

#### *Lecture 11* **BC, IC, and Other MU Channels** *May 9, 2024*

#### **J OHN M. C IOFFI**

Hitachi Professor Emeritus of Engineering

Instructor EE379B – Spring 2024

### **Announcements & Agenda**

- Announcements
	- PS #5 due May 14

#### ■ Agenda

- Maximum BC rate sum
- Scalar Duality (BC and MAC)
- Continuous-time scalar BC
- MAC-set approach to IC
	- Examples
- IC maximum rate sum and energy sums
- Nesting, DAS, cellfree, & relay



# **Worst-Case Noise Examples/Uses – for algebra, see text.**

**Summary of algebra appears in L10:30-32.**

# **The (single-user) best receiver with WCN**

- The MMSE receiver is block diagonal(!)
	- for WCN only, but is
	- just what the BC needs

$$
W = \underbrace{S_0^{-1} \cdot G^{-*}}_{1 - to - 1} \cdot \underbrace{A^* \cdot H^* \cdot R_{wcn}^{-1}}_{\text{noise-white-match}} \cdot \underbrace{Q_{wcn}^* \cdot Q_{wcn}}_{I}
$$
\n
$$
= S_0^{-1} \cdot G^{-*} \cdot \Phi^* \cdot Q^* \cdot Q \cdot R^* \cdot Q_{wcn}
$$
\n
$$
= S_0^{-1} \cdot G^{-*} \cdot \Phi^* \cdot R^* \cdot Q_{wcn}
$$
\n
$$
= S_0^{-1} \cdot G^{-*} \cdot G^{-1} \cdot D_A \cdot Q_{wcn}
$$
\n
$$
= S_0^{-1} \cdot D_A \cdot Q_{wcn},
$$

§ Design has same bias removal as with all MMSE.





# **BC WCN-Design Steps Summary (2.8.3.3)**

#### **Special Square Root**

- Find  $R_{wcn}$  this step also finds  $S_{wcn}$  and also the primary/secondary users and  $b_{max}(R_{xx})$ .
	- Delete rows/columns (secondary sub user dimensions) with zeros from  $S_{wcn}$ , and correspondingly then in  $R_{wcn}$ .
- If  $\mathcal{S}_{wcn}$  is non-trivial (block diagonal MIMO), form  $\mathcal{S}_{wcn}$ = $Q^*_{wcn} \cdot \mathcal{S}'_{wcn} \cdot Q_{wcn}$  (eigen decomp).
- Perform QR factorization on  $Q_{wcn} \cdot R_{wcn}^{-1} \cdot H = R \cdot Q^*$  where R is upper triangular, and Q is unitary.
- **•** Perform Cholesky Factorization on  $Q^* \cdot R_{rr} \cdot Q = \Phi \cdot \Phi^*$  where  $\Phi$  is also upper triangular.
- And now, the special square root is  $R_{xx}^{1/2} = Q \cdot \Phi$  (see diagram L10:22 = A).

#### **Precoder and Diagonal Receiver**

- Find the diagonal matrix  $D_4 = \text{Diag} \{ R \cdot \Phi \}.$
- Find the (primary sub-user) precoder  $G = D_A^{-1} \cdot R \cdot \Phi$  (monic upper triangular).
- Find the backward MMSE (block) diagonal matrix  $S_0 = D_A \cdot (S')_{wcn}^{-1} \cdot D_A$  (note,  $R_b^{-1} = G^* \cdot S_0 \cdot G$ ).
- Block diagonal (unbiased) receiver is  $W_{unb} = (S_0^{-1} I)^{-1} \cdot D_A \cdot Q_{wcn}$ .
- Can check, but  $b_{max}(R_{xx})$  from WCN will be  $\mathcal{I}_{wen}(x; y) = \log_2|S_0| = \sum_{u=1}^{U^o} \log_2(1 + SNR_{BC,wen,u}).$

#### **Other data rate vectors b then share this system between primary/secondary.**

# **Example – all primary**

Energy  $\mathcal{E}_r = 2$ ,  $L_r = 2$ 

 $\gg H = [80 70 ; 50 60];$ >>Rxx=[1 .8 ; .8 1];

>> [Rwcn,b]=wcnoise(Rxx,H,1)  $Rwcn =$ 1.0000 0.0232 0.0232 1.0000  $b = 9.6430$  $\Rightarrow$  Swcn = inv(Rwcn)-inv(H\*Rxx\*H'+Rwcn) = 0.9835 0.0000 0.0000 0.9688 >> Htilde=inv(Rwcn)\*H = 78.8817 68.6440 48.1687 58.4064 >> [R,Q,P]=rq(Htilde)  $R =$  -12.4389 -74.6780 0 -104.5673  $Q =$  0.6565 -0.7544 -0.7544 -0.6565  $P = 2 1$ **ORDER IS REVERSED SO SWITCH USERS!** J=[0 1; 1 0]; **Nonsingular Rwcn** >> Rxxrot=Q'\*Rxx\*Q; >> Phi=lohc(Rxxrot) = 0.4482 0.0825 0 1.3388 >> DA=diag(diag(R\*Phi));  $\gg$  G=inv(DA)\*R\*Phi = 1.0000 18.1182 0 1.0000  $\Rightarrow$  A=Q\*inv(R)\*DA\*G = 0.2942 -0.9557 -0.3381 -0.9411 >> S0=DA\*inv(Swcn)\*DA = 1.0e+04 \* 0.0032 -0.0000 -0.0000 2.0229 Wunb=inv((S0)-eye(2))\*DA\***J** -0.0000 -0.1822 -0.0069 -0.0000 Indeed diagonal with order switch!  $\Rightarrow$  Gunb=eye(2)+S0\*inv(S0-eye(2))\*(G-eye(2)) = 1.0000 18.7103 0 1.0000  $\Rightarrow$  b=0.5\*log2(diag(S0))' = 2.4909 7.1521  $\gg$  sum(b) = 9.6430 (checks)



Try different Input Rxx, See text, Ex 2.8.7



# **Return to Design**

**•** The design can allocate  $R_{xx}$  energy to secondary and primary users as





- The receivers are GDFE for primary users.
- Just decode with user-specific MMSE design for secondary receivers, for which corresponding users have codes with lower rates than could be decoded at the primary receivers.

$$
y_1 \longrightarrow W_1
$$
  

$$
y_2 \longrightarrow W_2
$$
  

$$
\vdots
$$
  

$$
y_{U^o} \longrightarrow W_{U^o}
$$

 $\mathbf{v}$ 



### **Another example – singular 3x3 BC (Ex 2.8.8)**

>> H=[80 60 40 60 45 30 20 20 20];  $\gg$  rank(H) = 2  $\gg$  Rxx=diag([3 4 2]); >> [Rwcn1, b]=wcnoise(Rxx, H, 1, 1e-5 , 1e-4); >> Rwcn1 1.0000 0.7500 0.0016 0.7500 1.0000 0.0012 0.0016 0.0012 1.0000  $\Rightarrow$  b = 11.3777 >> Swcn=inv(Rwcn1)-inv(H\*Rxx\*H'+Rwcn1) = 0.9995 0.0000 0.0000 0.0000 -0.0000 0.0000 0.0000 0.0000 0.9948 User 2 is secondary – remove for now  $\Rightarrow$  H1=[H(1,1:3)  $H(3,1:3)$ ] = 80 60 40 20 20 20 >> [Rwcn, b]=wcnoise(Rxx, H1, 1, 1e-5 , 1e-4);  $\geq$  Rwcn = 1.0000 0.0016 0.0016 1.0000  $>> h = 11.3777$ >> Swcn=inv(Rwcn)-inv(H1\*Rxx\*H1'+Rwcn) = 0.9995 0.0000 0.0000 0.9948 Primary/Secondary



May 7, 2024

>> [R,Q,P]=rq(inv(Rwcn)\*H1)  $R =$  0 9.1016 -33.2537 0 0 -107.6507  $Q =$  0.4082 -0.5306 -0.7429 -0.8165 0.1517 -0.5571 0.4082 0.8340 -0.3713  $P = 2 1$ ORDER IS REVERSED (Here it is order of users 1 and 3 since 2 was eliminated)  $\Rightarrow$  R1=R(1:2,2:3);  $\gg$  Q1=Q(1:3,2:3); >> Rxxrot=Q1' \*Rxx\*Q1 = 2.3275 0.2251 0.2251 3.1725 >> Phi=lohc(Rxxrot);  $\geq$  DA=diag(diag(R1\*Phi)) = 13.8379 0 0 -191.7414  $\geq$  G=inv(DA)\*R1\*Phi = 1.0000 -4.1971 0 1.0000  $>>$  A=Q1\*inv(R1)\*DA\*G = -0.8067 -1.3902 0.2306 -0.9730 1.2679 -0.5559  $>> A^*A' =$  2.5833 1.1667 -0.2500 1.1667 1.0000 0.8333 -0.2500 0.8333 1.9167

>> S0=DA\*inv(Swcn)\*DA = 1.0e+04 \* 0.0192 0.0000 0.0000 3.6957  $\Rightarrow$  MSWMFunb=inv((S0)-eye(2))\*DA\*J = 0.0000 0.0726 -0.0052 -0.0000 >> Gunb=eye(2)+S0\*inv(S0-eye(2))\*(G-eye(2)) = 1.0000 -4.2191 -0.0000 1.0000  $\frac{1}{2}$  = 0.5\*log2(diag(S0))' = 3.7909 7.5868  $\Rightarrow$  sum(b) = 11.3777 checks  $\Rightarrow$  H\*A = 0.0219 -191.8333 0.0164 -143.8749 13.8379 -58.3825 See Example 2.8.8 or details of below Assign 1 energy unit to User 1, 1/3 to user 3, and now squeeze in 2/3 energy on user 2  $\Rightarrow$  b=0.5\*log2(diag([11/3]) \*diag(S0)) = 3.7909 6.7943 Crosstalk is  $>>$  ct= $1/3*143.9^2$  = 6.8928e+03  $\frac{1}{2}$  >> b2=0.5\*log2(1+(2/3)\*60^2/6892.8) = 0.2155  $\Rightarrow$  b2+sum(b) = 10.8007 < 11.3777 **Energy on secondary reduces rate sum** Rcvr & Data Rate 1  $\sqrt{2/3} + \sqrt{1/3}$ 

Section 2.8.3.5 **May 7, 2024 Sq Root & Precoder** that is in pass space) L11: 8 **Not equal to Rxx Energy not inserted into null space (same on part**

# **System Diagram for this WCN design**





L10: 9

## **Gaussian Vector BC System Diagram**



This design is for any  $R_{xx}$ , but the square root  $Q \cdot \Phi$  is very special and unique; this design is for the  $R_{wcn}$ , no matter the real correlation between receiver noises; U' is number of user components.



Section 2.8.3.4 May 7, 2024 **L10: 10** 

# **Maximum BC rate sum**

#### **Maximum BC rate sum**

- **Maximize**  $\mathcal{I}(x; y)$  through water-filling (but ... presumes receivers can coordinate).
	- This is concave problem that always can be solved for the best input autocorrelation  $R_{xx}$ .
- **Minimize**  $\mathcal{I}_{min}(x; y)$  through worst-case-noise to get  $\mathcal{I}_{wcn}(x; y)$ .
	- This is a convex problem that always can be solved for worst noise autocorrelation  $R_{wcn}$ .

This is a saddle-point problem that produces a max-min = min-max:



#### **bcmax.m**

function [Rxx, Rwcn, bmax] = bcmax(iRxx, H, Lyu) Uses cvx\_wcnoise.m and rate-adaptive waterfill.m (Lagrange Multiplier based) Inputs: - iRxx: initial input autocorrelation array, size is Lx x Lx x N. Only the sum of traces matters, so can initialize to any valid autocorrelation matrix Rxx to run wcnoise. needs to include factor  $N/(N+nu)$  if nu  $\sim=0$  - H: channel response, size is Ly x Lx x N, w/o sqrt(N) normalization - Lyu: array number of antennas at each user; scalar Lyu means same for all Outputs: - Rxx: optimized input autocorrelation, Lx x Lx x N - Rwcn: optimized worst-case noise autocorrelation, with white local noise Ly x Ly x N so IF H is noise-whitened for Rnn, then actual noise is Rwcn^(1/2)\*Rnn\*Rwcn^(\*/2) - b: maximum sum rate/real-dimension - user must mult by 2 for complex case

§ Revisit example from slide L11:8



§ Secondary components' energy is zeroed for this maximum rate sum.



### **Revisit example from L11:6**

```
H = \begin{bmatrix} 80 & 70 \end{bmatrix} 50 60 ];
>> iRxx=[1 .8
     .8 1];
\Rightarrow [Rxx, Rwcn, bmax] = bcmax(iRxx, H, Lyu)
Rxx = 1.0001 0.0082
   0.0082 0.9999
Rwcn =
   1.0000 0.0049
   0.0049 1.0000
bmax = 10.3517 > 9.6430
```
■ Usually converges pretty quickly, not always though – CVX can get finicky when singularity involved.



### **New Example – Singular Rwcnopt**



May 7, 2024

Section 2.8.3.3

L<sub>10</sub>: 15

**Stanford University** 

**igher** 

# **Scalar Duality (BC and MAC)**

PS5.2 - 2.29 scalar BC region

### **Scalar Dual Channels – Same**  $I(x; y)$

- § Dual Channels have:
	- $H_{MAC}$  for MAC

 $\ddot{\cdot}$ 

 $\boldsymbol{\chi}$ 



 $\mathcal{I}(x; y) = \mathcal{I}_{MAC} = \log_2 \left| H_{MAC} \cdot diag\left\{ \mathcal{E}_U^{MAC} \cdots \mathcal{E}_2^{MAC} \mathcal{E}_1^{MAC} \right\} \cdot H_{MAC}^* + I \right|$ 



## **Scalar Duality**

■ Set data rates equal and solve for  $\mathcal{E}_u^{MAC/BC}$ :

$$
\overline{b}_1 = \frac{1}{2} \log_2 \left( 1 + \frac{\mathcal{E}_1^{MAC} \cdot g_1}{1 + \mathcal{E}_2^{MAC} \cdot g_2 + \ldots + \mathcal{E}_U^{MAC} \cdot g_U} \right) \quad \overline{b}_1 = \frac{1}{2} \log_2 \left( 1 + \frac{\mathcal{E}_1^{BC} \cdot g_1}{1} \right)
$$
\n
$$
\overline{b}_2 = \frac{1}{2} \log_2 \left( 1 + \frac{\mathcal{E}_2^{MAC} \cdot g_2}{1 + \mathcal{E}_3^{MAC} \cdot g_3 + \ldots + \mathcal{E}_U^{MAC} \cdot g_U} \right) \quad \overline{b}_2 = \frac{1}{2} \log_2 \left( 1 + \frac{\mathcal{E}_2^{BC} \cdot g_2}{1 + \mathcal{E}_1^{EC} \cdot g_2} \right)
$$
\n
$$
\vdots
$$
\n
$$
\overline{b}_U = \frac{1}{2} \log_2 \left( 1 + \frac{\mathcal{E}_U^{MAC} \cdot g_U}{1} \right)
$$
\n
$$
\vdots
$$
\n
$$
\overline{b}_U = \frac{1}{2} \log_2 \left( 1 + \frac{\mathcal{E}_U^{MAC} \cdot g_U}{1} \right)
$$
\n
$$
\overline{b}_U = \frac{1}{2} \log_2 \left( 1 + \frac{\mathcal{E}_U^{BC} \cdot g_U}{1 + [\mathcal{E}_1^{BC} + \ldots + \mathcal{E}_U^{BC} \cdot \right] \cdot g_U}{1 + \ldots + \mathcal{E}_U^{BC} \cdot g_U}
$$



May 9, 2024

Section 2.8.4 **L11:18** 

# **Corresponding Energies**

$$
\mathcal{E}_1^{BC} = \mathcal{E}_1^{MAC} \cdot \frac{1}{1 + \mathcal{E}_2^{MAC} \cdot g_2 + \dots + \mathcal{E}_U^{MAC} \cdot g_U}
$$
\n
$$
\mathcal{E}_2^{BC} = \mathcal{E}_2^{MAC} \cdot \frac{1 + \mathcal{E}_1^{BC} \cdot g_2}{1 + \mathcal{E}_3^{MAC} \cdot g_3 + \dots + \mathcal{E}_U^{MAC} \cdot g_U}
$$
\n
$$
\vdots = \vdots
$$
\n
$$
\mathcal{E}_U^{BC} = \mathcal{E}_U^{MAC} \cdot \left(1 + \left[\mathcal{E}_1^{BC} + \dots + \mathcal{E}_{U-1}^{BC}\right] \cdot g_U\right)
$$

- By selecting these energies, all user rates are the same (with the order reversal) and running through all such energies that sum to total produces the SAME energy-sum capacity region.
- See proof in text (Theorem 2.8.2 in Section 2.8.4).



Section 2.8.4 **L11: 19** 

## **Revisit Scalar Example**

- § Total energy is 1, instead use dual MAC to investigate BC with:
	- $\mathcal{E}_2^{BC} = 0.25$  (bottom of BC),
	- $\varepsilon_1^{BC} = 0.75$  (top BC), &
	- reversing order  $g_1 = 6400$  and  $g_2 = 2500$ .

$$
\mathcal{E}_2^{MAC} = \frac{\varepsilon_2^{BC}}{1 + \varepsilon_1^{BC} \cdot g_2} = \frac{.25}{1 + 2500 \cdot (.75)} = \frac{1}{7504} = 1.3326 \times 10^{-4} \text{ (top MAC)}
$$
  

$$
\mathcal{E}_1^{MAC} = \mathcal{E}_1^{BC} \cdot \left(1 + g_2 \cdot \mathcal{E}_2^{MAC}\right) = .75 \cdot (1 + 2500/7504) = \frac{7503}{7504} = .9999 = 1 - \mathcal{E}_2^{MAC} \text{ (bottom MAC)}
$$

■ User data rates for this combination are (and were in earlier table found directly for BC).

$$
b_1 = \frac{1}{2} \cdot \log_2 \left( 1 + \frac{\mathcal{E}_1^{MAC} \cdot g_1}{1 + \mathcal{E}_2^{MAC} \cdot g_2} \right) = 6.1144
$$

$$
b_2 = \frac{1}{2} \cdot \log_2 \left( 1 + \frac{\mathcal{E}_2^{MAC} \cdot g_2}{1} \right) = .2074
$$

■ Can use the easier MAC developments to analyze the BC through duality.



Section 2.8.4

L11: 20

# **Continuous-time Scalar BC**

Section 2.8.5

May 9, 2024 21

### **Continuous time/freq Scalar BC**





Design for this MAC, and then find dual

Section 2.8.5 **L11:22** 

## **Scalar Duality**

Replace with integrals and  $\mathcal{E}_{u}^{MAC/BC} \rightarrow S_{u}^{MAC/BC} (f)$ 

May 9, 2024



Section 2.8.5 **L11:23** 

## **Corresponding PSD's**

$$
\bullet \quad \mathcal{E}_u^{MAC/BC} \rightarrow \mathcal{S}_u^{MAC/BC}(f)
$$

■ See proof in notes (Theorem 2.8.2 in Section 2.8.4), but execute with PSD's  $S_u^{MAC/BC}(f)$ 

$$
\begin{aligned}\n\mathcal{E}_1^{BC} &= \mathcal{E}_1^{MAC} \cdot \frac{1}{1 + \mathcal{E}_2^{MAC} \cdot g_2 + \dots + \mathcal{E}_U^{MAC} \cdot g_U} \\
\mathcal{E}_2^{BC} &= \mathcal{E}_2^{MAC} \cdot \frac{1 + \mathcal{E}_2^{BC} \cdot g_2}{1 + \mathcal{E}_3^{MAC} \cdot g_3 + \dots + \mathcal{E}_U^{MAC} \cdot g_U} \\
\vdots &= \vdots \\
\mathcal{E}_U^{BC} &= \mathcal{E}_U^{MAC} \cdot \left(1 + \left[\mathcal{E}_1^{BC} + \dots + \mathcal{E}_{U-1}^{BC}\right] \cdot g_U\right)\n\end{aligned}
$$



Section 2.8.5 **L11: 24** 

# **MAC-set Approach to IC**

Sec 2.9

May 9, 2024 25

## **The Interference Channel (IC)**



## **Prior-User Set (repeat from L7)**

Order vector and its inverse are:

$$
\boldsymbol{\pi}_u = \begin{bmatrix} \pi(U') \\ \vdots \\ \pi(1) \end{bmatrix} \quad \boldsymbol{\pi}_u^{-1} = \begin{bmatrix} U' \\ \vdots \\ 1 \end{bmatrix} \qquad j = \pi(i) \ \rightarrow i = \pi^{-1}(j)
$$

- **•** Prior-User Set is  $\mathbb{P}_{\nu}(\pi) = \{j \mid \pi^{-1}(j) < \pi^{-1}(\nu)\}\$ 
	- That is "all the users before the desired user u in the given order  $\pi$ .
	- Receiver  $u$  best decodes these "prior" users and removes them, while "post" users are noise
	- $\pi$  can be any order in  $\mathbb{P}_{\nu}(\pi)$ , but the most interesting is usually  $\pi_{\nu}$  (receiver  $u$ 's order)



§ Data rates (mutual information bounds) average only those users who are not cancelled are xtalk noise.



$$
\mathcal{I}_{min}(\boldsymbol{\varPi},p_{\mathbf{xy}})=\begin{bmatrix} 4\\20\\1\\2 \end{bmatrix}
$$

# **Maximum number of subusers** *U2*

User  $u$  has maximum bit rate, when all other users are given (cancelled):

$$
b_u \leq \mathcal{I}(x_u; y_u/x_{U\setminus u})
$$

**Stanford University** 

- Although -- this may not be good for the other users  $i \neq u$ .
	- $\mathcal{I}_{min}(\Pi, p_{xy})$  calculation, given a  $\Pi$ , precedes a subsequent **convex-hull** search over all  $\Pi$  to obtain the achievable region  $\mathcal{A}(b, p_{xy})$



- For  $U \geq 2$ , the other users' subuser components may be desirable to decode, but not all  $\rightarrow U'$ !  $\leq (U^2)$ ! for **each** receiver's order may need search/evaluation.
- $\boldsymbol{\Pi}$  maximally has  $(U')^U$  possible choices (in most general case).

At any receiver u, the subuser components separate into two groups for any given order  $\pi_u$ :

(1) those cancelled (or generally conditional probability has specific given values for those components), and

(2) those not cancelled, which are averaged out generally in marginal distributions.

If that choice is made for each user for each receiver, there are thus maximally U choices across all receivers into (1) or (2), so  $U' \leq U^2$ .



## **Example channel – Scalar Gaussian IC**



- Earlier  $H$ , but this time as an IC
	- $.8 > .7$  and  $.6 > .5$
	- Not complete set of orders (4 instead of  $(4!)^2 = 576$
- § Shaded points are interior to line formed by unshaded points





Sec 2.9.1  $\mu_{AV}$  9. 2024

# **IC Rate Region Examples**

Sec 2.9.1

May 9, 2024 30

## **Example continued**





Table 2.8: More example points with the best orders

- The dimension-sharing of the large- $b<sub>u</sub>$  points dominates the other points on the interior.
- § Also check of vertices' derivatives relative to the dimension-sharing line (-1.006) try upper.

$$
\ln(2) \cdot \frac{db_2}{d\mathcal{E}_2} = \frac{3200}{6400 \cdot \mathcal{E}_2 + 1} = 0.4999
$$
\n
$$
\ln(2) \cdot \frac{db_1}{d\mathcal{E}_2} = -\frac{3200 \cdot 2500 \cdot \mathcal{E}_1}{6400 \cdot (\mathcal{E}_2 + 2500 \cdot \mathcal{E}_1 + 1) \cdot (6400 \cdot \mathcal{E}_2 + 1)} = -.1404
$$
\n
$$
\frac{db_2}{db_1} = -3.56
$$

**Vertices all (o) inside pentagon for this example**

- 
- For 2 vertices if magnitude of slope is less than 1, then upper point and otherwise lower point (or whole line).
	- May 9, 2024 • Could check other vertex also, but if curvature is within the line already (convex), then no need. Sec 2.9.1 May 9, 2024 May 9, 2024

### **Energy-sum IC extension**



# **So-called "weak" symmetric IC**

 $b<sub>2</sub>$ 

2

Achievable Region when  $\mathcal{E}_1 = \mathcal{E}_2 = 1$ 

 $\alpha \geq 1$ 

 $\boldsymbol{\Pi} = \begin{bmatrix} 2 & 2 \\ 4 & 4 \end{bmatrix}$ 

 $1 + \frac{n-11}{2}$ 

 $\bm{b} =$ 

 $rac{1}{2}$ · $\log_2(1 +$ 

 $rac{1}{2}$  ·  $\log_2(1 +$ 

 $\bm{b} =$ 

!  $\overline{2}$ !  $\overline{2}$ 

 $\varepsilon_{\rm z}$  $1 + \alpha^2 \cdot \mathcal{E}_1$ 

 $\varepsilon_{1}$ 

- $y_2$  $\begin{bmatrix} 2 & 2 \\ y_1 \end{bmatrix}$  =  $\begin{bmatrix} 1 & \alpha \\ \alpha & 1 \end{bmatrix} \cdot \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} +$  $n<sub>2</sub>$  $n<sub>1</sub>$  $R_{nn} = I$  and  $\mathcal{E} \leq 1$
- When  $\alpha \rightarrow 0$  , there is no crosstalk and so  $\mathcal{C}_{IC}(\boldsymbol{b})$  is a square.
- When  $\alpha > 1$  ,  $\mathcal{C}_{IC}(\boldsymbol{b})$  is a pentagon.
- $1 + \alpha^2 \cdot \mathcal{E}_2$  $\boldsymbol{\Pi} = \begin{bmatrix} 1 & 2 \\ 2 & 4 \end{bmatrix}$ When  $0 < \alpha < 1$ ,  $\mathcal{C}_{IC}(\boldsymbol{b})$  is intermediate  $\alpha > 1$ 2 1  $\mathcal{E}_2$  $rac{1}{2} \cdot \log_2 \left(1 + \right)$ ,  $1 + \alpha^2 \cdot \mathcal{E}_1$  $\boldsymbol{\Pi} = \begin{bmatrix} 1 & 1 \\ 2 & 2 \end{bmatrix}$ min  $\alpha^2 \cdot \varepsilon_2$ 2 2  $rac{1}{2} \cdot \log_2 (1 +$  $\boldsymbol{0}$  $1 + \mathcal{E}_1$  $b<sub>1</sub>$  $\boldsymbol{0}$  $\mathcal{E}_1$  $rac{1}{2} \cdot \log_2 \left(1 + \right)$ , 1  $1 + \alpha^2 \cdot \mathcal{E}_2$ These two are same for  $\alpha = 1$  and equal energy, min 2  $\alpha^2 \cdot \varepsilon_1$ and determine Imin vector possibilities  $rac{1}{2} \cdot \log_2 (1 +$  $1 + \mathcal{E}_2$ **Stanford University** Sec 2.9.1 May 9, 2024

### **Vector Gaussian IC Example**

2 users and *H* is 4 x 2 
$$
y = \begin{bmatrix} y_2 \\ y_1 \end{bmatrix} = \begin{bmatrix} H_2 \\ H_1 \end{bmatrix} \cdot \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} + \begin{bmatrix} n_2 \\ n_1 \end{bmatrix}
$$

$$
R_{nn} = 01 \cdot I
$$

$$
H_2 = [\mathbf{h}_{22} \quad \mathbf{h}_{21}] = \begin{bmatrix} .9 & .3 \\ .3 & .8 \end{bmatrix}
$$

$$
H_1 = [\mathbf{h}_{12} \quad \mathbf{h}_{11}] = \begin{bmatrix} .8 & .7 \\ .6 & .5 \end{bmatrix}
$$

 $>> H2 = 19$  3 3 8 ]; >> Rb2inv=H2'\*H2+diag([1 1]); >> Gbar2=chol(Rb2inv); >> G2=inv(diag(diag(Gbar2)))\*Gbar2; >> S02=diag(diag(Gbar2))\*diag(diag(Gbar2));  $\gg 0.5^*$ log2(diag(S02)) =  $b2 = 3.2539$  $h1 = 2.7526$  $\gg$  H1 = [8 7 6 5 ]; >> Rb1inv=H1'\*H1+diag([1 1]); >> Gbar1=chol(Rb1inv); >> S01=diag(diag(Gbar1))\*diag(diag(Gbar1));  $\gg 0.5$ \*log2(diag(S01)) =  $h2 = 3.3291$  $b1 = 0.4128$ 

```
>> J2=hankel([0 1]);
>> Rb2inv=J2*H2'*H2*J2+diag([1 1]);
  74 51
  51 91
>> Gbar2=chol(Rb2inv);
>> S02=diag(diag(Gbar2))*diag(diag(Gbar2));
\gg 0.5^*log2(diag(S02)) =
h1 = 3.1047h2 = 2.9018>> Rb1inv=J2*H1'*H1*J2+diag([1 1]);
>> Gbar1=chol(Rb1inv);
>> S01=diag(diag(Gbar1))*diag(diag(Gbar1));
\gg 0.5*log2(diag(S01)) =
h1 = 3.1144b2 = 0.6275
```
Try various energy points With the same orders



## **4x2 IC Example continued**



Vary energies near max points to see if local points above or below dimension-sharing, so curved boundary or flat

Since the line has slope magnitude less than 1, then the curvature is above this line with max rate sum at magnitude 1

#### PS 5.4 (2.31) – IC channel has mix, one 2x2 user and one scalar user

May 9, 2024 Sec 2.9.2 May 9, 2024 **L11: 35** 

# **Chris AF's MU\_IC.m**



May 9, 2024

L11: 36



# **End Lecture 11**

# **Nesting, DAS, cellfree & Relay**

#### **Multiuser Nesting**



**CONTROL** 

**Stanford University** 

**But easier**

#### **Single-Stage Relay Channel**

Conceptually uses what we know already and introduces sub-users at  $k<sub>r</sub>$  relay points

Sec 2.11.1



# **Reflective Intelligent Surfaces (RIS)**



- The RIS matrix  $Q_H$  satisfies  $||Q_H||_F^2 \leq G_H$  , the RIS gain it may also satisfy
	- $Q_H$  is unitary matrix (preserves energy)
	- $Q_H$  is diagonal, and usually also unitary, to be phase/gain-only adjustment on each antenna port (in-to-out)
	- $Q_H$  has individual elements restricted
- For a given  $R_{xx}$ , maximize over  $Q_H$  $\mathcal{I}(\boldsymbol{y};\boldsymbol{x}) = \log_2 |R_{n, RIS} + H_{RIS} \cdot R_{\boldsymbol{xx}} \cdot H_{BIS}^*|$
- For a given  $Q_H$ , maximize the same over  $R_{xx}$ 
	- Iterate?  $\rightarrow$  not convex in the  $Q_H$ , this needs work.

Sec 2.11.4

May 9, 2024

$$
R_{\boldsymbol{n}\boldsymbol{n},RIS}=\left[\begin{array}{cc} R_{\boldsymbol{n}\boldsymbol{n}} & 0 \\ 0 & R_{\boldsymbol{n}\boldsymbol{n},out} + Q_H \cdot R_{\boldsymbol{n}\boldsymbol{n}in} \cdot Q_H^* \end{array}\right]
$$