



STANFORD

Lecture 10

BC, Interference & General MU Channels

May 4, 2026

JOHN M. CIOFFI

Hitachi Professor Emeritus of Engineering

Instructor EE379B – Spring 2026

Announcements & Agenda

- Announcements

- PS #5 due May 14

- Agenda

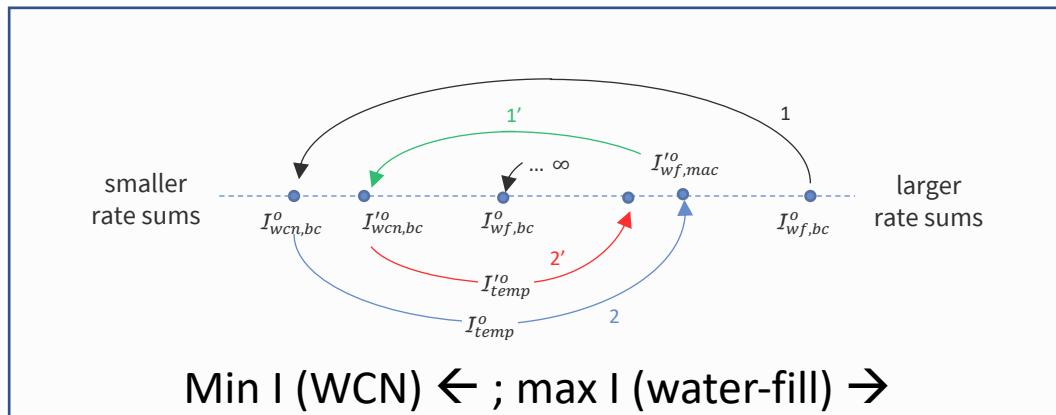
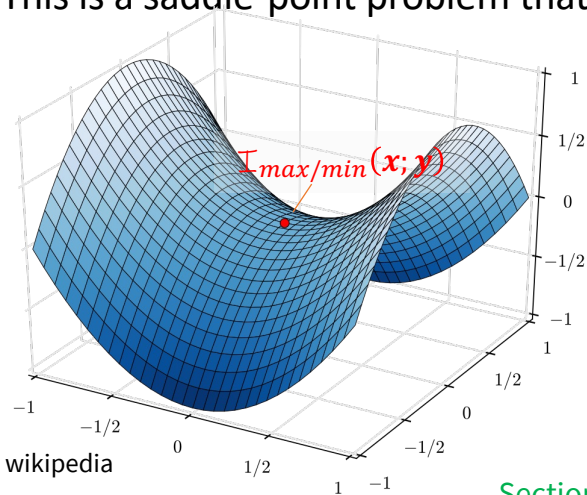
- Maximum BC rate sum
- Scalar Duality (BC and MAC)
- Continuous-time scalar BC
- MAC-set approach to IC
 - Examples
- IC maximum rate sum and energy sums
- Nesting, DAS, cellfree, & relay



Maximum BC rate sum

Maximum BC rate sum

- Maximize** $\mathcal{I}(\mathbf{x}; \mathbf{y})$ through water-filling (but ... presumes receivers can coordinate).
 - This is concave problem that always can be solved for the best input autocorrelation R_{xx} .
- Minimize** $\mathcal{I}_{min}(\mathbf{x}; \mathbf{y})$ through worst-case-noise to get $\mathcal{I}_{wcn}(\mathbf{x}; \mathbf{y})$.
 - This is a convex problem that always can be solved for worst noise autocorrelation R_{wcn} .
- This is a saddle-point problem that produces a max-min = min-max:



bcmax.m

```
function [Rxx, Rwcn, bmax] = bcmax(iRxx, H, Lyu)
```

Uses `cvx_wnoise.m` and `rate-adaptive waterfill.m` (Lagrange Multiplier based)

Inputs:

- `iRxx`: initial input autocorrelation array, size is $L_x \times L_x \times N$. Only the sum of traces matters, so can initialize to any valid autocorrelation matrix `Rxx` to run `wnoise`. needs to include factor $N/(N+nu)$ if $nu \approx 0$
- `H`: channel response, size is $L_y \times L_x \times N$, w/o \sqrt{N} normalization
- `Lyu`: array number of antennas at each user; scalar `Lyu` means same for all

Outputs:

- `Rxx`: optimized input autocorrelation, $L_x \times L_x \times N$
- `Rwcn`: optimized worst-case noise autocorrelation, with white local noise $L_y \times L_y \times N$
so IF `H` is noise-whitened for `Rnn`, then actual noise is $Rwcn^{(1/2)} * Rnn * Rwcn^{(*/2)}$
- `b`: maximum sum rate/real-dimension - user must mult by 2 for complex case

- Revisit example from slide L9:24

```
H = [ 80 70
      50 60 ];
>> iRxx=[1 .8
         .8 1];
>> [Rxx, Rwcn, bmax] = bcmax(iRxx, H, Lyu)

Rxx =
    1.0001    0.0082
    0.0082    0.9999
Rwcn =
    1.0000    0.0049
    0.0049    1.0000
bmax = 10.3517 > 9.6430
```

- Usually converges pretty quickly, not always though – CVX can get finicky when singularity involved.



Revisit example from L9:26

```
iRxx =  
  3  0  0  
  0  4  0  
  0  0  2  
>> H =  
  80  60  40  
  60  45  30  
  20  20  20  
[RxxA, RwcA, bmax] = bcmax(iRxx, H, 1)  
RxxA =  
  3.7515  1.5032 -0.7451  
  1.5032  1.5019  1.5007  
 -0.7451  1.5007  3.7465  
RwcA =  
  1.0000  0.7500  0.0008  
  0.7500  1.0000  0.0006  
  0.0008  0.0006  1.0000  
bmax = 12.1084 (> 11.3777 that occurred earlier)
```

- Secondary components' energy is zeroed for this maximum rate sum.



New Example – Singular Rwcnopt

```
>> H
    1.0719 -0.8627 -0.1901  0.2952
    1.0498 -0.7245  0.2568  0.2757
   -0.4586  0.5595  1.0027  0.0530
    0.4107  0.0496  0.3965 -0.7740
[F,L,M]=svd(H);
L =
    2.0671    0    0    0
    0    1.1449    0    0
    0    0    0.8130    0
    0    0    0    0.0000
H=F(:,1)*M(:,1)*L(1,1)+F(:,2)*M(:,2)*L(2,2)+F(:,3)*M(:,3)*L(3,3);

[Rxxopt, Rwcnopt, bmax] = bccmax(eye(4), H, 1)

Rxxopt =
    1.1559 -0.7381  0.0995 -0.0691
   -0.7381  0.6399  0.2862 -0.2207
    0.0995  0.2862  1.3091 -0.0326
   -0.0691 -0.2207 -0.0326  0.8951
Rwcnopt =
    1.0000  0.9163 -0.4792 -0.0191
    0.9163  1.0000 -0.0991  0.1251
   -0.4792 -0.0991  1.0000  0.1044
   -0.0191  0.1251  0.1044  1.0000
bmax = 2.1911
>> det(Rwcnopt) % = 1.5337e-09
```

Singular Rwcn

```
inv(Rwcnopt) - inv(H*Rxxopt*H'+Rwcnopt)
   -69.9931  62.4186 -26.7687  -6.3512
    62.4186 -54.9656  23.8385  5.6560
   -26.7687  23.8385  -9.5873  -2.4256
   -6.3512  5.6560  -2.4256  -0.1050
[Rwcnopt, sumRatebar, S1, S2, S3, S4] =
cvx_wcnoise(Rxxopt, H, [1 1 1 1])
Rwcnopt =
    1.0000  0.9163 -0.4792 -0.0191
    0.9163  1.0000 -0.0991  0.1251
   -0.4792 -0.0991  1.0000  0.1044
   -0.0191  0.1251  0.1044  1.0000
sumRatebar = 2.1911
rank(H) = 3
>> S3+S4(1:4,1:4) =
    0.0978    0    0    0
    0    0.6204    0    0
    0    0    0.6360    0
    0    0    0    0.4705
rank(H) = 3
>> Htilde=pinv(Rwcnopt)*H;
>> [R,Q,P]=rq(Htilde);
R =
    0.0000  0.0211 -0.2595  0.1213
    0    -0.9411 -0.0435  0.0071
    0    0    -1.0542  0.0496
    0    0    0    -1.1499
P = 1  4  2  3
```

```
>> [V,D]=eig(Rxxopt)
```

V =

```
   -0.5342  -0.7457  -0.3635  -0.1626
   -0.7873  0.6052  -0.0344  -0.1124
    0.2071  0.2556  -0.9213  0.2073
   -0.2278  -0.1108  0.1336  0.9581
```

D =

```
    0.0000    0    0    0
    0  1.7105    0    0
    0    0  1.3639    0
    0    0    0  0.9256
```

Confirms that best input is also singular – it should never have higher rank than number of primary user (components).

Scalar Duality (BC and MAC)

PS5.2 - 2.29 scalar BC region

Scalar Dual Channels – Same $I(x; y)$

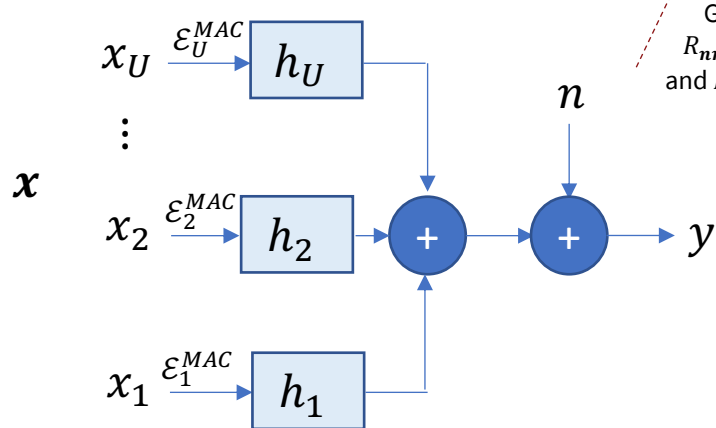
■ Dual Channels have:

- H_{MAC} for MAC
- $H_{BC} = H_{MAC}^* \cdot J$ for dual-of-MAC as a BC.

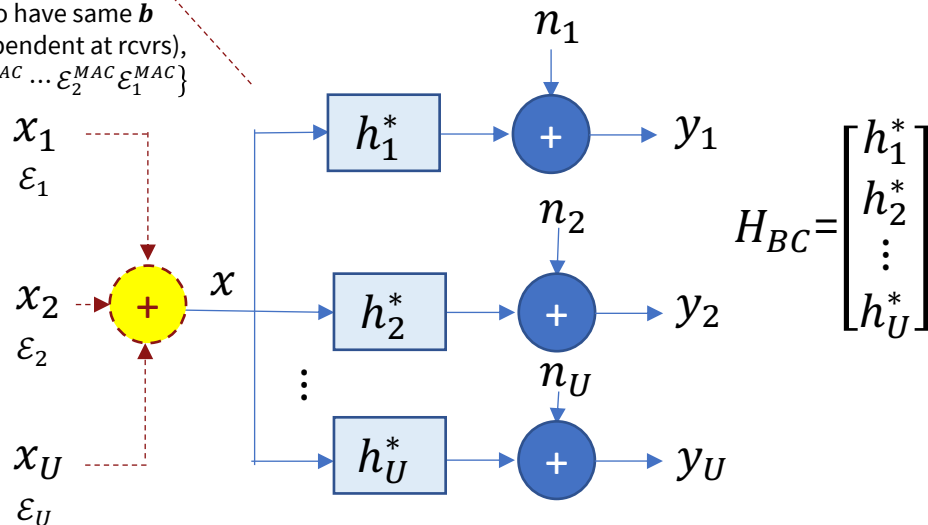
$$I(x; y) = I_{MAC} = \log_2 |H_{MAC} \cdot \text{diag}\{\varepsilon_U^{MAC} \dots \varepsilon_2^{MAC} \varepsilon_1^{MAC}\} \cdot H_{MAC}^* + I|$$

$$I_{BC}(x; y) = \sum_{u=1}^U \log_2 \left(\frac{\varepsilon_u^{BC} \cdot |h_u|^2}{1 + |h_u|^2 \cdot \sum_{i=u+1}^U \varepsilon_i^{BC}} + 1 \right)$$

$$H_{MAC} = [h_U \dots h_2 h_1]$$



Goal: For duals to have same \mathbf{b}
 R_{nn} is white (independent at rcvrs),
 and $R_{xx} = \text{diag}\{\varepsilon_U^{MAC} \dots \varepsilon_2^{MAC} \varepsilon_1^{MAC}\}$



$$\begin{bmatrix} \varepsilon_U^{MAC} \\ \vdots \\ \varepsilon_1^{MAC} \end{bmatrix} \preceq \varepsilon$$

$$\sum_{u=1}^U \varepsilon_u^{BC} = \varepsilon_x = \sum_{u=1}^U \varepsilon_u^{MAC} \quad \text{(scalar case)}$$



Scalar Duality

- Set data rates equal and solve for $\epsilon_u^{MAC/BC}$:

	MAC	BC
<div style="border: 1px solid black; background-color: yellow; padding: 5px; display: inline-block;">order intentionally reversed</div>	$\bar{b}_1 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_1^{MAC} \cdot g_1}{1 + \epsilon_2^{MAC} \cdot g_2 + \dots + \epsilon_U^{MAC} \cdot g_U} \right)$	$\bar{b}_1 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_1^{BC} \cdot g_1}{1} \right)$
	$\bar{b}_2 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_2^{MAC} \cdot g_2}{1 + \epsilon_3^{MAC} \cdot g_3 + \dots + \epsilon_U^{MAC} \cdot g_U} \right)$	$\bar{b}_2 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_2^{BC} \cdot g_2}{1 + \epsilon_1^{BC} \cdot g_2} \right)$
	\vdots	\vdots
	$\bar{b}_U = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_U^{MAC} \cdot g_U}{1} \right)$	$\bar{b}_U = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_U^{BC} \cdot g_U}{1 + [\epsilon_1^{BC} + \dots + \epsilon_{U-1}^{BC}] \cdot g_U} \right)$



Corresponding Energies

$$\begin{aligned}\mathcal{E}_1^{BC} &= \mathcal{E}_1^{MAC} \cdot \frac{1}{1 + \mathcal{E}_2^{MAC} \cdot g_2 + \dots + \mathcal{E}_U^{MAC} \cdot g_U} \\ \mathcal{E}_2^{BC} &= \mathcal{E}_2^{MAC} \cdot \frac{1 + \mathcal{E}_1^{BC} \cdot g_2}{1 + \mathcal{E}_3^{MAC} \cdot g_3 + \dots + \mathcal{E}_U^{MAC} \cdot g_U} \\ &\vdots = \vdots \\ \mathcal{E}_U^{BC} &= \mathcal{E}_U^{MAC} \cdot (1 + [\mathcal{E}_1^{BC} + \dots + \mathcal{E}_{U-1}^{BC}] \cdot g_U)\end{aligned}$$

- By selecting these energies, all user rates are the same (with the order reversal) and running through all such energies that sum to total produces the SAME energy-sum capacity region.
- See proof in text (Theorem 2.8.2 in Section 2.8.4).



Revisit Scalar Example

- Total energy is 1, instead use dual MAC to investigate BC with:

- $\varepsilon_2^{BC} = 0.25$ (bottom of BC),
- $\varepsilon_1^{BC} = 0.75$ (top BC), &
- reversing order $g_1 = 6400$ and $g_2 = 2500$.

$$\varepsilon_2^{MAC} = \frac{\varepsilon_2^{BC}}{1 + \varepsilon_1^{BC} \cdot g_2} = \frac{.25}{1 + 2500 \cdot (.75)} = \frac{1}{7504} = 1.3326 \times 10^{-4} \text{ (top MAC)}$$

$$\varepsilon_1^{MAC} = \varepsilon_1^{BC} \cdot (1 + g_2 \cdot \varepsilon_2^{MAC}) = .75 \cdot (1 + 2500/7504) = \frac{7503}{7504} = .9999 = 1 - \varepsilon_2^{MAC} \text{ (bottom MAC)}$$

- User data rates for this combination are (and were in earlier table found directly for BC).

$$b_1 = \frac{1}{2} \cdot \log_2 \left(1 + \frac{\varepsilon_1^{MAC} \cdot g_1}{1 + \varepsilon_2^{MAC} \cdot g_2} \right) = 6.1144$$

$$b_2 = \frac{1}{2} \cdot \log_2 \left(1 + \frac{\varepsilon_2^{MAC} \cdot g_2}{1} \right) = .2074$$

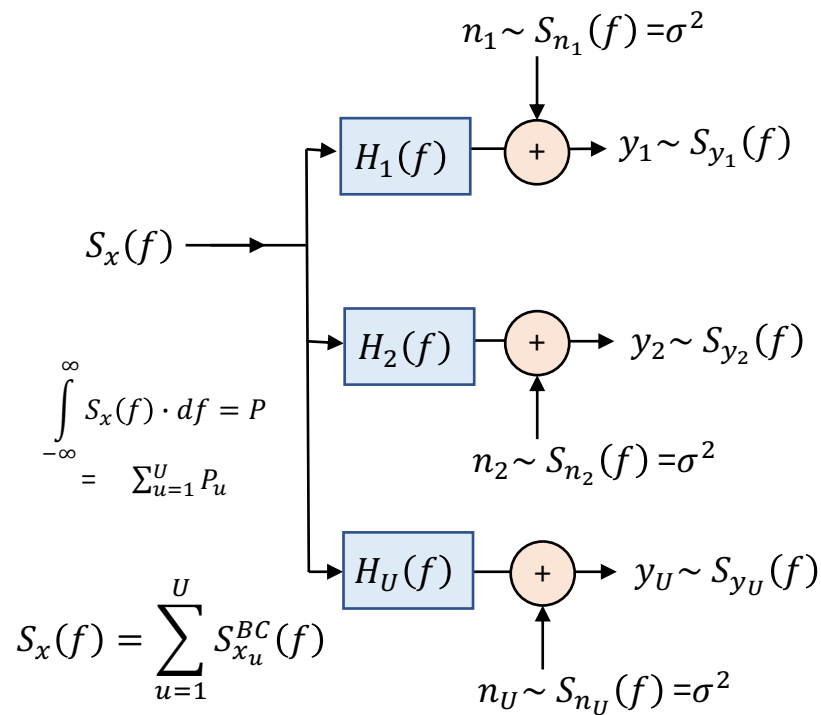
- Can use the easier MAC developments to analyze the BC through duality.



Continuous-time Scalar BC

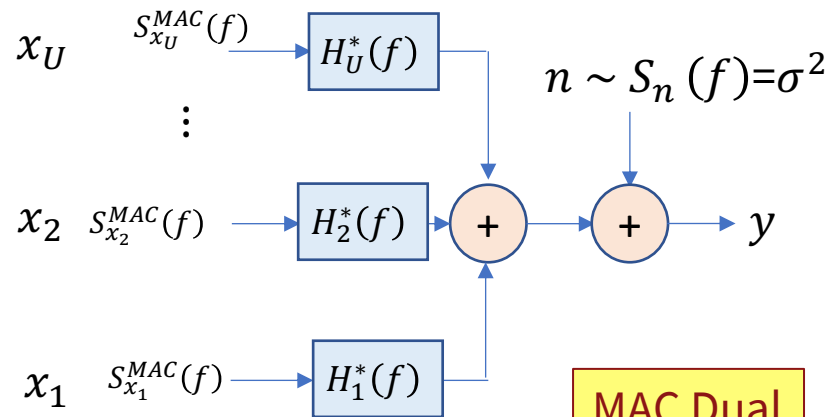
Section 2.8.5

Continuous time/freq Scalar BC



$$\ln S_{0,u} = \int_{-\infty}^{\infty} \ln \left[1 + \frac{|H_u(f)|^2}{S_{n_u}(f)} \right] \cdot df ; (if < \infty)$$

$$SNR_{geo,u} = P_u \cdot S_{0,u}$$



$$\mathbf{x} \quad \sum_{u=1}^U \int_{-\infty}^{\infty} S_{x_u}^{MAC}(f) \cdot df = P$$

Design for this MAC, and then find dual



Scalar Duality

- Replace with integrals and $\epsilon_u^{MAC/BC} \rightarrow S_u^{MAC/BC}(f)$

MAC	BC
$b_1 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_1^{MAC} \cdot g_1}{1 + \epsilon_2^{MAC} \cdot g_2 + \dots + \epsilon_U^{MAC} \cdot g_U} \right)$	$b_1 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_1^{BC} \cdot g_1}{1} \right)$ 1 ND user
$b_2 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_2^{MAC} \cdot g_2}{1 + \epsilon_3^{MAC} \cdot g_3 + \dots + \epsilon_U^{MAC} \cdot g_U} \right)$	$b_2 = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_2^{BC} \cdot g_2}{1 + \epsilon_1^{BC} \cdot g_2} \right)$
\vdots	\vdots
$b_U = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_U^{MAC} \cdot g_U}{1} \right)$ 1 ND user	$b_U = \frac{1}{2} \log_2 \left(1 + \frac{\epsilon_U^{BC} \cdot g_U}{1 + [\epsilon_1^{BC} + \dots + \epsilon_{U-1}^{BC}] \cdot g_U} \right)$

order intentionally reversed



Corresponding PSD's

- $\mathcal{E}_u^{MAC/BC} \rightarrow S_u^{MAC/BC}(f)$
- See proof in notes (Theorem 2.8.2 in Section 2.8.4), but execute with PSD's $S_u^{MAC/BC}(f)$

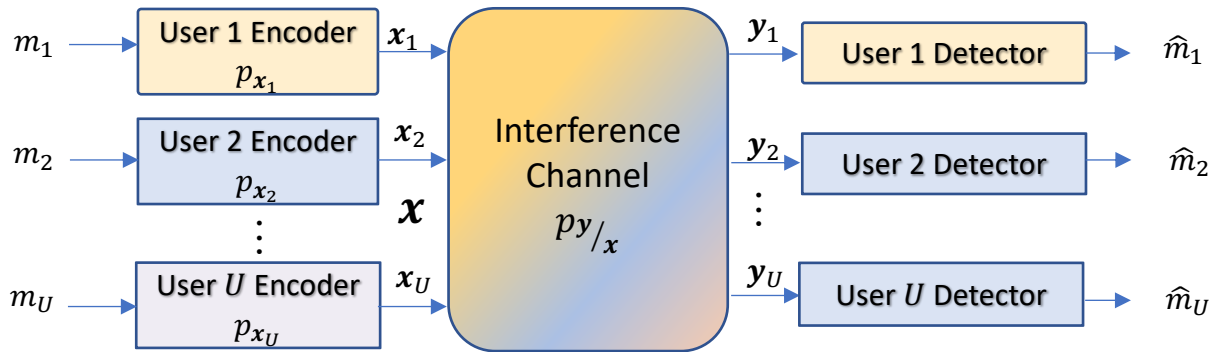
$$\begin{aligned}\mathcal{E}_1^{BC} &= \mathcal{E}_1^{MAC} \cdot \frac{1}{1 + \mathcal{E}_2^{MAC} \cdot g_2 + \dots + \mathcal{E}_U^{MAC} \cdot g_U} \\ \mathcal{E}_2^{BC} &= \mathcal{E}_2^{MAC} \cdot \frac{1 + \mathcal{E}_2^{BC} \cdot g_2}{1 + \mathcal{E}_3^{MAC} \cdot g_3 + \dots + \mathcal{E}_U^{MAC} \cdot g_U} \\ &\vdots = \vdots \\ \mathcal{E}_U^{BC} &= \mathcal{E}_U^{MAC} \cdot (1 + [\mathcal{E}_1^{BC} + \dots + \mathcal{E}_{U-1}^{BC}] \cdot g_U)\end{aligned}$$



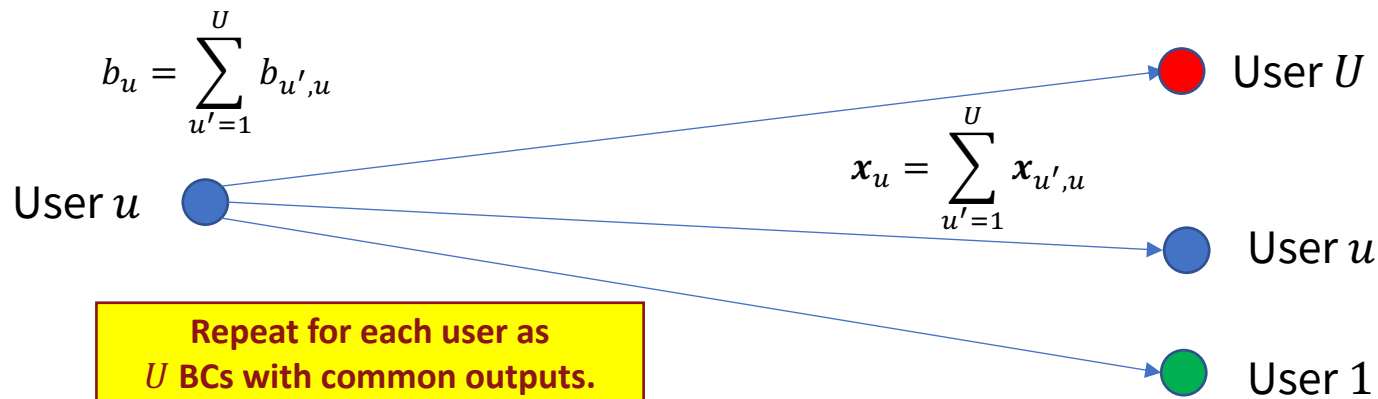
MAC-set Approach to IC

Sec 2.9

The Interference Channel (IC)



Refine Capacity Region Def:
with order-approach:
Assume each user must
use a single-user $\Gamma = 0$ dB
Code.



U^2 potential
sub-user
channels/
decodes

Repeat for each user as
 U BCs with common outputs.

Also could be
 U MACs with
common inputs.
 u' is receiver index



Prior-User Set (repeat from L7)

Order vector and its inverse are:

$$\boldsymbol{\pi}_u = \begin{bmatrix} \pi(U') \\ \vdots \\ \pi(1) \end{bmatrix} \quad \boldsymbol{\pi}_u^{-1} = \begin{bmatrix} U' \\ \vdots \\ 1 \end{bmatrix} \quad j = \pi(i) \rightarrow i = \pi^{-1}(j)$$

Prior-User Set is $\mathbb{P}_u(\boldsymbol{\pi}) = \{j \mid \boldsymbol{\pi}^{-1}(j) < \boldsymbol{\pi}^{-1}(u)\}$

- That is “all the users before the desired user u in the given order $\boldsymbol{\pi}$.”
- Receiver u best decodes these “prior” users and removes them, while “post” users are noise
- $\boldsymbol{\pi}$ can be any order in $\mathbb{P}_u(\boldsymbol{\pi})$, but the most interesting is usually $\boldsymbol{\pi}_u$ (receiver u 's order)

rcvr/ User i	$\pi_4(i)$	$\pi_3(i)$	$\pi_2(i)$	$\pi_1(i)$
$i = 4$	3	3	4	3
$i = 3$	4	2	3	2
$i = 2$	1	4	2	1
$i = 1$	2	1	1	4
$\mathbb{P}_u(\boldsymbol{\pi}_u)$	{1,2}	{2,4,1}	{1}	{4}

$$\boldsymbol{\Pi} = \begin{bmatrix} 3 & 3 & 4 & 3 \\ 4 & 2 & 3 & 2 \\ 1 & 4 & 2 & 1 \\ 2 & 1 & 1 & 4 \end{bmatrix}$$

really should be 16
with full subusers,
but simplify here to 4

- Data rates (mutual information bounds) average only those users who are not cancelled are xtalk noise.

\mathfrak{I}	\mathfrak{I}_4	\mathfrak{I}_3	\mathfrak{I}_2	\mathfrak{I}_1
top	∞	$\mathfrak{I}_3(3/1,2,4)$ 20	∞	∞
	$\mathfrak{I}_4(4/1,2)$ 10	$\mathfrak{I}_3(2/1,4)$ 9	∞	∞
	$\mathfrak{I}_4(1/2)$ 5	$\mathfrak{I}_3(4/1)$ 4	$\mathfrak{I}_2(2/1)$ 4	$\mathfrak{I}_1(1/4)$ 2
bottom	$\mathfrak{I}_4(2)$ 1	$\mathfrak{I}(1)$ 2	$\mathfrak{I}_2(1)$ 2	$\mathfrak{I}_1(4)$ 5

$$\mathfrak{I}_{\min}(\boldsymbol{\Pi}, p_{xy}) = \begin{bmatrix} 4 \\ 20 \\ 1 \\ 2 \end{bmatrix}$$



Maximum number of subusers U^2

- User u has maximum bit rate, when all other users are given (cancelled):

$$b_u \leq \mathcal{I}(x_u; y_u/x_{U \setminus u})$$

- Although -- this may not be good for the other users $i \neq u$.

- $\mathcal{I}_{\min}(\Pi, p_{xy})$ calculation, given a Π , precedes a subsequent **convex-hull** search over all Π to obtain the achievable region $\mathcal{A}(\mathbf{b}, p_{xy})$

$$\mathbf{b} = \sum_{i=1}^U \alpha_i \cdot \mathbf{b}_i \quad 1 = \sum_{i=1}^U \alpha_i \quad \text{convex-hull}$$

No more than U terms needed for U -dimensional \mathbf{b}

$$\mathbf{b} = \begin{bmatrix} b_U \\ \vdots \\ b_1 \end{bmatrix} \quad b_u = \sum_{i=1}^U b_{i,u} \quad \text{Sum of subcomponents}$$

Will do this for $\mathcal{I}_{\min}(\Pi, p_{xy})$ values over possible Π 's

- For $U \geq 2$, the other users' subuser components may be desirable to decode, but not all $\rightarrow U'! \leq (U^2)!$ for **each** receiver's order may need search/evaluation.
- Π maximally has $(U'!)^U$ possible choices (in most general case).

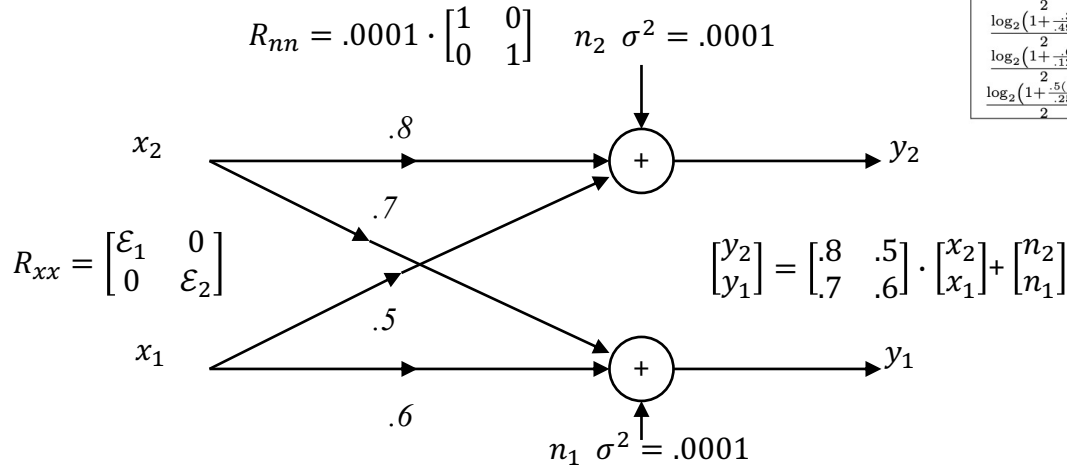
At any receiver u , the subuser components separate into two groups for any given order π_u :

- (1) those cancelled (or generally conditional probability has specific given values for those components), and
- (2) those not cancelled, which are averaged out generally in marginal distributions.

If that choice is made for each user for each receiver, there are thus maximally U choices across all receivers into (1) or (2), so $U' \leq U^2$.



Example channel – Scalar Gaussian IC



$\frac{\log_2(1 + \frac{.64}{.0001})}{2} = 6.3220$	$\frac{\log_2(1 + \frac{.36}{.0001})}{2} = 5.9071$	$\frac{\log_2(1 + \frac{.64}{.2501})}{2} = 0.9157$	$\frac{\log_2(1 + \frac{.25}{.6401})}{2} = 0.2378$
$\frac{\log_2(1 + \frac{.36}{.4901})}{2} = 0.3973$	$\frac{\log_2(1 + \frac{.49}{.3601})}{2} = 0.6196$	$\frac{\log_2(1 + \frac{(.5) \cdot .64}{.0001})}{2} = 5.8222$	$\frac{\log_2(1 + \frac{(.5) \cdot .36}{.0001})}{2} = 5.4073$
$\frac{\log_2(1 + \frac{.64}{.1251})}{2} = 1.3063$	$\frac{\log_2(1 + \frac{.18}{.4901})}{2} = 0.2257$	$\frac{\log_2(1 + \frac{.49}{.1801})}{2} = 0.9478$	$\frac{\log_2(1 + \frac{.36}{.2451})}{2} = 0.6519$
$\frac{\log_2(1 + \frac{.5 \cdot (.64)}{.2501})}{2} = 0.5944$	$\frac{\log_2(1 + \frac{.5 \cdot (.49)}{.3601})}{2} = 0.3744$	$\frac{\log_2(1 + \frac{.5 \cdot (.25)}{.6401})}{2} = 0.1287$	$\frac{\log_2(1 + \frac{.25}{.3201})}{2} = 0.4163$

Table 2.6: Some useful calculations for the upcoming example

**Convex combos of
 These 2 vertices
 Probably includes
 Other sub user combos
 (because they are big)**

- Earlier H , but this time as an IC
 - $.8 > .7$ and $.6 > .5$
 - Not complete set of orders (4 instead of $(4!)^2 = 576$)
- **Shaded points** are interior to line formed by unshaded points

Π	u	\mathcal{E}_u	$\mathbb{P}_u(\pi_u)$	$\mathcal{I}_u(x_u; y_u / \mathbb{P}_u(\pi_u))$	$\mathcal{I}_{-u}(x_u; y_{-u} / \mathbb{P}_{-u}(\pi_{-u}))$	$\mathcal{I}_{min,u}(\Pi, \mathcal{E})$
$\begin{bmatrix} 2 & 2 \\ 1 & 1 \end{bmatrix}$	2	1	$\{1\}$	6.322	∞	6.322
$\begin{bmatrix} 1 & 1 \\ 2 & 2 \end{bmatrix}$	1	1	\emptyset	.3973	.2378	.2378
$\begin{bmatrix} 1 & 1 \\ 2 & 2 \end{bmatrix}$	2	1	\emptyset	.9157	.6196	.6196
$\begin{bmatrix} 2 & 2 \\ 1 & 1 \end{bmatrix}$	1	1	$\{2\}$	5.9071	∞	5.9071
$\begin{bmatrix} 1 & 2 \\ 2 & 1 \end{bmatrix}$	2	1	\emptyset	.9157	∞	.9157
$\begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix}$	1	1	\emptyset	.3973	∞	.3973
$\begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix}$	2	1	$\{1\}$	6.322	.2378	.2378
$\begin{bmatrix} 1 & 2 \\ 2 & 1 \end{bmatrix}$	1	1	$\{2\}$	5.9071	.6196	.6196

Table 2.7: Evaluation of \mathcal{I}_{min} for different orders.



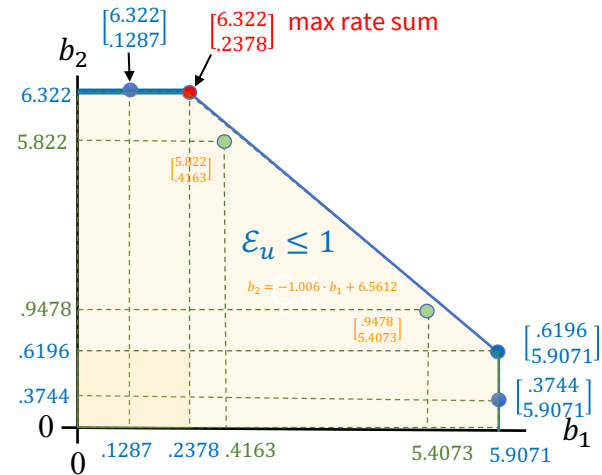
IC Rate Region Examples

Sec 2.9.1

Example continued

Π	u	\mathcal{E}_u	$\mathbb{P}_u(\pi_u)$	$\mathcal{I}_u(x_u; y_u / \mathbb{P}_u(\pi_u))$	$\mathcal{I}_u(x_{-u}; y_{-u} / \mathbb{P}_{-u}(\pi_{-u}))$	$\mathcal{I}_{\min, u}(\Pi, \mathcal{E})$
[2 2]	2	1	{1}	6.322	∞	6.322
[1 1]	1	0.5	\emptyset	.2257	.1287	.1287
[1 1]	2	1	\emptyset	1.3063	.9478	.9478
[2 2]	1	0.5	{2}	5.4073	∞	5.4073
[2 2]	2	0.5	{1}	5.822	∞	5.822
[1 1]	1	1	\emptyset	.6519	.4163	.4163
[1 1]	2	0.5	\emptyset	5.822	.3744	.3744
[2 2]	1	1	{1}	5.9071	∞	5.9071
[2 2]	2	1	{1, 2}	6.322	∞	6.322
[1 1]	1	0.95	{1}	.3818	.2276	.2276
[1 1]	2	1	{2}	.9425	.6412	.6412
[2 2]	1	0.95	{1, 2}	5.8701	∞	5.8701

Table 2.8: More example points with the best orders



- The dimension-sharing of the large- b_u points dominates the other points on the interior.
- Also – check of vertices' derivatives relative to the dimension-sharing line (-1.006) – try upper.

$$\ln(2) \cdot \frac{db_2}{d\mathcal{E}_2} = \frac{3200}{6400 \cdot \mathcal{E}_2 + 1} = 0.4999$$

$$\ln(2) \cdot \frac{db_1}{d\mathcal{E}_2} = -\frac{3200 \cdot 2500 \cdot \mathcal{E}_1}{6400 \cdot (\mathcal{E}_2 + 2500 \cdot \mathcal{E}_1 + 1) \cdot (6400 \cdot \mathcal{E}_2 + 1)} = -.1404$$

$$\frac{db_2}{db_1} = -3.56$$

Vertices all (●) inside pentagon for this example

- For 2 vertices if magnitude of slope is less than 1, then upper point and otherwise lower point (or whole line).

- Could check other vertex also, but if curvature is within the line already (convex), then no need.

Energy-sum IC extension

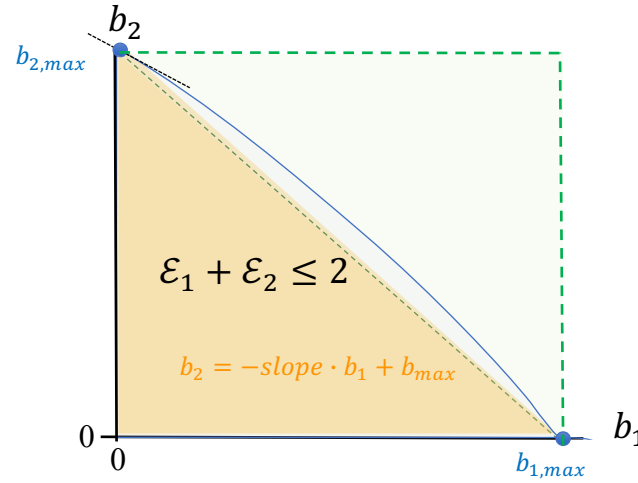
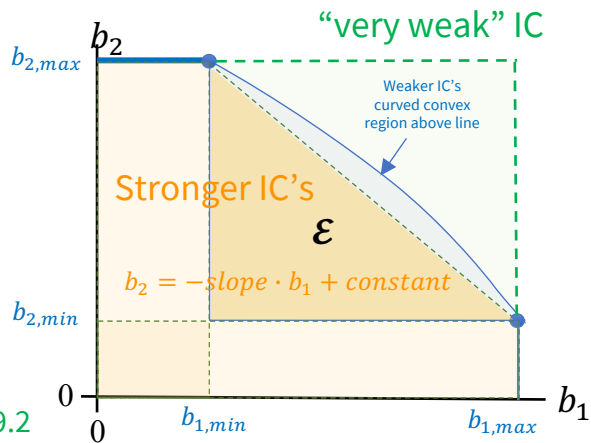
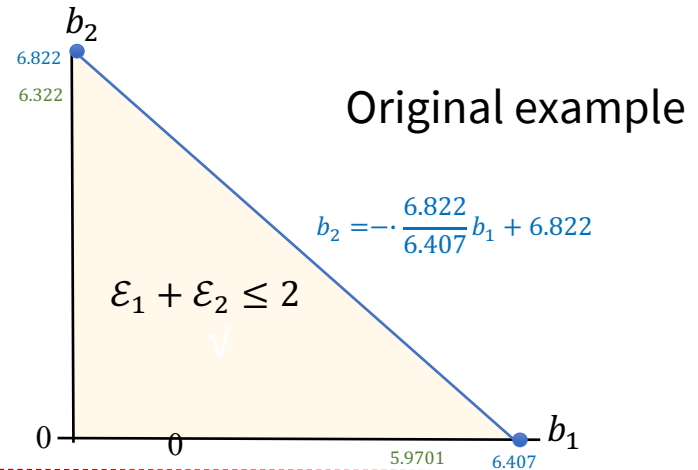
- Same channel with only energy-sum constraint:
- Same derivative test shows dimension-sharing line is boundary

General derivative test

$$\ln(2) \cdot \frac{db_2}{d\mathcal{E}_2} = \frac{g_{22}/2}{g_{22} \cdot \mathcal{E}_2 + 1}$$

$$\ln(2) \cdot \frac{db_1}{d\mathcal{E}_2} = -\frac{g_{22}/2 \cdot g_{12} \cdot \mathcal{E}_1}{g_{22} \cdot (\mathcal{E}_2 + g_{12} \cdot \mathcal{E}_1 + 1) \cdot (g_{22} \cdot \mathcal{E}_2 + 1)} \quad \text{or}$$

$$\ln(2) \cdot \frac{db_1}{d\mathcal{E}_2} = -\frac{g_{21}/2 \cdot g_{11} \cdot \mathcal{E}_1}{g_{21} \cdot (\mathcal{E}_2 + g_{11} \cdot \mathcal{E}_1 + 1) \cdot (g_{21} \cdot \mathcal{E}_2 + 1)},$$



Symmetric 2x2 IC

$$\begin{bmatrix} y_2 \\ y_1 \end{bmatrix} = \begin{bmatrix} 1 & \alpha \\ \alpha & 1 \end{bmatrix} \cdot \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} + \begin{bmatrix} n_2 \\ n_1 \end{bmatrix}$$

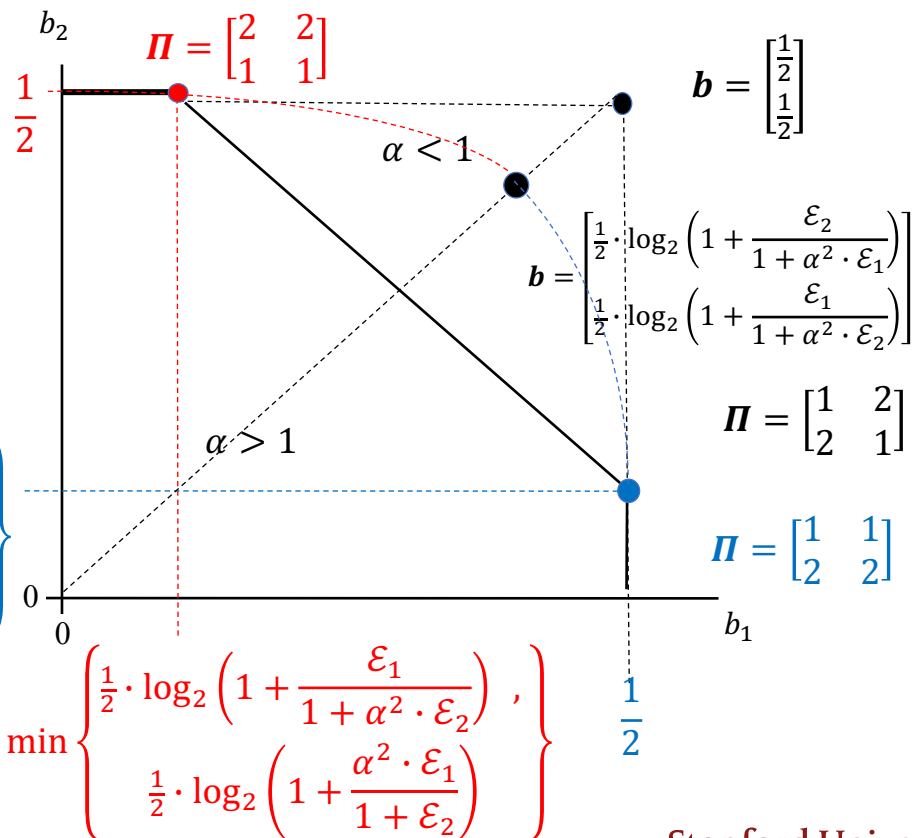
$$R_{nn} = I \text{ and } \mathcal{E} \leq 1$$

- When $\alpha \rightarrow 0$, there is no crosstalk and so $\mathcal{C}_{IC}(\mathbf{b})$ is a square.
- When $\alpha > 1$, $\mathcal{C}_{IC}(\mathbf{b})$ is a pentagon.
- When $0 < \alpha < 1$, $\mathcal{C}_{IC}(\mathbf{b})$ is intermediate

$$\min \left\{ \begin{array}{l} \frac{1}{2} \cdot \log_2 \left(1 + \frac{\mathcal{E}_2}{1 + \alpha^2 \cdot \mathcal{E}_1} \right) \\ \frac{1}{2} \cdot \log_2 \left(1 + \frac{\alpha^2 \cdot \mathcal{E}_2}{1 + \mathcal{E}_1} \right) \end{array} \right\}$$

These two are same for $\alpha = 1$ and equal energy, and determine Imin vector possibilities

Achievable Region when $\mathcal{E}_1 = \mathcal{E}_2 = 1$



Vector Gaussian IC Example, $L_{x,u} \equiv 1$; $L_{y,u} \equiv 2$

▪ 2 users and H is 4×2 : $\mathbf{y} = \begin{bmatrix} \mathbf{y}_2 \\ \mathbf{y}_1 \end{bmatrix} = \begin{bmatrix} H_2 \\ H_1 \end{bmatrix} \cdot \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} + \begin{bmatrix} \mathbf{n}_2 \\ \mathbf{n}_1 \end{bmatrix}$.

$$H_2 = [\mathbf{h}_{22} \quad \mathbf{h}_{21}] = \begin{bmatrix} .9 & .3 \\ .3 & .8 \end{bmatrix}$$

$$R_{nm} = .01 \cdot I$$

$$H_1 = [\mathbf{h}_{12} \quad \mathbf{h}_{11}] = \begin{bmatrix} .8 & .7 \\ .6 & .5 \end{bmatrix}$$

```
>> H2 = [9 3
         3 8];
>> Rb2inv=H2*H2+diag([1 1]);
>> Gbar2=chol(Rb2inv);
>> G2=inv(diag(diag(Gbar2)))*Gbar2;
>> S02=diag(diag(Gbar2))*diag(diag(Gbar2));
>> 0.5*log2(diag(S02)) =
```

b2 = 3.2539

b1 = 2.7526

```
>> H1 = [8 7
         6 5];
```

```
>> Rb1inv=H1*H1+diag([1 1]);
>> Gbar1=chol(Rb1inv);
>> S01=diag(diag(Gbar1))*diag(diag(Gbar1));
>> 0.5*log2(diag(S01)) =
```

b2 = 3.3291

b1 = 0.4128

Two MAC sets

Coincides with $\mathcal{I}_{min} = \begin{bmatrix} 3.25 \\ .413 \end{bmatrix}$

```
>> J2=hankel([0 1]);
>> Rb2inv=J2*H2*H2*J2+diag([1 1]);
         74 51
         51 91
>> Gbar2=chol(Rb2inv);
>> S02=diag(diag(Gbar2))*diag(diag(Gbar2));
>> 0.5*log2(diag(S02)) =
```

b1 = 3.1047

b2 = 2.9018

Two MAC Sets with Reversed order

```
>> Rb1inv=J2*H1*H1*J2+diag([1 1]);
>> Gbar1=chol(Rb1inv);
>> S01=diag(diag(Gbar1))*diag(diag(Gbar1));
>> 0.5*log2(diag(S01)) =
```

b1 = 3.1144

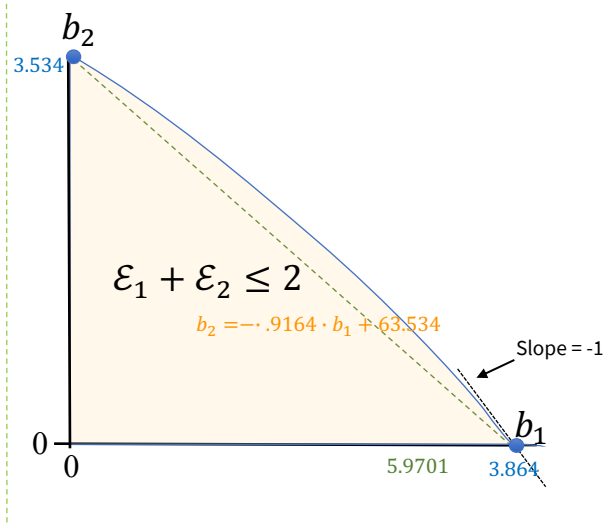
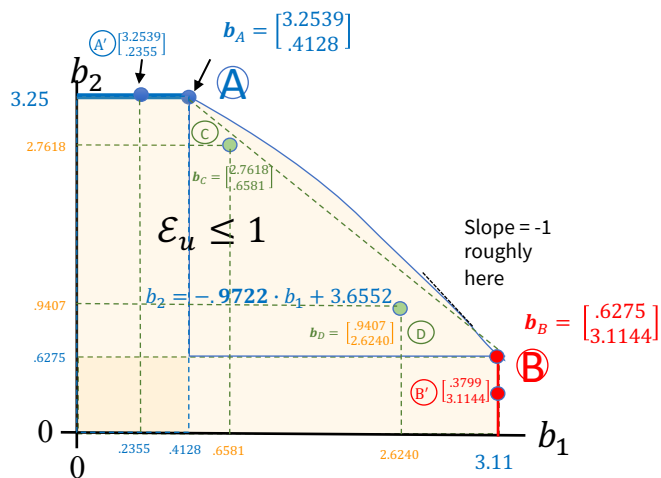
b2 = 0.6275

Coincides with $\mathcal{I}_{min} = \begin{bmatrix} 3.11 \\ .628 \end{bmatrix}$

Try various energy points
With the same orders.



4x2 IC Example continued



- Vary energies near max points to see if local points above or below dimension-sharing, so curved boundary or flat
- Since the line has slope magnitude less than 1, then the curvature is above this line with max rate sum at magnitude 1.

PS 5.4 (2.31) – IC channel has mix, one 2x2 user and one scalar user



Chris Adelico-Ferrarin's MU_IC.m

```
function [b, GU, WU, S0, MSWMFU] = mu_ic(H, A, Lxu, Lyu, cb)
```

Per-tonal (temporal dimension) multiuser interference channel receiver and per-user bits - Chris Adelico Ferrarin - 2023

Inputs: H, A, Lxy, Lyu, cb
Outputs: b, GU, WU, S0, MSWMFU

Definitions:

H: noise-whitened channel matrix [HUU ... HU1] sum-Lyu x sum-Lxu

$$\begin{bmatrix} | & \cdot & \cdot & | \\ |H2U \dots H21| \\ |H1U \dots H11| \end{bmatrix}$$

A: Block Diag sq-root sum-Lxu x sum-Lxu discrete modulators, blkdiag([AU ... A1]); The Au entries derive from each IC user's Lxu x Lxu input autocorrelation matrix, where the trace of each such autocorrelation matrix is user u's energy/symbol. This is per-tone.

Lxu: # of input dimensions for each user U ... 1 in 1 x U row vector
Lyu: # of output dimensions for each user U ... 1 in 1 x U row vector
cb: = 1 if complex baseband or 2 if real baseband channel

GU: unbiased feedback matrix sum-Lxu x sum-Lxu x U with matrices indexed from user U (e.g. GU(:,2) gives GU for user U-1). with matrices indexed from user U (e.g. WU(:,2) gives WU for user U-1).
S0: sub-channel channel gains sum-Lxu x sum-Lxu x U with matrices indexed from user U (e.g. GU(:,2) gives S0 for user U-1).
MSWMFU: unbiased mean-squared whitened matched filter, sum-Lxu x Ly x U with matrices indexed from user U (e.g. MSWMFU(:,2) gives

WU: unbiased feedforward linear equalizer sum-Lxu x sum-Lxu x U MSWMFU for user U-1).

b - user u's bits/symbol 1 x U
the user should recompute b if there is a cyclic prefix

Does not Find I_{min}

- IC needs Lxu and Lyu
- A is like mu_mac
- Per tone
- Arrange input Hu matrices according to order Π

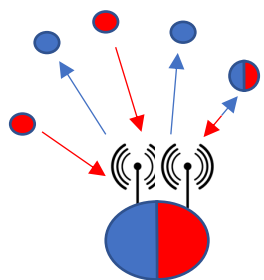
```
H2 = [0.9 0.3; 0.3 0.8]; % From L12:6
H1 = [0.8 0.7; 0.6 0.5];
sigma2 = 0.01;
Ht = [H2; H1] / sqrt(sigma2); % this is 4x2 matrix (2 outputs/input)
A = [1 0; 0 1];
Lxu = [1 1];
Lyu = [2 2];
cb = 2;
>> [b_A, GU_A, WU_A, S0_A, MSWMFU_A] = mu_ic(Ht, A, Lxu, Lyu, cb);
USER 2                USER 1
>> b_A % =
    3.2539            0.4128
GU_A(:,1) =          GU_A(:,2) =
    1.0000    0.5667          1.0000    0.8600
         0    1.0000              0    1.0000
WU_A(:,1) =          WU_A(:,2) =
    0.0111         0            0.0100         0
   -0.0126    0.0225          -1.1026    1.2949
S0_A(:,1) =          S0_A(:,2) =
   91.0000         0            101.0000         0
         0   45.4176              0    1.7723
MSWMFU_A(:,1) =      MSWMFU_A(:,2) =
    0.1000    0.0333          0.0800    0.0600
   -0.0460    0.1423          0.2436   -0.1410
```

Order implied by H's column index

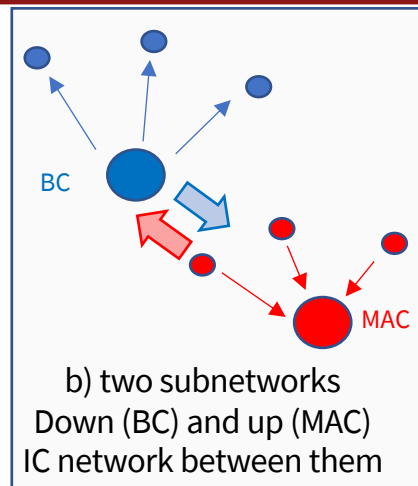


Nesting, DAS, cellfree & Relay

Multiuser Nesting

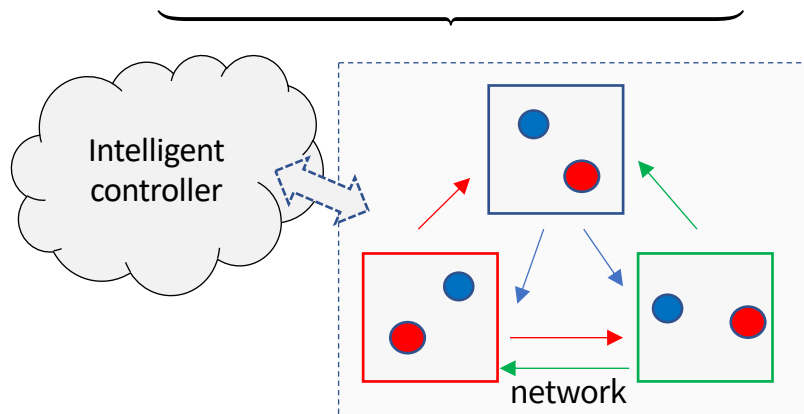


a) radio node edge
(base station or
Access point)



b) two subnetworks
Down (BC) and up (MAC)
IC network between them

Macro User



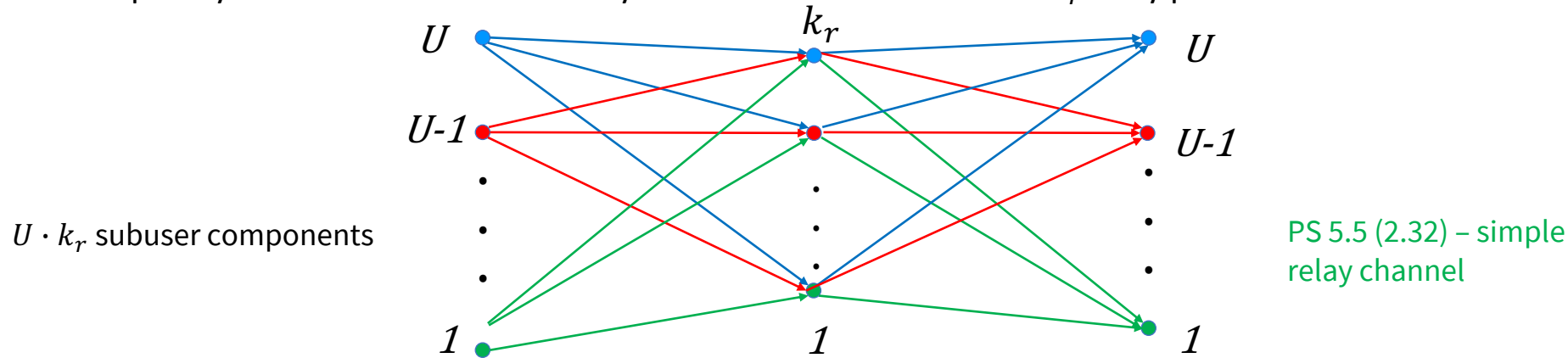
c) three nested IC: (BC, MAC)
Like those in b), nested into
3x3 IC network

**Will be suboptimal,
But easier
Section 5.6.2.3 later**



Single-Stage Relay Channel

- Conceptually uses what we know already and introduces sub-users at k_r relay points



$\mathbf{B}^{(1)}$ is $U \times k_r$ MAC user set

$\mathbf{B}^{(2)}$ is $k_r \times U$ BC user set

$$\mathcal{A}^{(1)}(\mathbf{B}^{(1)}, R_{xx}) = \bigcup_{\Pi^{(1)}}^{conv} \left\{ \mathbf{B}_{min}^{(1)}(\Pi^{(1)}, R_{xx}) \right\}$$

$$\mathcal{A}^{(2)}(\mathbf{B}^{(2)}, R_{xx}) = \bigcup_{\Pi^{(2)}}^{conv} \left\{ \mathbf{B}_{min}^{(2)}(\Pi^{(2)}, R_{xx}) \right\}$$

$$b_u(R_{xx}^{(1)}, R_{xx}^{(2)}) = \left\{ b_u \mid b_u \in \sum_{i=1}^{k_r} \min_{\{\beta_k^{(1)} \in \mathcal{B}_k^{(1)} \wedge \beta_k^{(2)} \in \mathcal{B}_k^{(2)}\}} \left[\beta_k^{(1)}(u, R_{xx}^{(1)}) ; \beta_u^{(2)}(k, R_{xx}^{(2)}) \right] \right\}$$

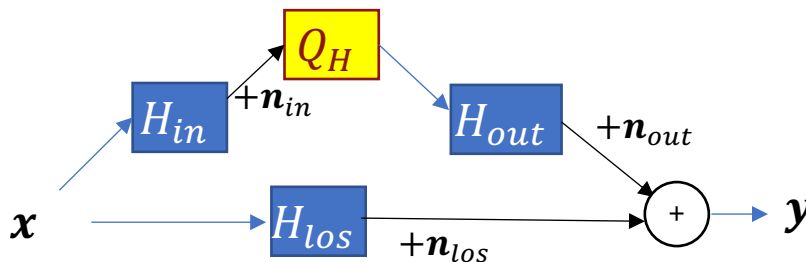
$$\mathcal{C}(\mathbf{b}) = \sum_{u'=1}^U \left[\bigcup_{R_{xx}^{(1)} \wedge R_{xx}^{(2)}}^{hull} \bigcup_{\Pi^{(1)} \otimes \Pi^{(2)}}^{hull} \mathbf{b}(R_{xx}^{(1)}, R_{xx}^{(2)}) \right]_{u, u'}$$

- Multi-stage tedious, but same principles apply recursively

Reflective Intelligent Surfaces (RIS)

Posed Project/Research
"maxRIS" or "minRIS"

$$\mathbf{y} = \underbrace{\begin{bmatrix} H_{los} \\ H_{out} \cdot Q_H \cdot H_{in} \end{bmatrix}}_{H_{RIS}} \cdot \mathbf{x} + \underbrace{\begin{bmatrix} \mathbf{n}_{los} \\ \mathbf{n}_{out} + Q_H \cdot \mathbf{n}_{in} \end{bmatrix}}_{\mathbf{n}_{RIS}}$$



- The RIS matrix Q_H satisfies $\|Q_H\|_F^2 \leq G_H$, the RIS gain – it may also satisfy
 - Q_H is unitary matrix (preserves energy)
 - Q_H is diagonal, and usually also unitary, to be phase/gain-only adjustment on each antenna port (in-to-out)
 - Q_H has individual elements restricted

- For a given R_{xx} , maximize over Q_H $\mathcal{I}(\mathbf{y}; \mathbf{x}) = \log_2 |R_{n,RIS} + H_{RIS} \cdot R_{xx} \cdot H_{RIS}^*|$

- For a given Q_H , maximize the same over R_{xx}

$$R_{nn,RIS} = \begin{bmatrix} R_{nn} & 0 \\ 0 & R_{nn,out} + Q_H \cdot R_{nn,in} \cdot Q_H^* \end{bmatrix}$$

- Iterate? \rightarrow not convex in the Q_H , this needs work.

$$H_{in} \cdot R_{xx} \cdot H_{in}^* + R_{n_{in}} \mathbf{n}_{in} = V_{in} \cdot \Lambda_{in} \cdot V_{in}^*$$

$$R_{out}^{-1/2} \cdot H_{out} \cdot \Lambda_{in} \cdot H_{out}^* \cdot R_{out}^{-*/2} + I = V_{out} \cdot \Lambda_{out} \cdot V_{out}^*$$

- However this instructor postulates that $Q_H^{opt} = V_{out}^* \cdot V_{in}^*$
- where





STANFORD

End Lecture 10