) = \log_2 |R_{n,RIS} + H_{RIS} \cdot R_{\boldsymbol{x}\boldsymbol{x}} \cdot H_{RIS}^*|$
- For a given Q_H , maximize the same over R_{xx}
 - Iterate? \rightarrow not convex in the Q_H , this needs work.

Sec 2.11.4 May

May 9, 2024

$$R_{\boldsymbol{n}\boldsymbol{n},RIS} = \left[egin{array}{cc} R_{\boldsymbol{n}\boldsymbol{n}} & 0 \ 0 & R_{\boldsymbol{n}\boldsymbol{n},out} + Q_H \cdot R_{\boldsymbol{n}\boldsymbol{n}in} \cdot Q_H^* \end{array}
ight]$$

L11: 41 Stanford University